

WIRELESS COMMUNICATIONS

EC412 WIRELESS COMMUNICATIONS & CELLULAR NETWORKS

UNIT – I

INTRODUCTION TO MOBILE COMMUNICATION: Evolution of Mobile Radio Communication, Mobile Radio Telephony in US and around the world, Examples of Wireless Communication Systems: Paging system, Cordless telephones systems, Cellular telephone Systems, Trends in Cellular Radio and personal Communications. The Cellular concept: Frequency reuse, Channel Assignment strategies, Hand off Strategies, Interference and system capacity, improving coverage and capacity in cellular systems.

UNIT – II

MOBILE RADIO PROPAGATION: Large Scale Fading: Introduction, Free space propagation model, Relating power to electric field, The Three basic propagation mechanisms: Reflection, Ground reflection (Two-Ray) model, Diffraction, scattering, Practical Link budget design using path loss models.

Small Scale Fading: Small-scale Multipath Propagation, Impulse response model of a multipath channel, Parameters of mobile multipath channels, Types of small scale fading: Fading effects due to multipath time delay spread and Doppler spread Rayleigh and Ricean distributions.

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Equalization: Fundamentals of equalizers, Equalizers in a communication receiver, Linear equalizers, Nonlinear equalizers: Decision feedback equalizers, Maximum likelihood sequence Estimation (MLSE) equalizer. Diversity Techniques: Space diversity: Selection diversity, feedback, MRC, EGC diversity, Polarization diversity, Frequency diversity, Time diversity, Rake Receiver.

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Multiple Access in Wireless communications: Principle and applications of Multiple Access Techniques FDMA, TDMA, CDMA, Spread Spectrum Multiple Access.

UNIT – V

Wireless Generations Technologies up to 3G:1G, TDMA-based 2G, IS-95, 2.5G, 3G development, Air interface technologies, Internet speeds of 2G, 2.5G, and 3G technologies, Limitations of 3G, Quality of services (QOS) in 3G. 4GTechnology:4G evolution,

TEXT / REFERENCE BOOKS:

- 1) WCY Lee, Mobile Cellular Telecommunications Systems, McGraw Hill, 1990.
 - 2) WCY Lee, Mobile Communications Design Fundamentals, Prentice Hall, 1993.
 - 3) Raymond Steele, Mobile Radio Communications, IEEE Press, New York, 1992.
 - 4) AJ Viterbi, CDMA: Principles of Spread Spectrum Communications, Addison Wesley, 1995.
 - 5) VK Garg &JE Wilkes, Wireless & Personal Communication Systems, Prentice Hall, 1996.
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UNIT-1

Introduction to Wireless Communication Systems

The ability to communicate with people on the move has evolved remarkably since Guglielmo Marconi first demonstrated radio's ability to provide continuous contact with ships sailing the English Channel in 1897. Since then new wireless communications methods and services have been enthusiastically adopted by people throughout the world. Particularly during the past ten years, the mobile radio communications industry has grown by orders of magnitude, fueled by digital and RF circuit fabrication improvements, new large-scale circuit integration, and other miniaturization technologies which make portable radio equipment smaller, cheaper, and more reliable. Digital switching techniques have facilitated the large scale deployment of affordable, easy-to-use radio communication networks.

Evolution of Mobile Radio Communications

The ability to provide wireless communications to an entire population was not even conceived until Bell Laboratories developed the cellular concept in the 1960s and 1970s. With the development of highly reliable, miniature, solid-state radio frequency hardware in the 1970s, the wireless communications era was born. The following market penetration data show how wireless communications in the consumer sector has grown in popularity. Figure 1.1 illustrates how mobile telephony has penetrated our daily lives compared with other popular inventions of the 20th century. Figure 1.1 shows that the first 35 years of mobile telephony saw little market penetration due to high cost and the technological challenges involved, but however, in the past decade, wireless communications has been accepted by consumers at rates comparable to television and the video cassette recorder.

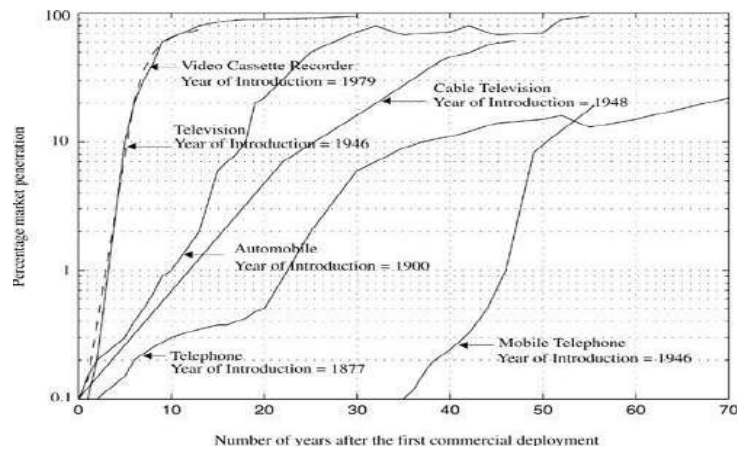


Figure 1.1. The growth of mobile telephony as compared with other popular inventions of the 20th century.

- By 1934, 194 municipal police radio systems and 58 state police stations had adopted amplitude modulation (AM) mobile communication systems for public safety in the U.S.
- In 1935, Edwin Armstrong demonstrated frequency modulation (FM) for the first time, and since the late 1930s, FM has been the primary modulation technique used for mobile communication systems throughout the world.
- With the boom in CB radio and cordless appliances such as garage door openers and telephones, the number of users of mobile and portable radio in 1995 was about 100 million, or 37% of the U.S. population
- The number of worldwide cellular telephone users grew from 25,000 in 1984 to about 25 million in 1993, and since then subscription-based wireless services have been experiencing customer growth rates well in excess of 50% per year. At the beginning of the 21st century, over 1% of the worldwide wireless subscriber population had already abandoned wired telephone service for home use, and had begun to rely solely on their cellular service provider for telephone access.

Mobile Radiotelephony in the U.S.

In 1946, the first public mobile telephone service was introduced in twenty-five major American cities. Each system used a single, high-powered transmitter and a large tower in order to cover distances of over 50 km in a particular market. During the 1950s and 1960s, AT&T Bell Laboratories and other telecommunications companies throughout the world developed the theory and techniques of cellular radiotelephony—the concept of breaking a coverage zone (market) into small cells, each of which reuses portions of the spectrum to increase spectrum usage at the expense of greater system infrastructure. AT&T proposed the concept of a cellular mobile system to the FCC in 1968, although technology was not available to implement cellular telephony until the late 1970s. In 1983, the FCC finally allocated 666 duplex channels (40 MHz of spectrum in the 800 MHz band, each channel having a one-way bandwidth of 30 kHz for a total spectrum occupancy of 60 kHz for each duplex channel) for the U.S. Advanced Mobile Phone System (AMPS).

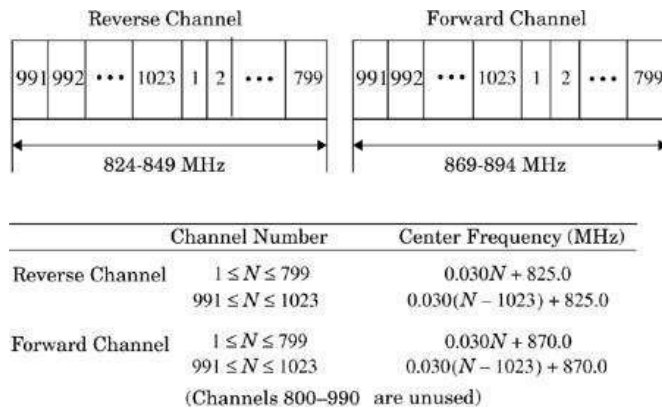


Figure 1.2. Frequency spectrum allocation for the U.S. cellular radio service

In late 1991, the first US Digital Cellular (USDC) system hardware was installed in major U.S. cities. The USDC standard (Electronic Industry Association Interim Standard IS-54 and later IS-136) allowed cellular operators to replace gracefully some single-user analog channels with digital channels which support three users in the same 30 kHz bandwidth. In this way, U.S. carriers gradually phased out AMPS as more users accepted digital phones.

A cellular system based on code division multiple access (CDMA) has been developed by Qualcomm, Inc. and standardized by the Telecommunications Industry Association (TIA) as an Interim Standard (IS-95). This system supports a variable number of users in 1.25 MHz wide channels using direct sequence spread spectrum. CDMA systems can operate at much larger interference levels because of their inherent interference resistance properties. The ability of CDMA to operate with a much smaller signal-to-noise ratio (SNR) than conventional narrowband FM techniques allows CDMA systems to use the same set of frequencies in every cell, which provides a large improvement in capacity.

Personal Communication Service (PCS) licenses in the 1800/1900 MHz band were auctioned by the U.S. Government to wireless providers in early 1995, and these have spawned new wireless services that complement, as well as compete with, cellular and SMR.

Mobile Radio Systems Around the World

Many mobile radio standards have been developed for wireless systems throughout the world, and more standards are likely to emerge. Tables 1.1 through 1.3 list the most common paging, cordless, cellular, and personal communications standards used in North America, Europe, and Japan.

Table 1.1. Major Mobile Radio Standards in North America

Standard	Type	Year Introduction	Multiple Access	Frequency Band	Channel Modulation	Bandwidth
AMPS	Cellular	1983	FDMA	824-894 MHz	FM	30 kHz
NAMPS	Cellular	1992	FDMA	824-894 MHz	FM	10 kHz
USDC	Cellular	1991	TDMA	824-894 MHz	$\pi/4$ -DQPSK	30 kHz
CDPD	Cellular	1993	FH/Packet	824-894 MHz	GMSK	30 kHz
IS-95	Cellular/PCS	1993	CDMA	824-894 MHz	QPSK/BPSK	1.25 MHz
GSC	Paging	1970s	Simplex	Several	FSK	12.5 kHz
POCSAG	Paging	1970s	Simplex	Several	FSK	12.5 kHz
FLEX	Paging	1993	Simplex	Several	4-FSK	15 kHz
DCS-1900 (GSM)	PCS	1994	TDMA	1.85-1.99 GHz	GMSK	200 kHz
PACS	Cordless/PCS	1994	TDMA/FDMA	1.85-1.99 GHz	$\pi/4$ -DQPSK	300 kHz
MIRS	SMR/PCS	1994	TDMA	Several	16-QAM	25 kHz
iDen	SMR/PCS	1995	TDMA	Several	16-QAM	25 kHz

Table 1.2. Major Mobile Radio Standards in Europe

Standard	Type	Year Introduction	Multiple Access	Frequency Band	Channel Modulation	Bandwidth
ETACS	Cellular	1985	FDMA	900 MHz	FM	25 kHz
NMT-450	Cellular	1981	FDMA	450-470 MHz	FM	25 kHz
NMT-900	Cellular	1986	FDMA	890-960 MHz	FM	12.5 kHz
GSM	Cellular/PCS	1990	TDMA	890-960 MHz	GMSK	200 kHz
C-450	Cellular	1985	FDMA	450-465 MHz	FM	20 kHz/10 kHz

Standard	Type	Year of Introduction	Multiple Access	Frequency Band	Modulation	Channel Bandwidth
ERMES	Paging	1993	FDMA	Several	4-FSK	25kHz
CT2	Cordless	1989	FDMA	864-868 MHz	GFSK	100kHz
DECT	Cordless	1993	TDMA	1880-1900MHz	GFSK	1.728 MHz
DCS-1800	Cordless/PCS	1993	TDMA	1710-1880MHz	GMSK	200kHz

Table 1.3. Major Mobile Radio Standards in Japan

Standard	Type	Year of Introduction	Multiple Access	Frequency Band	Modulation	Channel Bandwidth
JTACS	Cellular	1988	FDMA	860-925MHz	FM	25kHz
PDC	Cellular	1993	TDMA	810-1501 MHz	$\pi/4$ -DQPSK	25kHz
NTT	Cellular	1979	FDMA	400/800 MHz	FM	25kHz
NTACS	Cellular	1993	FDMA	843-925MHz	FM	12.5kHz
NTT	Paging	1979	FDMA	280MHz	FSK	12.5kHz
NEC	Paging	1979	FDMA	Several	FSK	10kHz
PHS	Cordless	1993	TDMA	1895-1907MHz	$\pi/4$ -DQPSK	300kHz

The world's first cellular system was implemented by the Nippon Telephone and Telegraph company (NTT) in Japan. The system, deployed in 1979, uses 600 FM duplex channels (25 kHz for each one-way link) in the 800 MHz band. In Europe, the Nordic Mobile Telephone system (NMT 450) was developed in 1981 for the 450 MHz band and uses 25 kHz channels. The European Total Access Cellular System (ETACS) was deployed in 1985 and is virtually identical to the U.S. AMPS system, except that the smaller bandwidth channels result in a slight degradation of signal-to-noise ratio (SNR) and coverage range. In Germany, a cellular standard called C-450 was introduced in 1985. The first generation European cellular systems are generally incompatible with one another because of the different frequencies and communication protocols used. These systems are now being replaced by the Pan European digital cellular standard GSM (Global System for Mobile) which was first deployed in 1990 in a new 900 MHz band which all of Europe dedicated for cellular telephone service. The GSM standard has gained worldwide acceptance as the first universal digital cellular system with modern network features extended to each mobile user, and it is the leading digital air interface for PCS services above 1800 MHz throughout the world. In Japan, the Pacific Digital Cellular (PDC) standard provides digital cellular coverage using a system similar to North America's USDC.

Examples of Wireless Communication Systems

Most people are familiar with a number of mobile radio communication systems used in everyday life. Garage door openers, remote controllers for home entertainment equipment, cordless telephones, hand-held walkie-talkies, pagers (also called paging receivers or "beepers"), and cellular telephones are all examples of mobile radio communication systems. However, the cost, complexity, performance, and types of services offered by each of these mobile systems are vastly different.

Table 1.4 lists definitions of terms used to describe elements of wireless communication systems.

Table 1.4. Wireless Communications System Definitions

Base Station	A fixed station in a mobile radio system used for radio communication with mobile stations. Base stations are located at the center or on the edge of a coverage region and consist of radio channels and transmitter and receiver antennas mounted on a tower.
Control Channel	Radio channel used for transmission of call setup, call request, call initiation, and other beacon or control purposes.
Forward Channel	Radio channel used for transmission of information from the base station to the mobile.
Full Duplex Systems	Communication systems which allow simultaneous two-way communication. Transmission and reception is typically on two different channels (FDD) although new cordless/PCS systems are using TDD.
Half Duplex Systems	Communication systems which allow two-way communication by using the same radio channel for both transmission and reception. At any given time, the user can only either transmit or receive information.
Handoff	The process of transferring a mobile station from one channel or base station to another.
Mobile Station	A station in the cellular radio service intended for use while in motion at unspecified locations. Mobile stations may be hand-held personal units (portables) or installed in vehicles (mobiles).
Mobile Switching Center	Switching center which coordinates the routing of calls in a large service area. In a cellular radio system, the MSC connects the cellular base stations and the mobiles to the PSTN. An MSC is also called a mobile telephone switching office (MTSO).
Page	A brief message which is broadcast over the entire service area, usually in a simulcast fashion by many base stations at the same time.

ReverseChannel	Radiochannelusedfortransmissionofinformationfromthemobiletobasestation.
Roamer	A mobilestationwhichoperatesina servicearea (market)other thanthatfromwhich service has been subscribed.
SimplexSystems	Communication systemswhichprovideonlyone-waycommunication.
Subscriber	Auserwhopayssubscriptionchargesforusinga mobilecommunicationssystem. A
Transceiver	device capable of simultaneously transmitting and receiving radio signals.

Mobile radio transmission systems may be classified as *simplex*, *half-duplex* or *full-duplex*. In simplex systems, communication is possible in only one direction. Paging systems, in which messages are received but not acknowledged, are simplex systems. Half-duplex radio systems allow two-way communication, but use the same radio channel for both transmission and reception. This means that at any given time, a user can only transmit or receive information. Constraints like “push-to-talk” and “release-to-listen” are fundamental features of half-duplex systems. Full duplex systems, on the other hand, allow simultaneous radio transmission and reception between a subscriber and a base station, by providing two simultaneous but separate channels (frequency division duplex, or FDD) or adjacent time slots on a single radio channel (time division duplex, or TDD) for communication to and from the user.

In FDD, a pair of simplex channels with a fixed and known frequency separation is used to define a specific radiochannelintheprogram. The channelused toconvey traffic tothemobileuser from a basestationiscalled the*forward channel*, whilethe channelusedtocarry traffic from themobileusertoa basestationis called the *reverse channel*. In the U.S. AMPS standard,the reverse channel has a frequency which is exactly 45 MHz lower than that of the forward channel. Full duplex mobile radio systems provide many of the capabilities of the standard telephone, with the added convenience of mobility. Full duplex and half-duplex systems use *transceivers* for radio communication. FDD is used exclusively in analog mobile radio systems.

Time division duplexing (TDD) uses the fact that it is possible to share a single radio channel in time, so that a portion of the time is used to transmit from the base station to the mobile, and the remaining time is used to transmit from the mobile to the base station. If the data transmission rate in the channel is muchgreater than the end-user’s data rate, it is possible to store information bursts and provide the appearance of full duplex operation to a user, even though there are *not* two simultaneous radio transmissions at any instant. TDD is only possible with digital transmissionformatsand digital modulation, andis very sensitive to timing. It isfor this reason that TDD has only recently been used, and only for indoor or small area wireless applications where the physical coverage distances (and thus the radio propagation time delay)are much smaller than the many kilometers used in conventional cellular telephone systems.

PagingSystems

Paging systems are communication systems that send brief messages to a subscriber. Depending on the type of service, the message may be either a numeric message, an alphanumeric message, or a voice message. Paging systems are typically used to notify a subscriber of the need to call a particular telephone number or travel to a known location to receive further instructions. In modern paging systems, news headlines, stock quotations, and faxes may be sent. A message is sent to a paging subscriber via the paging system access number (usually a toll-free telephone number) with a telephone keypad or modem. The issued message is called a *page*. The paging system then transmits the page throughout the service area using base stations which broadcast the page on a radio carrier.

Paging systems vary widely in their complexity and coverage area. While simple paging systems may cover a limited range of 2 to 5 km, or may even be confined to within individual buildings, wide area paging systems can provide worldwide coverage. Though paging receivers are simple and inexpensive, the transmission system required is quite sophisticated. Wide area paging systems consist of a network of telephone lines, many base station transmitters, and large radio towers that simultaneously broadcast a page from each base station (this is called *simulcasting*). Simulcast transmitters may be located within the same service area or in different cities or countries. Paging systems are designed to provide reliable communication to subscribers wherever they are; whether inside a building, driving on a highway, or flying in an airplane. This necessitates large transmitter powers (on the order of kilowatts) and low data rates (a couple ofthousand bits persecond) for maximum coverage from each base station. [Figure 1.3](#) shows a diagram of a wide area paging system.

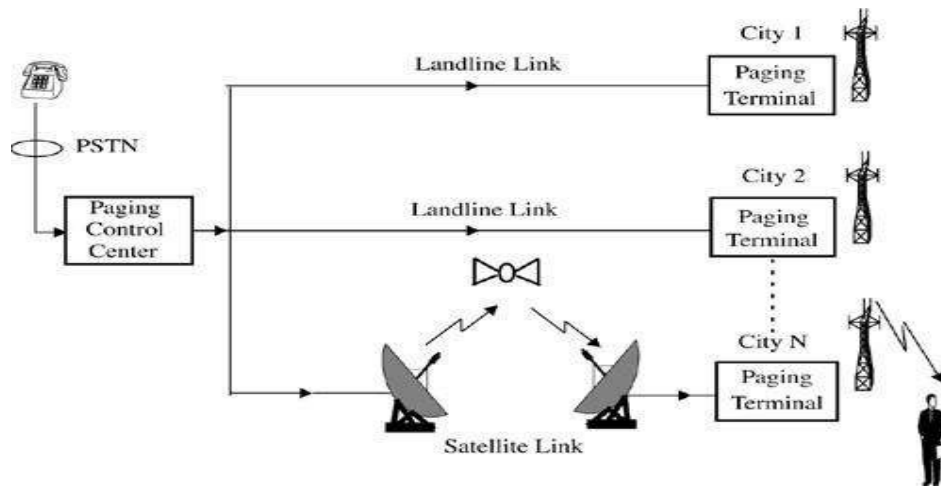


Figure 1.3. A wide area paging system. The paging control center dispatches pages received from the PSTN throughout several cities at the same time.

Paging systems are designed to provide ultra-reliable coverage, even inside buildings. Buildings can attenuate radio signals by 20 or 30 dB, making the choice of base station locations difficult for the paging companies. For this reason, paging transmitters are usually located on tall buildings in the center of a city, and simulcasting is used in conjunction with additional base stations located on the perimeter of the city to flood the entire area. Small RF bandwidths are used to maximize the signal-to-noise ratio at each paging receiver, so low data rates (6400 bps or less) are used.

Cordless Telephone Systems

Cordless telephone systems are full duplex communication systems that use radio to connect a portable handset to a dedicated base station, which is then connected to a dedicated telephone line with a specific telephone number on the public switched telephone network (PSTN). In first generation cordless telephone systems (manufactured in the 1980s), the portable unit communicates only to the dedicated base unit and only over distances of a few tens of meters. Early cordless telephones operate solely as extension telephones to a transceiver connected to a subscriber line on the PSTN and are primarily for in-home use.

Second generation cordless telephones have recently been introduced which allow subscribers to use their handsets at many outdoor locations within urban centers such as London or Hong Kong. Modern cordless telephones are sometimes combined with paging receivers so that a subscriber may first be paged and then respond to the page using the cordless telephone. Cordless telephone systems provide the user with limited range and mobility, as it is usually not possible to maintain a call if the user travels outside the range of the base station. Typical second generation base stations provide coverage ranges up to a few hundred meters. [Figure 1.4](#) illustrates a cordless telephone system.

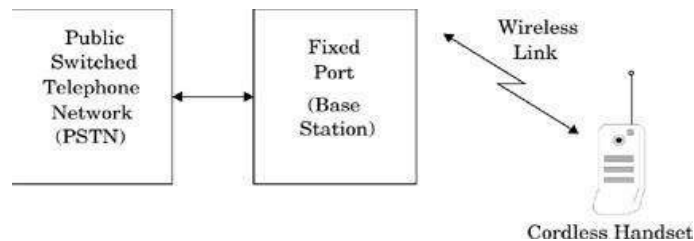


Figure 1.4. A cordless telephone system.

Cellular Telephone Systems

A cellular telephone system provides a wireless connection to the PSTN for any user location within the radio range of the system. Cellular systems accommodate a large number of users over a large geographic area, within a limited frequency spectrum. Cellular radio systems provide high quality service that is often comparable to that of the landline telephone systems. High capacity is achieved by limiting the coverage of each base station transmitter to a small geographic area called a *cell* so that the same radio channels may be reused by another base station located some distance away. A sophisticated switching technique called a *handoff* enables a call to proceed uninterrupted when the user moves from one cell to another.

[Figure 1.5](#) shows a basic cellular system which consists of *mobile stations*, *base stations* and a *mobile switching center* (MSC). The mobile switching center is sometimes called a *mobile telephone switching office* (MTSO), since it is responsible for connecting all mobiles to the PSTN in a cellular system. Each mobile communicates

via radio with one of the base stations and may be handed-off to any number of base stations throughout the duration of a call. The mobile station contains a transceiver, an antenna, and control circuitry, and may be mounted in a vehicle or used as a portable hand-held unit. The base stations consist of several transmitters and receivers which simultaneously handle full duplex communications and generally have towers which support several transmitting and receiving antennas. The base station serves as a bridge between all mobile users in the cell and connects the simultaneous mobile calls via telephone lines or microwave links to the MSC. The MSC coordinates the activities of all of the base stations and connects the entire cellular system to the PSTN. A typical MSC handles 100,000 cellular subscribers and 5,000 simultaneous conversations at a time, and accommodates all billing and system maintenance functions, as well. In large cities, several MSCs are used by a single carrier.

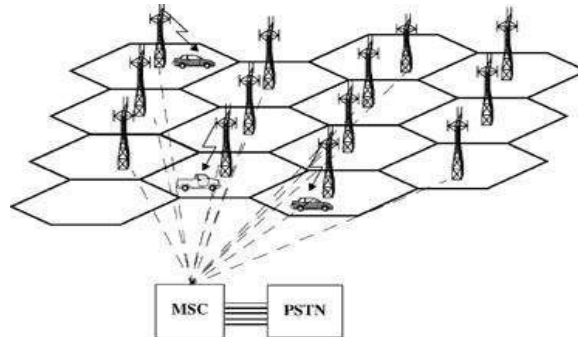


Figure 1.5. A cellular system. The towers represent base stations which provide radio access between mobile users and the mobile switching center (MSC).

Communication between the base station and the mobile is defined by a standard *common air interface* (CAI) that specifies four different channels. The channels used for voice transmission from the base station to mobiles are called *forward voice channels* (FVC), and the channels used for voice transmission from mobiles to the base station are called *reverse voice channels* (RVC). The two channels responsible for initiating mobile calls are the *forward control channels* (FCC) and *reverse control channels* (RCC). Control channels are often called *setup channels* because they are only involved in setting up a call and moving it to an unused voice channel. Control channels transmit and receive data messages that carry call initiation and service requests, and are monitored by mobiles when they do not have a call in progress. Forward control channels also serve as beacons which continually broadcast all of the traffic requests for all mobiles in the system.

Cellular systems rely on the frequency reuse concept, which requires that the forward control channels (FCCs) in neighboring cells be different. By defining a relatively small number of FCCs as part of the common air interface, cellular phones can be manufactured by many companies which can rapidly scan all of the possible FCCs to determine the strongest channel at any time. Once finding the strongest signal, the cellular phone receiver stays "camped" to the particular FCC. By broadcasting the same setup data on all FCCs at the same time, the MSC is able to signal all subscribers within the cellular system and can be certain that any mobile will be signaled when it receives a call via the PSTN.

How a Cellular Telephone Call is Made

When a cellular phone is turned on, but is not yet engaged in a call, it first scans the group of forward control channels to determine the one with the strongest signal, and then monitors that control channel until the signal drops below a usable level. At this point, it again scans the control channels in search of the strongest base station signal. When a telephone call is placed to a mobile user, the MSC dispatches the request to all base stations in the cellular system. The *mobile identification number* (MIN), which is the subscriber's telephone number, is then broadcast as a paging message over all of the forward control channels throughout the cellular system. The mobile receives the paging message sent by the base station which it monitors, and responds by identifying itself over the reverse control channel. The base station relays the acknowledgment sent by the mobile and informs the MSC of the handshake. Then, the MSC instructs the base station to move the call to an unused voice channel within the cell (typically, between ten to sixty voice channels and just one control channel are used in each cell's base station). At this point, the base station signals the mobile to change frequencies to an unused forward and reverse voice channel pair, at which point another data message (called an *alert*) is transmitted over the forward voice channel to instruct the mobile telephone to ring, thereby instructing the mobile user to answer the phone. [Figure 1.6](#) shows the sequence of events involved with connecting a call to a mobile user in a cellular telephone system. All of these events occur within a few seconds and are not noticeable by the user.

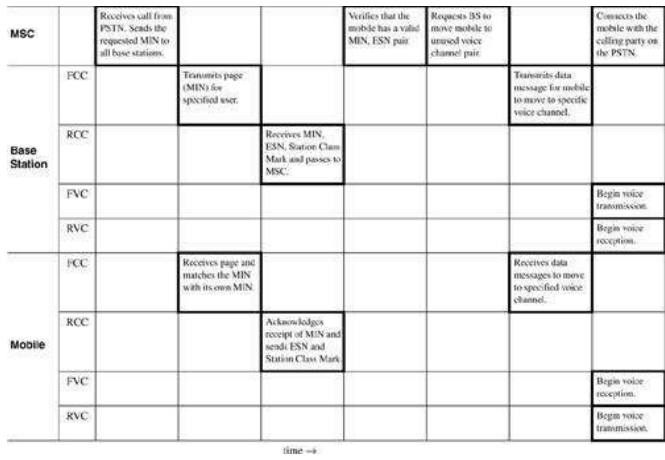


Figure 1.6. Timing diagram illustrating how a call to a mobile user initiated by a landline subscriber is established.

Once a call is in progress, the MSC adjusts the transmitted power of the mobile and changes the channel of the mobile unit and base stations in order to maintain call quality as the subscriber moves in and out of range of each base station. This is called a *handoff*. Special control signaling is applied to the voice channels so that the mobile unit may be controlled by the base station and the MSC while a call is in progress.

When a mobile originates a call, a call initiation request is sent on the reverse control channel. With this request the mobile unit transmits its telephone number (MIN), *electronic serial number* (ESN), and the telephone number of the called party. The mobile also transmits a *station class mark* (SCM) which indicates what the maximum transmitter power level is for the particular user. The cell base station receives this data and sends it to the MSC. The MSC validates the request, makes connection to the called party through the PSTN, and instructs the base station and mobile user to move to an unused forward and reverse voice channel pair to allow the conversation to begin. [Figure 1.7](#) shows the sequence of events involved with connecting a call which is initiated by a mobile user in a cellular system.

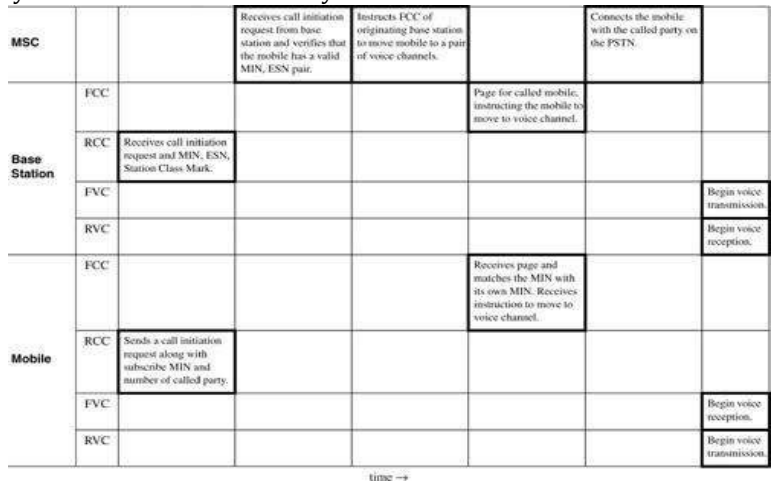


Figure 1.7. Timing diagram illustrating how a call initiated by a mobile is established.

All cellular systems provide a service called *roaming*. This allows subscribers to operate in service areas other than the one from which service is subscribed. When a mobile enters a city or geographic area that is different from its home service area, it is registered as a roamer in the new service area. If a particular roamer has roaming authorization for billing purposes, the MSC registers the subscriber as a valid roamer. Once registered, roaming mobiles are allowed to receive and place calls from that area, and billing is routed automatically to the subscriber's home service provider.

Comparison of Common Wireless Communication Systems

[Tables 1.5](#) and [1.6](#) illustrate the types of service, level of infrastructure, cost, and complexity required for the subscriber segment and base station segment of each of the five mobile or portable radio systems discussed earlier in this chapter. For comparison purposes, common household wireless remote devices are shown in the table. It is important to note that each of the five mobile radio systems given in [Tables 1.5](#) and [1.6](#) use a fixed base station, and for good reason. Virtually all mobile radio communication systems strive to connect a moving terminal to a fixed distribution system of some sort and attempt to look invisible to the distribution system.

Table 1.5. Comparison of Mobile Communication Systems—Mobile Station

Service	Coverage Range	Required Infrastructure	Complexity	Hardware Cost	Carrier Frequency	Functionality
TV Remote Control	Low	Low	Low	Low	Infrared	Transmitter
Garage Door Opener	Low	Low	Low	Low	<100MHz	Transmitter
Paging System	High	High	Low	Low	<1GHz	Receiver
Cordless Phone	Low	Low	Moderate	Low	<1GHz	Transceiver
Cellular Phone	High	High	High	Moderate	<2GHz	Transceiver

Table 1.6. Comparison of Mobile Communication Systems—Base Station

Service	Coverage Range	Required Infrastructure	Complexity	Hardware Cost	Carrier Frequency	Functionality
TV Remote Control	Low	Low	Low	Low	Infrared	Receiver
Garage Door Opener	Low	Low	Low	Low	<100MHz	Receiver
Paging System	High	High	High	High	<1GHz	Transmitter
Cordless Phone	Low	Low	Low	Moderate	<1GHz	Transceiver
Cellular Phone	High	High	High	High	<2GHz	Transceiver

Notice that the expectations vary widely among the services, and the infrastructure costs are dependent upon the required coverage area. For the case of low power, hand-held cellular phones, a large number of base stations are required to insure that any phone is in close range to a base station within a city. If base stations were not within close range, a great deal of transmitter power would be required of the phone, thus limiting the battery life and rendering the service useless for hand-held users.

Trends in Cellular Radio and Personal Communications

Since 1989, there has been enormous activity throughout the world to develop personal wireless systems that combine the network intelligence of today's PSTN with modern digital signal processing and RF technology. The concept, called Personal Communication Services (PCS), originated in the United Kingdom when three companies were given spectrum in the 1800 MHz range to develop Personal Communication Networks (PCN) throughout Great Britain. PCN was seen by the U.K. as a means of improving its international competitiveness in the wireless field while developing new wireless systems and services for citizens.

Indoor wireless networking products are rapidly emerging and promise to become a major part of the telecommunications infrastructure within the next decade. An international standards body, IEEE 802.11, is developing standards for wireless access between computers inside buildings. The European Telecommunications Standard Institute (ETSI) is also developing the 20 Mbps HIPERLAN standard for indoor wireless networks. Products have emerged that allow users to link their phone with their computer within an office environment, as well as in a public setting, such as an airport or train station.

A worldwide standard, the Future Public Land Mobile Telephone System (FPLMTS)—renamed International Mobile Telecommunication 2000 (IMT-2000) in mid-1995—has been formulated by the International Telecommunications Union (ITU) which is the standards body for the United Nations, with headquarters in Geneva, Switzerland. FPLMTS (now IMT-2000) is a third generation universal, multi-function, globally compatible digital mobile radio system that will integrate paging, cordless, and cellular systems, as well as low earth orbit (LEO) satellites, into one universal mobile system.

In emerging nations, where existing telephone service is almost nonexistent, fixed cellular telephone systems are being installed at a rapid rate. This is due to the fact that developing nations are finding it is quicker and more affordable to install cellular telephone systems for fixed home use, rather than install wires in neighborhoods which have not yet received telephone connections to the PSTN.

Modern Wireless Communication Systems

Since the mid 1990s, the cellular communications industry has witnessed explosive growth. Wireless communications networks have become much more pervasive than anyone could have imagined when the cellular concept was first developed in the 1960s and 1970s. The widespread adoption of wireless communications was accelerated in the mid 1990s, when governments throughout the world provided increased competition and new radio spectrum licenses for personal communications services (PCS) in the 1800–2000 MHz frequency bands.

The rapid worldwide growth in cellular telephone subscribers has demonstrated conclusively that wireless communications is a robust, viable voice and data transport mechanism. New standards and technologies are being implemented to allow wireless networks to replace fiber optic or copper lines between fixed points several kilometers apart (*fixed wireless access*). Similarly, wireless networks have been increasingly used as a replacement for wires within homes, buildings, and office settings through the deployment of *wireless local area networks* (WLANs). The evolving *Bluetooth* modem standard promises to replace troublesome appliance communication cords with invisible wireless connections within a person's personal workspace. Used primarily within buildings, WLANs and Bluetooth use low power levels and generally do not require a license for spectrum use.

Second Generation (2G) Cellular Networks

Most of today's ubiquitous cellular networks use what is commonly called *second generation* or *2G* technologies which conform to the second generation cellular standards. Unlike first generation cellular systems that relied exclusively on FDMA/FDD and analog FM, second generation standards use digital modulation formats and TDMA/FDD and CDMA/FDD multiple access techniques.

The most popular second generation standards include three TDMA standards and one CDMA standard: (a) *Global System Mobile (GSM)*, which supports eight time slotted users for each 200 kHz radio channel and has been deployed widely by service providers in Europe, Asia, Australia, South America, and some parts of the US (in the PCS spectrum band only); (b) *Interim Standard 136 (IS-136)*, also known as North American Digital Cellular (NADC), which supports three time slotted users for each 30 kHz radio channel and is a popular choice for carriers in North America, South America, and Australia (in both the cellular and PCS bands); (c) *Pacific Digital Cellular (PDC)*, a Japanese TDMA standard that is similar to IS-136 with more than 50 million users; and (d) the popular 2G CDMA standard *Interim Standard 95 Code Division Multiple Access (IS-95)*, also known as *cdmaOne*, which supports up to 64 users that are orthogonally coded and simultaneously transmitted on each 1.25 MHz channel. CDMA is widely deployed by carriers in North America (in both cellular and PCS bands), as well as in Korea, Japan, China, South America, and Australia.

In many countries, 2G wireless networks are designed and deployed for conventional mobile telephone service, as a high capacity replacement for, or in competition with, existing older first generation cellular telephone systems. Modern cellular systems are also being installed to provide fixed (non-mobile) telephone service to residences and businesses in developing nations—this is particularly cost effective for providing *plain old telephone service (POTS)* in countries that have poor telecommunications infrastructure and are unable to afford the installation of copper wire to all homes.

Table 2.1. Key Specifications of Leading 2G Technologies (adapted from [Lib99])

	cdmaOne, IS-95, ANSI J	GSM, DCS-1900, ANSI J-STD 007	NADC, IS-54/IS-136, ANSI J STD-011, PDC
Uplink Frequencies	824-849 MHz (US Cellular) 1850-1910 MHz (US PCS)	890-915 MHz (Europe) 1850-1910 MHz (US PCS)	800 MHz, 1500 MHz (Japan) 1850-1910 MHz (US PCS)
Downlink Frequencies	869-894 MHz (US Cellular) 1930-1990 MHz (US PCS)	935-960 MHz (Europe) 1930-1990 MHz (US PCS)	869-894 MHz (US Cellular) 1930-1990 MHz (US PCS) 800 MHz, 1500 MHz (Japan)
Duplexing	FDD	FDD	FDD
Multiple Access Technology	CDMA	TDMA	TDMA
Modulation	BPSK with Quadrature Spreading	GMSK with $BT=0.3$	$\pi/4$ DQPSK
Carrier Separation	1.25 MHz	200 kHz	30 kHz (IS-136) (25 kHz for PDC)
Channel Data Rate	1.2288 Mchips/sec	270.833 kbps	48.6 kbps (IS-136) (42 kbps for PDC)
Voice channels per carrier	64	8	3
Speech Coding	Code Excited Linear Prediction (CELP) @ 13 kbps, Enhanced Variable Rate Codec (EVRC) @ 8 kbps	Residual Pulse Excited Long Term Prediction (RPE-LTP) @ 13 kbps	Vector Sum Excited Linear Predictive Coder (VSELP) @ 7.95 kbps

Evolution to 2.5G Mobile Radio Networks

Since the mid 1990s, the 2G digital standards have been widely deployed by wireless carriers for cellular and PCS, even though these standards were designed before the widespread use of the Internet. Consequently, 2G

technologies use circuit-switched data modems that limit data users to a single circuit-switched voice channel.

In an effort to retrofit the 2G standards for compatibility with increased throughput data rates that are required to support modern Internet applications, new data-centric standards have been developed that can be overlaid upon existing 2G technologies. These new standards represent 2.5G technology and allow existing 2G equipment to be modified and supplemented with new base station add-ons and subscriber unit software upgrades to support higher data rate transmissions for web browsing, e-mail traffic, mobile commerce (m-commerce), and location-based mobile services. The 2.5G technologies also support a popular new web browsing format language, called Wireless Applications Protocol (WAP), that allows standard webpages to be viewed in a compressed format specifically designed for small, portable hand held wireless devices, a wide range of 2.5G standards have been developed to allow each of the major 2G technologies (GSM, CDMA, and IS-136) to be upgraded incrementally for faster Internet data rates. Figure 2.3 illustrates the various 2.5G and 3G upgrade paths for the major 2G technologies.

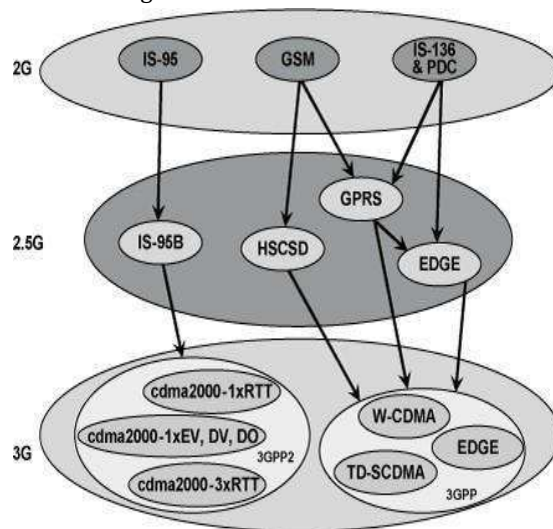


Figure 2.3. Various upgrade paths for 2G technologies.

Table 2.2 describes the required changes to the network infrastructure (e.g., the base station and the switch) and the subscriber terminals (e.g., the handset) for the various upgrade options for 2.5G and 3G. The technical features of each 2.5G upgrade path are described below.

Table 2.2. Current and Emerging 2.5G and 3G Data Communication Standards

Wireless Data Technologies	Channel BW	Duplex	Infrastructure change	Requires New Spectrum	Requires New Handsets
HSCSD	200KHz	FDD	Requires software upgrade at base station.	No	Yes New HSCSD handsets provide 57.6Kbps on HSCSD networks, and 9.6Kbps on GSM networks with dual mode phones. GSM-only phones will not work in HSCSD networks.
GPRS	200KHz	FDD	Requires new packet overlay including routers and gateways.	No	Yes New GPRS handsets work on GPRS networks at 171.2Kbps, 9.6Kbps on GSM networks with dual mode phones. GSM-only phones will not work in GPRS networks.
EDGE	200KHz	FDD	Requires new transceiver at base station. Also, software upgrade to the base station controller and base station.	No	Yes New handsets work on EDGE networks at 384 Kbps, GPRS networks at 144Kbps, and GSM networks at 9.6Kbps with tri mode phones. GSM and GPRS-only phones will not work in EDGE networks.

Wireless Data Technologies	Channel BW	Duplex	Infrastructure change	Requires New Spectrum	Requires New Handsets
W-CDMA	5MHz	FDD	Requires completely new base stations.	Yes	Yes New W-CDMA handsets will work on W-CDMA at 2 Mbps, EDGE networks at 384Kbps, GPRS networks at 144 Kbps, GSM networks at 9.6Kbps. Older handsets will not work in W-CDMA.
IS-95B	1.25 MHz	FDD	Requires new software in base station controller.	No	Yes New handsets will work on IS-95B at 64Kbps and IS-95A at 14.4Kbps. CdmaOne phones can work in IS-95B at 14.4Kbps.
cdma2000 1xRTT	1.25 MHz	FDD	Requires new software in backbone and new channel cards at base station. Also need to build a new packet service node.	No	Yes New handsets will work on 1xRTT at 144Kbps, IS-95B at 64 Kbps, IS 95A at 14.4Kbps. Older handsets can work in 1xRTT but at lower speeds.
cdma2000 1xEV (DO and DV)	1.25 MHz	FDD	Requires software and digital carrier upgrade on 1xRTT networks.	No	Yes New handsets will work on 1xEV at 2.4Mbps, 1xRTT at 144 Kbps IS-95B at 64Kbps, IS-95A at 14.4 Kbps. Older handsets can work in 1xEV but at lower speeds.
cdma2000 3xRTT	3.75 MHz	FDD	Requires backbone modifications and new channel cards at base station.	Maybe	Yes New handsets will work on IS-95A at 14.4 Kbps, IS-95B at 64 Kbps, 1xRTT at 144Kbps, 3xRTT at 2 Mbps. Older handsets can work in 3X but at lower speeds.

Evolution for 2.5G TDMA Standards

Three different upgrade paths have been developed for GSM carriers, and two of these solutions also support IS-136. The three TDMA upgrade options include: (a) High Speed Circuit Switched Data (HSCSD); (b) General Packet Radio Service (GPRS); and (c) Enhanced Data Rates for GSM Evolution (EDGE). These options provide significant improvements in Internet access speed over today's GSM and IS-136 technology and support the creation of new Internet-ready cell phones.

HSCSD for 2.5G GSM

As the name implies, High Speed Circuit Switched Data is a circuit switched technique that allows a single mobile subscriber to use consecutive user time slots in the GSM standard. That is, instead of limiting each user to only one specific time slot in the GSM TDMA standard, HSCSD allows individual data users to commandeer consecutive time slots in order to obtain higher speed data access on the GSM network. HSCSD relaxes the error control coding algorithms originally specified in the GSM standard for data transmissions and increases the available application data rate to 14,400 bps, as compared to the original 9,600 bps in the GSM specification. HSCSD is ideal for dedicated streaming Internet access or real-time interactive web sessions and simply requires the service provider to implement a software change at existing GSM base stations.

GPRS for 2.5G GSM and IS-136

General Packet Radio Service is a packet-based data network, which is well suited for non-real time Internet usage, including the retrieval of email, faxes, and asymmetric web browsing, where the user downloads much more data than it uploads on the Internet. Unlike HSCSD, which dedicates circuit switched channels to specific users, GPRS supports multi-user network sharing of individual radio channels and time slots. Similar to the Cellular Digital Packet Data (CDPD) standard developed for the North American AMPS systems in the early 1990s, the GPRS standard provides a packet network on dedicated GSM or IS-136 radio channels. GPRS retains the original modulation formats specified in the original 2G TDMA standards, but uses a completely redefined air interface in order to better handle packet data access.

As is the case for any packet network, the data throughput experienced by an individual GPRS user decreases substantially as more users attempt to use the network or as propagation conditions become poor for particular users. It is worth noting that GPRS was originally designed to provide a packet data access overlay solely for GSM networks, but at the request of North American IS-136 operators (see UWC-136 Air Interface in Table 2.3), GPRS was extended to include both TDMA standards.

Table 2.3. Leading IMT-2000 Candidate Standards as of 1998 (adapted from [Lib99])

Air Interface	Mode of Operation	Duplexing Method	Key Features
cdma2000U STIA TR45.5	Multi-Carrier and Direct Spreading DS-SS CDMA with $N=1, 3, 6, 9, 12$	FDD and TDD Modes	Backward compatibility with IS-95A and IS-95B. Downlink can be implemented using either Multi-Carrier or Direct Spreading. Uplink can support simultaneous combination of Multi-Carrier or Direct Spreading Auxiliary carrier to help with downlink channel estimation in forward link beamforming.
UTRA (UMTS Terrestrial Radio Access) ETSS MG2	DS-SS CDMA at Rates of $N \times 0.960$ Mcps with $N=4, 8, 16$	FDD and TDD Modes	Wideband DS-SS CDMA System. Backward compatibility with GSM/DCS-1900. Upto 2.048 Mbps on Downlink in FDD Mode. Minimum forward channel bandwidth of 5 GHz. The collection of proposed standards represented here each exhibit unique features but support a common set of chip rates, 10 ms frame structure, with 16 slots per frame. Connection-dedicated pilots assist in downlink beamforming.
W-CDMA/NA (Wideband CDMA North America) USAT1P1-ATIS			
W-CDMA/Japan (Wideband CDMA Japan) ARIB			
CDMA II South Korea TTA			
WIMS/W-CDMA USATIA TR46.1			
CDMA I South Korea TTA	DS-SS CDMA at $N \times 0.9216$ Mcps with $N=1, 4, 16$	FDD and TDD Modes	Upto 512 kbps per spreading code, code aggregation upto 2.048 Mbps.
UWC-136 (Universal Wireless Communications Consortium) USA TIA TR45.3	TDMA - Up to 722.2 kbps (Outdoor/Vehicular), Upto 5.2 Mbps (Indoor Office)	FDD (Outdoor/Vehicular), TDD (Indoor Office)	Backward compatibility and upgrade path for both IS-136 and GSM. Fits into existing IS-136 and GSM. Explicit plans to support adaptive antenna technology.
TD-SSCDMA China Academy of Telecommunication Technology (CATT)	DS-SS CDMA 1.1136 Mcps	TDD	RF channel bitrate upto 2.227 Mbps. Use of smart antenna technology is fundamental (but not strictly required) in TDSSCDMA.
DECT ETSI Project (EPDECT)	1150-3456 kbps TDMA	TDD	Enhanced version of 2G DECT technology.

EDGE for 2.5G GSM and IS-136

Enhanced Data rates for GSM (or Global) Evolution is a more advanced upgrade to the GSM standard, and

requires the addition of new hardware and software at existing base stations. EDGE introduces a new digital

modulation format, 8-PSK (octal phase shift keying), which is used in addition to GSM's standard GMSK modulation. EDGE allows for nine different (autonomously and rapidly selectable) air interface formats, known as *multiple modulation and coding schemes* (MCS), with varying degrees of error control protection. Each MCS state may use either GMSK (low data rate) or 8-PSK (high data rate) modulation for network access, depending on the instantaneous demands of the network and the operating conditions. Because of the higher data rates and relaxed error control covering in many of the selectable air interface formats, the coverage range is smaller in EDGE than in HSDRC or GPRS. EDGE is sometimes referred to as Enhanced GPRS, or EGPRS.

IS-95B for 2.5G CDMA

Unlike the several GSM and IS-136 evolutionary paths to high speed data access, CDMA (often called *cdmaOne*) has a single upgrade path for eventual 3G operation. The interim data solution for CDMA is called IS-95B. Like GPRS, IS-95B is already being deployed worldwide, and provides high speed packet and circuit switched data access on a common CDMA radio channel by dedicating multiple orthogonal user channels (Walsh functions) for specific users and specific purposes. Each IS-95 CDMA radio channel supports up to 64 different user channels.

The 2.5G CDMA solution, IS-95B, supports *medium data rate* (MDR) service by allowing a dedicated user to command up to eight different user Walsh codes simultaneously and in parallel for an instantaneous throughput of 115.2 kbps per user (8×14.4 kbps). IS-95B also specifies hard handoff procedures that allow subscriber units to search different radio channels in the network without instruction from the switch so that subscriber units can rapidly tune to different base stations to maintain link quality.

Third Generation (3G) Wireless Networks

3G systems promise unparalleled wireless access in ways that have never been possible before. Multi-megabit Internet access, communications using Voice over Internet Protocol (VoIP), voice-activated calls, unparalleled network capacity, and ubiquitous "always-on" access are just some of the advantages being touted by 3G developers. Companies developing 3G equipment envision users having the ability to receive live music, conduct interactive web sessions, and have simultaneous voice and data access with multiple parties at the same time using a single mobile handset, whether driving, walking, or standing still in an office setting.

The eventual 3G evolution for CDMA systems leads to cdma2000. Several variants of CDMA 2000 are currently being developed, but they all are based on the fundamentals of IS-95 and IS-95B technologies. The eventual 3G evolution for GSM, IS-136, and PDC systems leads to Wideband CDMA (W-CDMA), also called Universal Mobile Telecommunications Service (UMTS). W-CDMA is based on the network fundamentals of GSM, as well as the merged versions of GSM and IS-136 through EDGE.

3G W-CDMA (UMTS)

The Universal Mobile Telecommunications System (UMTS) is a visionary air interface standard that has evolved since late 1996 under the auspices of the European Telecommunications Standards Institute (ETSI). European carriers, manufacturers, and government regulators collectively developed the early versions of UMTS as a competitive open air-interface standard for third generation wireless telecommunications.

UMTS, or W-CDMA, assures backward compatibility with the second generation GSM, IS-136, and PDC TDMA technologies, as well as all 2.5G TDMA technologies. The network structure and bit level packaging of GSM data is retained by W-CDMA, with additional capacity and bandwidth provided by a new CDMA air interface. The 3G W-CDMA air interface standard had been designed for "always-on" packet-based wireless service, so that computers, entertainment devices, and telephones may all share the same wireless network and be connected to the Internet, anytime, anywhere. W-CDMA requires a minimum spectrum allocation of 5 MHz, which is an important distinction from the other G standards. With W-CDMA data rates from as low as 8 kbps to as high as 2 Mbps will be carried simultaneously on a single W-CDMA 5 MHz radio channel, and each channel will be able to support between 100 and 350 simultaneous voice calls at once, depending on antenna sectoring, propagation conditions, user velocity, and antenna polarizations. W-CDMA employs variable/selectable direct sequence spread spectrum chip rates that can exceed 16 Megachips per second per user.

3G cdma2000

The cdma2000 vision provides a seamless and evolutionary high data rate upgrade path for current users of 2G and 2.5G CDMA technology, using a building block approach that centers on the original 2G CDMA channel bandwidth of 1.25 MHz per radio channel.

The first 3G CDMA air interface, cdma2000 1xRTT, implies that a single 1.25 MHz radio channel is used (the initials RTT stand for *Radio Transmission Technology*). The ultimate 3G solution for CDMA relies upon multicarrier techniques that gang adjacent cdmaOne radio channels together for increased bandwidth. The cdma2000 3xRTT standard uses three adjacent 1.25 MHz radio channels that are used together to provide packet data throughput speeds in excess of 2 Mbps per user, depending upon cell loading, vehicle speed, and propagation conditions.

3GTD-SCDMA

The China Academy of Telecommunications Technology (CATT) and Siemens Corporation jointly submitted an IMT-2000 3G standard proposal in 1998, based on Time Division-Synchronous Code Division Multiple Access (TD-SCDMA). This proposal was adopted by ITU as one of the 3G options in late 1999.

TD-SCDMA relies on the existing core GSM infrastructure and allows a 3G network to evolve through the addition of high data rate equipment at each GSM base station. TD-SCDMA combines TDMA and TDD techniques to provide a data-only overlay in an existing GSM network. By using TDD, different time slots within a single frame on a single carrier frequency are used to provide both forward channel and reverse channel transmissions. For the case of asynchronous traffic demand, such as when a user downloads a file, the forward link will require more bandwidth than the reverse link, and thus more time slots will be dedicated to providing forward link traffic than for providing reverse link traffic. TD-SCDMA proponents claim that the TDD feature allows this 3G standard to be very easily and inexpensively added to existing GSM systems.

Wireless Local Loop (WLL) and LMDS

Fixed wireless equipment is extremely well suited for rapidly deploying a broadband connection in many instances, and this approach is steadily becoming more popular for providing "last mile" broadband local loop access, as well as for emergency or redundant point-to-point or point-to-multipoint private networks.

Modern fixed wireless systems are usually assigned microwave or millimeter radio frequencies in the 28 GHz band and higher, which is greater than ten times the carrier frequency of 3G terrestrial cellular telephone networks. At these higher frequencies, the wavelengths are extremely small, which in turn allows very high gain directional antennas to be fabricated in small physical form factors. Microwave wireless links can be used to create a wireless local loop (WLL) such as the one shown in [Figure 2.4](#).

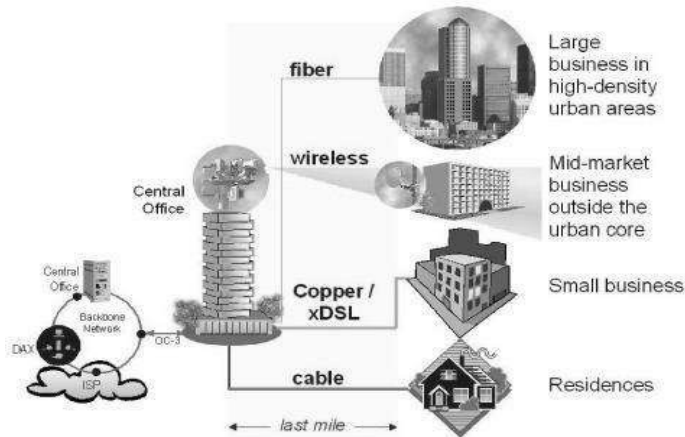


Figure 2.4. Example of the emerging applications and markets for broadband services.

The local loop can be thought of as the "last mile" of the telecommunication network that resides between the central office (CO) and the individual homes and businesses in close proximity to the CO.

Governments throughout the world have realized that WLL could greatly improve the efficiency of their citizens while stimulating competition that could lead to improved telecommunications services. A vast array of new services and applications have been proposed and are in the early stages of commercialization. These services include the concept of Local Multipoint Distribution Service (LMDS), which provides broadband telecommunications access in the local exchange.

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In 1998, 1300 MHz of unused spectrum in the 27–31 GHz band was auctioned by the US government to support LMDS. Similar auctions have been held in other countries around the world.

One of the most promising applications for LMDS is in a local exchange carrier (LEC) network. [Figure 2.7](#) shows a typical network configuration, where the LEC owns a very wide bandwidth asynchronous transfer mode (ATM) or Synchronous Optical Network (SONET) backbone switch, capable of connecting hundreds of megabits per second of traffic with the Internet, the PSTN, or to its own private network. As long as a LOS

path exists, LMDS will allow LECs to install wireless equipment on the premises of customers for rapid broadband connectivity without having to lease or install its own cables to the customers.

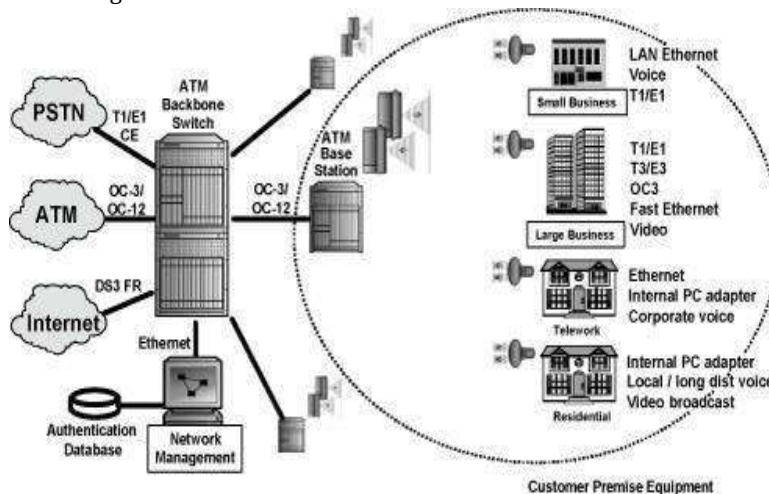


Figure 2.7. A wireless Competitive Local Exchange Carrier (CLEC) using Asynchronous Transfer Mode (ATM) distribution.

Wireless Local Area Networks (WLANs)

In 1997 the FCC allocated 300 MHz of unlicensed spectrum in the Industrial Scientific and Medical (ISM) bands of 5.150–5.350 GHz and 5.725–5.825 GHz for the express purpose of supporting low-power license-free spread spectrum data communication. This allocation is called the *Unlicensed National Information Infrastructure (UNII)* band.

The IEEE 802.11 Wireless LAN working group was founded in 1987 to begin standardization of spread spectrum WLANs for use in the ISM bands. Figure 2.10 illustrates the evolution of IEEE 802.11 Wireless LAN standards, which also include infrared communications. Figure 2.10 shows how both frequency hopping and direct sequence approaches were used in the original IEEE 802.11 standard (2 Mbps user throughput), but as of late 2001 only direct sequence spread spectrum (DS-SS) modems had thus far been standardized for high rate (11 Mbps) user data rates within IEEE 802.11.

The DS-SS IEEE 802.11b standard has been named *Wi-Fi* by the *Wireless Ethernet Compatibility Alliance (WECA)*, a group that promotes adoption of 802.11b DS-SS WLAN equipment and interoperability between vendors. IEEE 802.11g is developing *Complimentary Code Keying Orthogonal Frequency Division Multiplexing (CCK-OFDM)* standards in both the 2.4 GHz (802.11b) and 5 GHz (802.11a) bands, and will support roaming capabilities and dual-band use for public WLAN networks, while supporting backward compatibility with 802.11b technology.

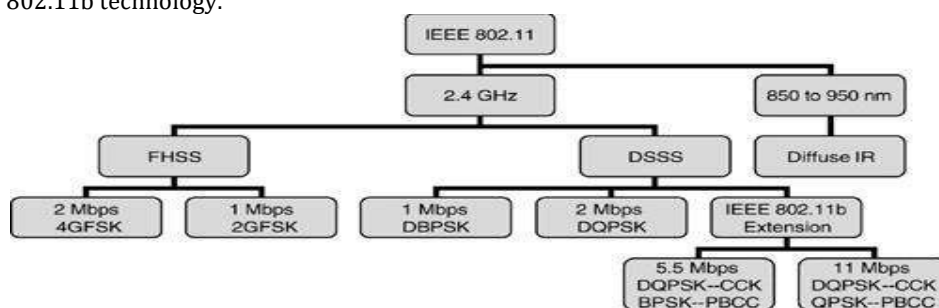


Figure 2.10. Overview of the IEEE 802.11 Wireless LAN standard.

The frequency-hopping spread spectrum (FH-SS) proponents of IEEE 802.11 have formed the HomeRF standard that supports frequency hopping equipment. It is worth noting that both DS and FH types of WLANs must operate in the same unlicensed bands that contain cordless phones, baby monitors, Bluetooth devices, and other WLAN users.

Table 2.4. IEEE 802.11b Channels for Both DS-SS and FH-SS WLAN Standards

Country	Frequency Range Available	DSSS Channels Available	FHSS Channels Available
United States	2.4 to 2.4835 GHz	1 through 11	2 through 80
Canada	2.4 to 2.4835 GHz	1 through 11	2 through 80
Japan	2.4 to 2.497 GHz	1 through 14	2 through 95
France	2.4465 to 2.4835 GHz	10 through 13	48 through 82
Spain	2.445 to 2.4835 GHz	10 through 11	47 through 73

Country	Frequency Range Available	DSSS Channels Available	FHSS Channels Available
Remainder of Europe	2.4 to 2.4835	1 through 13	2 through 80

Figure 2.12 illustrates the unique WLAN channels that are specified in the IEEE 802.11b standard for the 2400–2483.5 MHz band. All WLANs are manufactured to operate on any one of the specified channels and are assigned to a particular channel by the network operator when the WLAN system is first installed. The channelization scheme used by the network installer becomes very important for a high density WLAN installation, since neighboring access points must be separated from one another in frequency to avoid interference and significantly degraded performance.

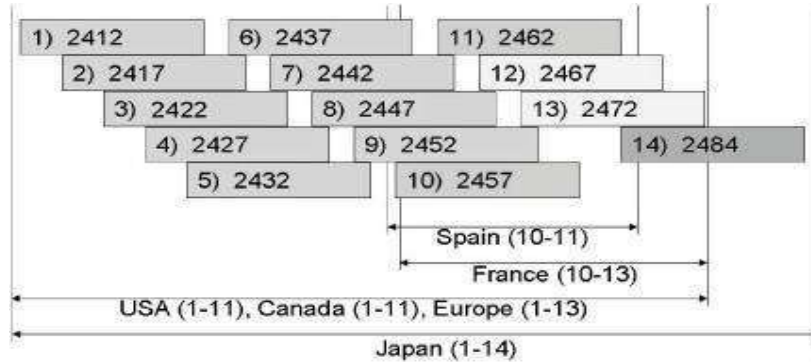


Figure 2.12. Channelization scheme for IEEE 802.11b throughout the world.

In Europe in the mid 1990s, the *High Performance Radio Local Area Network* (HIPERLAN) standard was developed to provide a similar capability to IEEE 802.11. HIPERLAN was intended to provide individual wireless LANs for computer communications and used the 5.2 GHz and the 17.1 GHz frequency bands. HIPERLAN provides asynchronous user data rates of between 1 to 20 Mbps, as well as time bounded messaging at rates of 64 kbps to 2.048 Mbps. HIPERLAN was designed to operate up to vehicle speeds of 35 km/hr, and typically provided 20 Mbps throughput at 50 m range.

In 1997, Europe's ETSI established a standardization committee for Broadband Radio Access Networks (BRANs). The goal of BRAN is to develop a family of broadband WLAN-type protocols that allow user interoperability, covering both short range (e.g., WLAN) and long range (e.g., fixed wireless) networking. HIPERLAN/2 has emerged as the next generation European WLAN standard and will provide up to 54 Mbps of user data to a variety of networks, including the ATM backbone, IP based networks, and the UMTS core.

Bluetooth and Personal Area Networks (PANs)

Bluetooth is an open standard that has been embraced by over 1,000 manufacturers of electronic appliances. It provides an ad-hoc approach for enabling various devices to communicate with one another within a nominal 10 meter range. Named after King Harald Bluetooth, the 10th century Viking who united Denmark and Norway, the Bluetooth standard aims to unify the connectivity chores of appliances within the personal workspace of an individual.

Bluetooth operates in the 2.4 GHz ISM Band (2400–2483.5 MHz) and uses a frequency hopping TDD scheme for each radio channel. Each Bluetooth radio channel has a 1 MHz bandwidth and hops at a rate of approximately 1600 hops per second. Transmissions are performed in 625 microsecond slots with a single packet transmitted over a single slot. For long data transmissions, particular users may occupy multiple slots using the same transmission frequency, thus slowing the instantaneous hopping rate to below 1600 hops/second. The frequency hopping scheme of each Bluetooth user is determined from a cyclic code of length $2^{27} - 1$, and each user has a channel symbol rate of 1 Mbps using GFSK modulation. The standard has been designed to support operation in very high interference levels and relies on a number of forward error control (FEC) coding and automatic repeat request (ARQ) schemes to support a raw channel bit error rate (BER) of about 10^{-3} .

Different countries have allocated various channels for Bluetooth operation. In the US and most of Europe, the FHSS 2.4 GHz ISM band is available for Bluetooth use. A detailed list of states is defined in the Bluetooth standard to support a wide range of applications, appliances, and potential uses of the Personal Area Network. Audio, text, data, and even video is contemplated in the Bluetooth standard [Tra01]. Figure 2.17 provides a depiction of the Bluetooth concept where a gateway to the Internet via IEEE 802.11b is shown as a conceptual possibility.

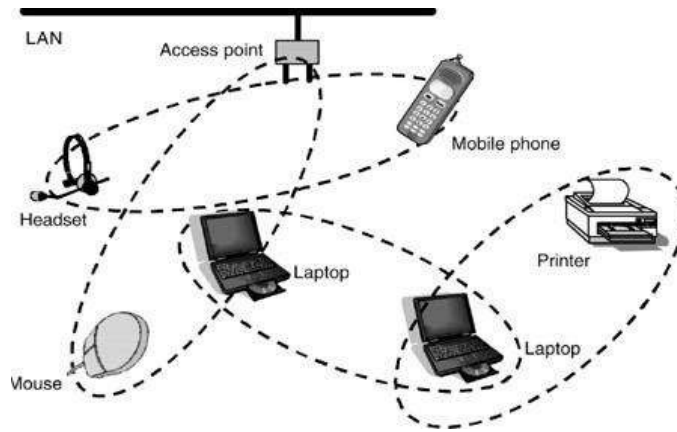


Figure 2.17. Example of a Personal Area Network (PAN) as provided by the Bluetooth standard.

The IEEE 802.15 standards committee has been formed to provide an international forum for developing Bluetooth and other PANs that interconnect pocket PCs, personal digital assistants (PDAs), cellphones, light projectors, and other appliances [Bra00]. With the rapid proliferation of *wearable computers*, such as PDAs, cellphones, smart cards, and position location devices, PANs may provide the connection to an entire new era of remote retrieval and monitoring of the world around us.

UNIT II

Mobile Radio Propagation (Large-Scale Path Loss)

Introduction

There are two basic ways of transmitting an electro-magnetic (EM) signal, through a guided medium or through an unguided medium. Guided mediums such as coaxial cables and fiber optic cables, are far less hostile toward the information carrying EM signal than the wireless or the unguided medium. It presents challenges and conditions which are unique for this kind of transmissions. A signal, as it travels through the wireless channel, undergoes many kinds of propagation effects such as reflection, diffraction and scattering, due to the presence of buildings, mountains and other such obstructions. Reflection occurs when the EM waves impinge on objects which are much greater than the wavelength of the traveling wave. Diffraction is a phenomena occurring when the wave interacts with a surface having sharp irregularities. Scattering occurs when the medium through the wave is traveling contains objects which are much smaller than the wavelength of the EM wave. These varied phenomena's lead to large scale and small scale propagation losses. Due to the inherent randomness associated with such channels they are best described with the help of statistical models. Models which predict the mean signal strength for arbitrary transmitter receiver distances are termed as large scale propagation models. These are termed so because they predict the average signal strength for large Tx-Rx separations, typically for hundreds of kilometers.

Free Space Propagation Model:

The free space propagation model is used to predict received signal strength when the transmitter and receiver have a clear, unobstructed line-of-sight path between them. Satellite communication systems and microwave line-of-sight radio links typically undergo free space propagation. As with most large-scale radio wave propagation models, the free space model predicts that received power decays as a function of the T-R separation distance raised to some power (i.e. a power law function). The free space power received by a receiver antenna which is separated from a radiating transmitter antenna by a distance d , is given by the Friis free space equation

$$P_r(d) = P_t G_t G_r \lambda^2 / (4\pi d)^2$$

Where P_t is the transmitted power, $P_r(d)$ is the received power which is a function of the T-R separation, G_t is the transmitter antenna gain, G_r is the receiver antenna gain, d is the T-R separation distance in meters and λ is the wavelength in meters. The gain of an antenna is related to its effective aperture, A_e by,

$$G = 4\pi A_e / \lambda^2$$

The effective aperture A_e is related to the physical size of the antenna, and λ is related to the carrier frequency by,

$$\lambda = c/f = 2\pi c/\omega_c$$

where f is the carrier frequency in Hertz, ω_c is the carrier frequency in radians per second and c is the speed of light given in meters/s. An isotropic radiator is an ideal antenna which radiates power with unit gain uniformly in all directions, and is often used to reference antenna gains in wireless systems. The effective isotropic radiated power (EIRP) is defined as **EIRP = $P_t G_t$**

it represents the maximum radiated power available from a transmitter in the direction of maximum antenna gain, as compared to an isotropic radiator. In practice, effective radiated power (ERP) is used instead of EIRP to denote the maximum radiated power as compared to a half-wave dipole antenna (instead of an isotropic antenna).

The path loss, which represents signal attenuation as a positive quantity measured in dB, is defined as the difference (in dB) between the effective transmitted power and the received power, and may or may not include the effect of the antenna gains. The path loss for the free space model when antenna gains are included is given by

$$PL(\text{dB}) = 10 \log(P_t/P_r) = -10 \log [G_t G_r \lambda^2 / (4\pi d)^2]$$

When antenna gains are excluded, the antennas are assumed to have unity gain, and path loss is given by **PL (dB) = $10 \log (P_t/P_r) = -10 \log [\lambda^2 / (4\pi d)^2]$**

The Friis free space model is only a valid predictor for P_r for values of d which are in the far-field of the transmitting antenna. The far-field or Fraunhofer region of a transmitting antenna is defined as the region beyond the far-field distance d_f , which is related to the largest linear dimension of the transmitter antenna aperture and the carrier wavelength. The Fraunhofer distance is given by

$$d_f = 2D^2/\lambda$$

Where D is the largest physical linear dimension of the antenna. The far-field region d_f must satisfy

$$d_f \gg D$$

What will be the far-field distance for a Base station antenna with Largest dimension $D=0.5\text{m}$ Frequency of operation $f_c=900\text{MHz}$

Sol:

$$\lambda = c/f = 3 \times 10^8 / 900 \times 10^6 = 0.33\text{m}$$

$$d_f = 2D^2/\lambda = 2(0.5)^2/0.33 = 1.5\text{m}$$

If a transmitter produces 50 watts of power, express the transmit power in units of (a) dBm, and (b) dBW. If 50 watts is applied to a unity gain antenna with a 900 MHz carrier frequency, find the received power in dBm at a free space distance of 100 m from the antenna. What is P_r (10 km)? Assume unity gain for the receiver antenna.

Given:

Transmitter power, $P_t = 50 \text{ W}$.

Carrier frequency, $f_c = 900 \text{ MHz}$

Using equation (3.9),

(a) Transmitter power,

$$\begin{aligned} P_t(\text{dBm}) &= 10 \log [P_t(\text{mW}) / (1 \text{ mW})] \\ &= 10 \log [50 \times 10^3] = 47.0 \text{ dBm}. \end{aligned}$$

(b) Transmitter power,

$$\begin{aligned} P_t(\text{dBW}) &= 10 \log [P_t(\text{W}) / (1 \text{ W})] \\ &= 10 \log [50] = 17.0 \text{ dBW}. \end{aligned}$$

The received power can be determined using equation (3.1).

$$P_r = \frac{P_t G_t G_r \lambda^2}{(4\pi)^2 d^2 L} = \frac{50 (1) (1) (1/3)^2}{(4\pi)^2 (100)^2 (1)} = 3.5 \times 10^{-6} \text{ W} = 3.5 \times 10^{-3} \text{ mW}$$

$$P_r(\text{dBm}) = 10 \log P_r(\text{mW}) = 10 \log (3.5 \times 10^{-3} \text{ mW}) = -24.5 \text{ dBm}.$$

The received power at 10 km can be expressed in terms of dBm using equation (3.9), where $d_0 = 100 \text{ m}$ and $d = 10 \text{ km}$

$$\begin{aligned} P_r(10 \text{ km}) &= P_r(100) + 20 \log \left[\frac{100}{10000} \right] = -24.5 \text{ dBm} - 40 \text{ dB} \\ &= -64.5 \text{ dBm}. \end{aligned}$$

Three Basic Propagation Mechanisms

Reflection, Diffraction and Scattering is the three basic propagation mechanisms which impact propagation in a mobile communication system.

Reflection occurs when a propagating electromagnetic wave impinges upon an object which has very large dimensions when compared to the wavelength of the propagating wave. Reflections occur from the surface of the earth and from buildings and walls.

Diffraction occurs when the radio path between the transmitter and receiver is obstructed by a surface that has sharp irregularities (edges). The secondary waves resulting from the obstructing surface are present throughout the space and even behind the obstacle, giving rise to a bending of waves around the obstacle, even when a line-of-sight path does not exist between transmitter and receiver. At high frequencies, diffraction, like reflection depends on the geometry of the object, as well as the amplitude, phase, and polarization of the incident wave at the point of diffraction.

Scattering occurs when the medium through which the wave travels consists of objects with dimensions that are small compared to the wavelength, and where the number of obstacles per unit volume is large. Scattered waves are produced by rough surfaces, small objects, or by other irregularities in the channel. In practice, foliage, street signs, and lamp posts induce scattering in a mobile communications system.

Reflection:

When a radio wave propagating in one medium impinges upon another medium having different electrical properties, the wave is partially reflected and partially transmitted. If the plane wave is incident on a perfect dielectric, part of the energy is transmitted into the second medium and part of the energy is reflected back into the first medium, and there is no loss of energy in absorption. If the second medium is a perfect conductor, then all incident energy is reflected back into the first medium without loss of energy. The electric field intensity of the reflected and transmitted waves may be related to the incident wave in the medium of origin through the Fresnel reflection coefficient (Γ). The reflection coefficient is a function of the material properties, and generally depends on the wave polarization, angle of incidence, and the frequency of the propagating wave.

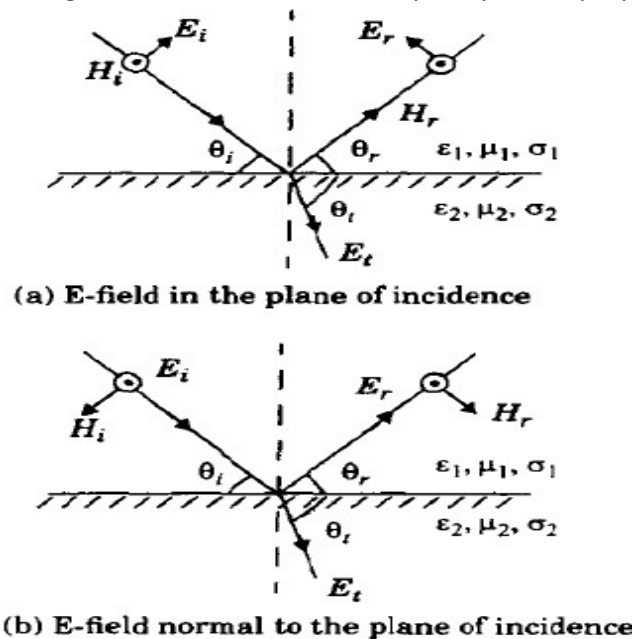


Figure 3.4
Geometry for calculating the reflection coefficients between two dielectrics.

Reflection from Dielectrics:

Figure 3.4 shows an electromagnetic wave incident at an angle θ_i with the plane of the boundary between two dielectric media. As shown in the figure, part of the energy is reflected back to the first media at an angle θ_r , and part of the energy is transmitted (refracted) into the second media at an angle θ_t . The nature of reflection varies with the direction of polarization of the E-field. The behavior for arbitrary directions of polarization can be studied by considering the two distinct cases shown in Figure

The plane of incidence is defined as the plane containing the incident, reflected, and transmitted rays. In Figure 3.4a, the E-field polarization is parallel with the plane of incidence (that is, the E-field has a vertical polarization, or normal component, with respect to the reflecting surface) and in Figure 3.4b, the E-field polarization is perpendicular to the plane of incidence (that is, the incident E-field is pointing out of the page towards the reader, and is perpendicular to the page and parallel

to the reflecting surface).

Because of superposition, only two orthogonal polarizations need be considered to solve general reflection problems. The reflection coefficients for the two cases of parallel and perpendicular E-field polarization at the boundary of two dielectrics are given by

$$\Gamma_{\parallel} = \frac{E_r}{E_i} = \frac{\eta_2 \sin \theta_t - \eta_1 \sin \theta_i}{\eta_2 \sin \theta_t + \eta_1 \sin \theta_i} \quad (\text{E-field in plane of incidence})$$

$$\Gamma_{\perp} = \frac{E_r}{E_i} = \frac{\eta_2 \cos \theta_t - \eta_1 \cos \theta_i}{\eta_2 \cos \theta_t + \eta_1 \cos \theta_i} \quad (\text{E-field not in plane of incidence})$$

Where η is the intrinsic impedance of the respective medium.

Where ϵ is the permittivity of the respective medium.

Brewster Angle:

$$\Gamma_{\parallel} = \frac{-\epsilon_r \sin \theta_i + \sqrt{\epsilon_r - \cos^2 \theta_i}}{\epsilon_r \sin \theta_i + \sqrt{\epsilon_r - \cos^2 \theta_i}}$$

$$\Gamma_{\perp} = \frac{\sin \theta_i - \sqrt{\epsilon_r - \cos^2 \theta_i}}{\sin \theta_i + \sqrt{\epsilon_r - \cos^2 \theta_i}}$$

The Brewster angle is the angle at which no reflection occurs in the medium of origin. It occurs when the incident angle θ_i is such that the reflection coefficient Γ_{\parallel} is equal to zero. The Brewster angle is given by the value of θ_B which satisfies

$$\sin(\theta_B) = \sqrt{\epsilon_1 / (\epsilon_1 + \epsilon_2)}$$

For the case when the first medium is free space and the second medium has a relative permittivity ϵ_r , above equation can be expressed as

$\sin(\theta_B) = \sqrt{(\epsilon_r - 1) / (\epsilon_r^2 - 1)}$ Note that the Brewster angle occurs only for vertical (i.e. parallel) polarization.

Reflection from Perfect Conductors:

Since electromagnetic energy cannot pass through a perfect conductor a plane wave incident on a conductor has all of its energy reflected. As the electric field at the surface of the conductor must be equal to zero at all times in order to obey Maxwell's equations, the reflected wave must be equal in magnitude to the incident wave. For the case when E-field polarization is in the plane of incidence, the boundary conditions require that $\theta_i = \theta_r$ and $E_i = -E_r$ (E-field in plane of incidence)

Similarly, for the case when the E-field is horizontally polarized, the boundary conditions require that $\theta_i = \theta_r$ and $E_i = -E_r$ (E-field not in plane of incidence)

Ground Reflection (2-ray) Model:

In a mobile radio channel, a single direct path between the base station and a mobile is seldom the only physical means for propagation, and hence the free space propagation model is in most cases inaccurate when used alone. The 2-ray ground reflection model shown in Figure 3.7 is a useful propagation model that is based on geometric optics, and considers both the direct path and a ground reflected propagation path between transmitter and receiver. This model has been found to be reasonably accurate for predicting the large-scale signal strength over distances of several kilometers for mobile radio systems that use tall towers (heights which exceed 50 m), as well as for line-of-sight, microcell channels in urban environments.

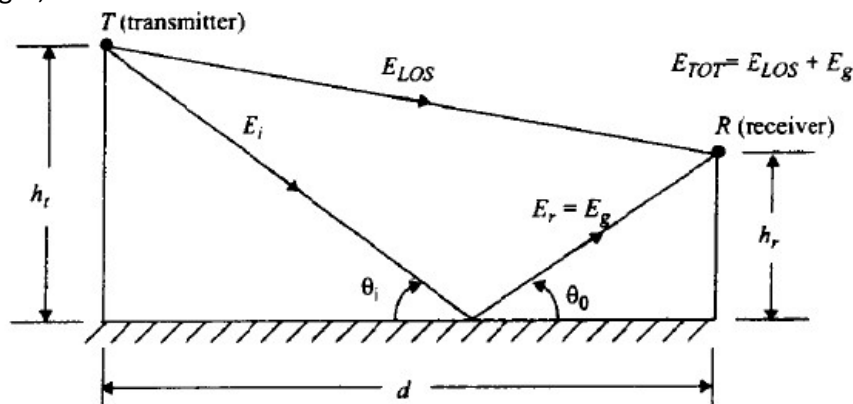


Figure 3.7 Two-ray ground reflection model.

Referring to Figure 3.7, h_t is the height of the transmitter and h_r is the height of the receiver. If E_0 is the free space E-field (in units of V/m) at a reference distance d_0 from the transmitter, then for $d > d_0$, the free space propagating E-field is given by

$$E(d, t) = \frac{E_0 d_0}{d} \cos\left(\omega_c \left(t - \frac{d}{c}\right)\right) \quad (d > d_0)$$

Two propagating waves arrive at the receiver: the direct wave that travels a distance d' ; and the reflected wave that travels a distance d'' .

The electric field $E_{TOT}(d, t)$ can be expressed as the sum of equations for distances d' and d'' (i.e. direct wave and reflected wave).

$$E_{TOT}(d, t) = \frac{E_0 d_0}{d'} \cos\left(\omega_c \left(t - \frac{d'}{c}\right)\right) + (-1) \frac{E_0 d_0}{d''} \cos\left(\omega_c \left(t - \frac{d''}{c}\right)\right)$$

Diffraction:

Diffraction allows radio signals to propagate around the curved surface of the earth, beyond the horizon, and to propagate behind obstructions. Although the received field strength decreases rapidly as a receiver moves deeper into the obstructed (shadowed) region, the diffraction field still exists and often has sufficient strength to produce a useful signal.

The phenomenon of diffraction can be explained by Huygens' principle, which states that all points on a wave front can be considered as point sources for the production of secondary wavelets, and that these wavelets combine to produce a new wave front in the direction of propagation. Diffraction is caused by the propagation of secondary wavelets into a shadowed region. The field strength of a diffracted wave in the shadowed region is the vector sum of the electric field components of all the secondary wavelets in the space around the obstacle.

Fresnel Zone Geometry:

Fresnel zones represent successive regions where secondary waves have a path length from the TX to the RX which are $\lambda/2$ greater in path length than of the LOS path. The plane below illustrates successive Fresnel zones.

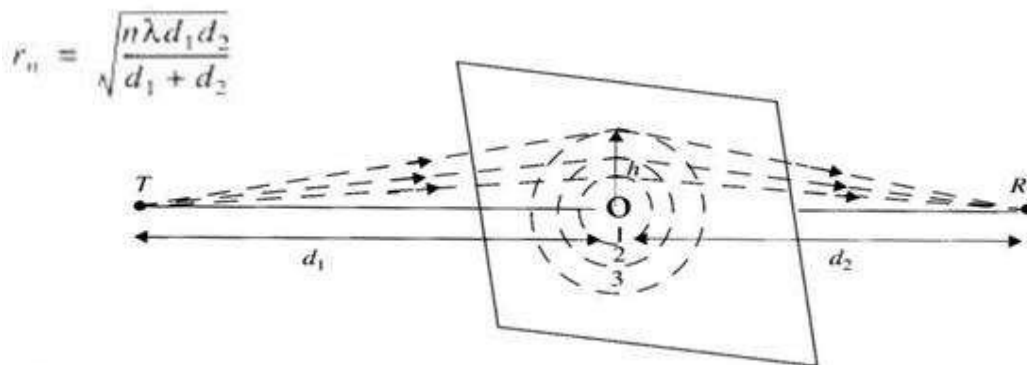


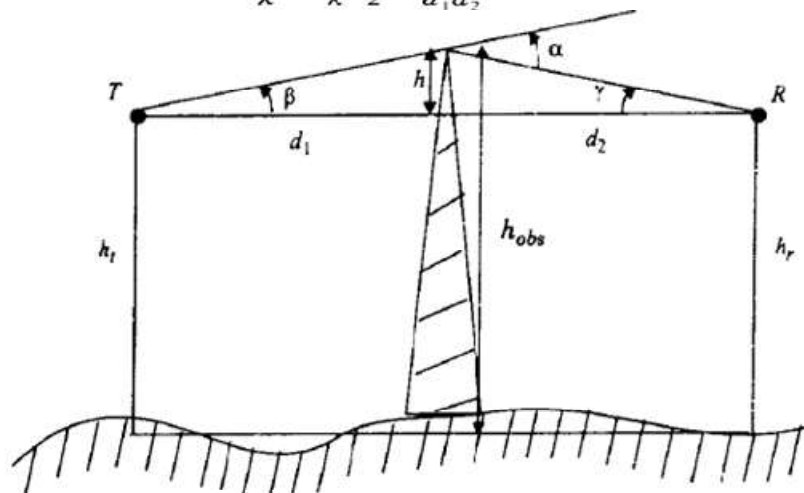
Figure 4.11 Concentric circles which define the boundaries of successive Fresnel zones.

Consider a transmitter and receiver separated in free space as shown in Figure 3.10a. Let an obstructing screen of effective height h within infinite width (going into and out of the paper,) be

$$\Delta = \frac{h^2(d_1 + d_2)}{2d_1d_2}$$

The corresponding phase difference is given by

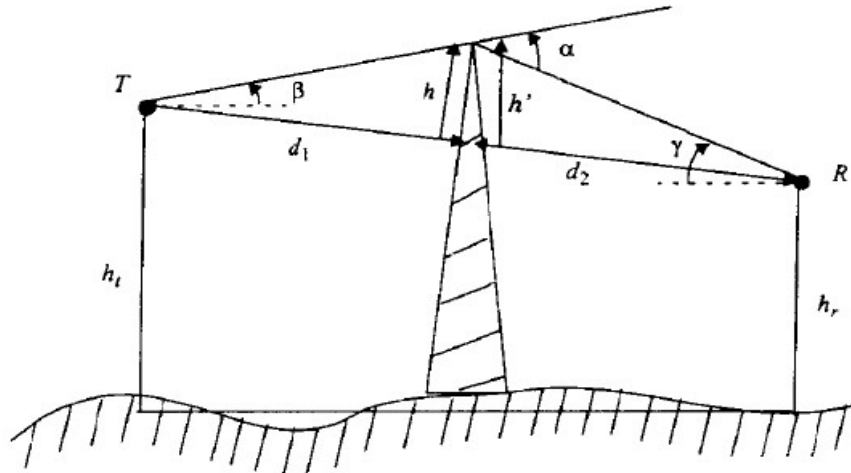
$$\phi = \frac{2\pi\Delta}{\lambda} = \frac{2\pi}{\lambda} \frac{h^2}{2} \frac{(d_1 + d_2)}{d_1d_2}$$



(a) Knife-edge diffraction geometry. The point T denotes the transmitter and R denotes the receiver, with an infinite knife-edge obstruction blocking the line-of-sight path.

placed between them at a distance d_1 from the transmitter and d_2 from the receiver. It is apparent that the wave propagating from the transmitter to the receiver via the top of the screen travels a longer distance than if a direct line-of-sight path (through the screen) existed. Assuming $h \ll d_1, d_2$

and $h \gg \lambda$, then the difference between the direct path and the diffracted path, called the excess path length (Δ), can be obtained from the geometry of Figure as



(b) Knife-edge diffraction geometry when the transmitter and receiver are not at the same height. Note that if α and β are small and $h \ll d_1$ and d_2 , then h and h' are virtually identical and the geometry may be redrawn as shown in Figure 3.10c.

Knife-edge Diffraction Model:

Estimating the signal attenuation caused by diffraction of radio waves over hills and buildings is essential in predicting the field strength in a given service area. Generally, it is impossible to make very precise estimates of the diffraction losses, and in practice prediction is a process of theoretical approximation modified by necessary empirical corrections. Though the calculation of diffraction losses over complex and irregular terrain is a mathematically difficult problem, expressions for diffraction losses for many simple cases have been derived. As a starting point, the limiting case of propagation over a knife-edge gives good insight into the order of magnitude of diffraction loss.

When shadowing is caused by a single object such as a hill or mountain, the attenuation caused by diffraction can be estimated by treating the obstruction as a diffracting knife edge. This is the simplest of diffraction models, and the diffraction loss in this case can be readily estimated using the classical Fresnel solution for the field behind a knife edge (also called a half-plane).

Multiple Knife-edge Diffraction:

In many practical situations, especially in hilly terrain, the propagation path may consist of more than one obstruction, in which case the total diffraction loss due to all of the obstacles must be computed. Burlington suggested that the series of obstacles be replaced by a single equivalent obstacle so that the path loss can be obtained using single knife-edge diffraction models. This method, illustrated in Figure 3.15, oversimplifies the calculations and often provides very optimistic estimates of the received signal strength. In a more rigorous treatment, Millington et al. gave a

wave-theory solution for the field behind two knife edges in series. This solution is very useful and can be applied easily for predicting diffraction losses due to two knife edges. However, extending this to more than two knife edges becomes a formidable mathematical problem. Many models that are mathematically less complicated have been developed to estimate the diffraction losses due to multiple obstructions.

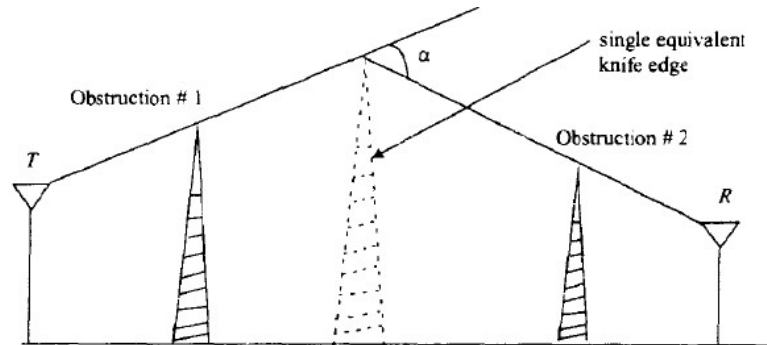


Figure 3.15
Bullington's construction of an equivalent knife edge [From [Bul47] © IEEE].

Scattering:

The actual received signal in a mobile radio environment is often stronger than what is predicted by reflection and diffraction models alone. This is because when a radio wave impinges on a rough surface, the reflected energy is spread out (diffused) in all directions due to scattering. Objects such as lamp posts and trees tend to scatter energy in all directions, thereby providing additional radio energy at a receiver. Flat surfaces that have much larger dimension than a wavelength may be modeled as reflective surfaces. However, the roughness of such surfaces often induces propagation effects different from the specular reflection described earlier in this chapter.

Rayleigh criterion: used for testing surface roughness. A surface is considered smooth if its minimum to maximum protuberance (bumps) h is less than critical height $h_c = \lambda / 8 \sin \theta_i$.

Scattering path loss factor ρ_s is given by $\rho_s = \exp[-8[(\pi \cdot \sigma_h \cdot \sin \theta_i) / \lambda]^2]$

Where h is surface height and σ_h is standard deviation of surface height about mean surface height. For rough surface, the flat surface reflection coefficient is multiplied by scattering loss factor ρ_s to account for diminished electric field.

Reflected E-fields for $h > h_c$ for rough surface can be calculated as $\Gamma_{\text{rough}} = \rho_s \Gamma$

A surface is considered smooth if its minimum to maximum protuberance h is less than h_c , and is considered rough if the protuberance is greater than h_c . For rough surfaces, the flat surface reflection coefficient needs to be multiplied by a scattering loss factor, ρ_s to account for the diminished reflected field.

OUTDOORPROPAGATIONMODELS

Based on the coverage area, the outdoor propagation environment may be divided into three categories

1. Propagation in Macrocells
2. Propagation in Microcells

Outdoor radio transmission takes place over an irregular terrain. The terrain profile must be taken into consideration for estimating the path loss.

e.g. trees, buildings and hills must be taken into consideration

Longley-Ryce Model:

The Longley-Ryce model is applicable to point-to-point communication systems in the frequency range from 40 MHz to 100 GHz, over different kinds of terrain. The median transmission loss is predicted using the path geometry of the terrain profile and the refractivity of the troposphere. Geometric optics techniques (primarily the 2-ray ground reflection model) are used to predict signal strengths within the radio horizon. Diffraction losses over isolated obstacles are estimated using the Fresnel-Kirchoff knife-edge models. Forward scatter theory is used to make troposcatter predictions over long distances.

The Longley-Ryce method operates in two modes. When a detailed terrain path profile is available, the path-specific parameters can be easily determined and the prediction is called a point-to-point mode prediction. On the other hand, if the terrain path profile is not available, the Longley-Ryce method provides techniques to estimate the path-specific parameters, and such a prediction is called an area mode prediction.

Okumura Model:

Okumura's model is one of the most widely used models for signal prediction in urban areas. This model is applicable for frequencies in the range 150 MHz to 1920 MHz (although it is typically extrapolated up to 3000 MHz) and distances of 1 km to 100 km. It can be used for base station antenna heights ranging from 30 m to 1000 m. Okumura developed a set of curves giving the median attenuation relative to free space (A_{mu}), in an urban area over a quasi-smooth terrain with a base station effective antenna height (h_{te}) of 200 m and a mobile antenna height (h_{re}) of 3 m. These curves were developed from extensive measurements using vertical omni-directional antennas at both the base and mobile, and are plotted as a function of frequency in the range 100 MHz to 1920 MHz and as a function of distance from the base station in the range 1 km to 100 km. To determine path loss using Okumura's model, the free space path loss between the points of interest is first determined, and then the value of $A_{mu}(f, d)$ (as read from the curves) is added to it along with correction factors to account for the type of terrain. The model can be expressed as

$$L_{50}(dB) = LF + A_{mu}(f, d) - G(te) - G(re) - G_{AREA}$$

where L_{50} is the 50th percentile (i.e., median) value of propagation path loss, L_F is the free space propagation loss, A_{mu} is the median attenuation relative to free space, $G(h_{te})$ is the base station antenna height gain factor, $G(h_{re})$ is the mobile antenna height gain factor, and $G_A R E A$ is the gain due to the type of environment. Note that the antenna height gains are strictly a function of height and have nothing to do with antenna patterns.

$$G(h_{te}) = 20 \log \left(\frac{h_{te}}{200} \right) \quad 1000 \text{ m} > h_{te} > 30 \text{ m}$$

$$G(h_{re}) = 10 \log \left(\frac{h_{re}}{3} \right) \quad h_{re} \leq 3 \text{ m}$$

$$G(h_{re}) = 20 \log \left(\frac{h_{re}}{3} \right) \quad 10 \text{ m} > h_{re} > 3 \text{ m}$$

Hata Model:

The Hata model [Hat90] is an empirical formulation of the graphical path loss data provided by Okumura, and is valid from 150 MHz to 1500 MHz. Hata presented the urban area propagation loss as a standard formula and supplied correction equations for application to other situations. The standard formula for median path loss in urban areas is given by

$$L_{50}(\text{urban})(\text{dB}) = 69.55 + 26.16 \log f_c - 13.82 \log h_{te} - a(h_{re}) + (44.9 - 6.55 \log h_{te}) \log d$$

where f_c is the frequency (in MHz) from 150 MHz to 1500 MHz, h_{te} is the effective transmitter (base station) antenna height (in meters) ranging from 30 m to 200 m, h_{re} is the effective receiver (mobile) antenna height (in meters) ranging from 1 m to 10 m, d is the T-R separation distance (in km), and $a(h_{re})$ is the correction factor for effective mobile antenna height which is a function of the size of the coverage area. For a small to medium sized city, the mobile antenna correction factor is given by

$$a(h_{re}) = (1.11 \log f_c - 0.7) h_{re} - (1.56 \log f_c - 0.8) \text{ dB for a large city, it is given by}$$

$$a(h_{re}) = 8.29 (\log 1.54 h_{re})^2 - 1.1 \text{ dB for } f_c \leq 300 \text{ MHz} \quad a(h_{re}) = 3.2 (\log 11.75 h_{re})^2 - 4.97 \text{ dB for } f_c > 300 \text{ MHz}$$

To obtain the path loss in a suburban area the standard Hata formula in equations are modified as $L_{50}(\text{dB}) = L_{50}(\text{urban}) - 2[\log(f_c/28)]^2 - 5.4$

and for path loss in open rural areas, the formula is modified as $L_{50}(\text{dB}) = L_{50}(\text{urban}) - 4.78(\log f_c)^2 + 18.33 \log f_c - 40.94$

PCS Extension to Hata Model

The European Co-operative for Scientific and Technical research (EURO-COST) formed the COST-231 working committee to develop an extended version of the Hata model. COST-231 proposed the following formula to extend Hata's model to 2 GHz. The proposed model for path loss is [EUR91]

$$L_{50}(urban) = 46.3 + 33.9 \log f_c - 13.82 \log h_{te} - a(h_{re}) + (44.9 - 6.55 \log h_{te}) \log d + C_M \quad (3.87)$$

where $a(h_{re})$ is defined in equations (3.83), (3.84.a), and (3.84.b) and

$$C_M = \begin{cases} 0 \text{ dB} & \text{for medium sized city and suburban areas} \\ 3 \text{ dB} & \text{for metropolitan centers} \end{cases} \quad (3.88)$$

The COST-231 extension of the Hata model is restricted to the following range of parameters:

$$\begin{aligned} f &: 1500 \text{ MHz to } 2000 \text{ MHz} \\ h_{te} &: 30 \text{ m to } 200 \text{ m} \\ h_{re} &: 1 \text{ m to } 10 \text{ m} \\ d &: 1 \text{ km to } 20 \text{ km} \end{aligned}$$

Walfisch and Bertoni Model

A model developed by Walfisch and Bertoni [Wal88] considers the impact of rooftops and building height by using diffraction to predict average signal strength at street level. The model considers the path loss, S , to be a product of three factors.

$$S = P_0 Q^2 P_1 \quad (3.89)$$

where P_0 represents free space path loss between isotropic antennas given by

$$P_0 = \left(\frac{\lambda}{4\pi R} \right)^2 \quad (3.90)$$

The factor Q^2 gives the reduction in the rooftop signal due to the row of buildings which immediately shadow the receiver at street level. The P_1 term is

based upon diffraction and determines the signal loss from the rooftop to the street.

In dB, the path loss is given by

$$S(\text{dB}) = L_0 + L_{rts} + L_{ms} \quad (3.91)$$

where L_0 represents free space loss, L_{rts} represents the “rooftop-to-street diffraction and scatter loss”, and L_{ms} denotes multiscreen diffraction loss due to the rows of buildings [Xia92]. Figure 3.25 illustrates the geometry used in the Walfisch Bertoni model [Wal88], [Mac93]. This model is being considered for use by ITU-R in the IMT-2000 standards activities.

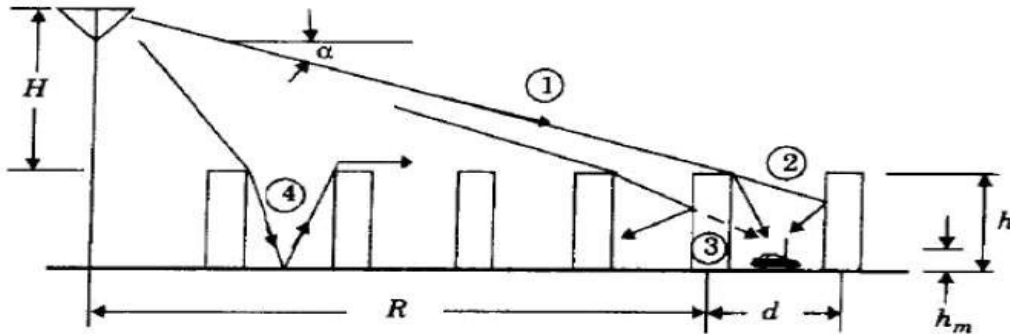


Figure 3.25
Propagation geometry for model proposed by Walfisch and Bertoni [From [Wal88] © IEEE].

Wideband PCS Microcell Model

Work by Feuerstein, et.al. in 1991 used a 20 MHz pulsed transmitter at 1900 MHz to measure path loss, outage, and delay spread in typical microcellular systems in San Francisco and Oakland. Using base station antenna heights of 3.7 m, 8.5 m, and 13.3 m, and a mobile receiver with an antenna height of 1.7 m above ground, statistics for path loss, multipath, and coverage area were developed from extensive measurements in line-of-sight (LOS) and obstructed (OBS) environments [Feu94]. This work revealed that a 2-ray ground reflection model (shown in Figure 3.7) is a good estimate for path loss in LOS microcells, and a simple log-distance path loss model holds well for OBS microcell environments.

For a flat earth ground reflection model, the distance d_f at which the first Fresnel zone just becomes obstructed by the ground (first Fresnel zone clearance) is given by

$$d_f = \frac{1}{\lambda} \sqrt{(\Sigma^2 - \Delta^2)^2 - 2(\Sigma^2 + \Delta^2) \left(\frac{\lambda}{2}\right)^2 + \left(\frac{\lambda}{2}\right)^4} \quad (3.92.a)$$

$$= \frac{1}{\lambda} \sqrt{16h_t^2 h_r^2 - \lambda^2 (h_t^2 + h_r^2) + \frac{\lambda^4}{16}}$$

For LOS cases, a double regression path loss model that uses a regression breakpoint at the first Fresnel zone clearance was shown to fit well to measurements. The model assumes omnidirectional vertical antennas and predicts average path loss as

$$PL(d) = \begin{cases} 10n_1 \log(d) + p_1 & \text{for } 1 < d < d_f \\ 10n_2 \log(d/d_f) + 10n_1 \log d_f + p_1 & \text{for } d > d_f \end{cases} \quad (3.92.b)$$

where p_1 is equal to $PL(d_0)$ (the path loss in decibels at the reference distance of $d_0 = 1$ m), d is in meters and n_1, n_2 are path loss exponents which are a function of transmitter height, as given in Figure 3.26. It can easily be shown that at 1900 MHz, $p_1 = 38.0$ dB.

For the OBS case, the path loss was found to fit the standard log-distance path loss law of equation (3.69.a)

$$PL(d) [dB] = 10n \log(d) + p_1 \quad (3.92.c)$$

INDOORPROPAGATIONMODELS

With the advent of Personal Communication Systems (PCS), there is a great deal of interest in characterizing radio propagation inside buildings. The indoor radio channel differs from the traditional mobile radio channel in two aspects - the distances covered are much smaller, and the variability of the environment is much greater for a much smaller range of separation distances. It has been observed that propagation within buildings is strongly influenced by specific features such as the layout of the building, the construction materials, and the building type. This section outlines models for path loss within buildings.

Indoor radio propagation is dominated by the same mechanisms as outdoor: reflection, diffraction, and scattering. However, conditions are much more variable. For example, signal levels vary greatly depending on whether interior doors are open or closed inside a building. Where antennas are mounted also impacts large-scale propagation. Antennas mounted at desk level in a partitioned office receive vastly different signals than those mounted on the ceiling. Also, the smaller propagation distances make it more difficult to insure far-field radiation for all receiver locations and types of antennas.

Partition Losses (same floor):

Buildings have a wide variety of partitions and obstacles which form the internal and external structure. Houses typically use a wood frame partition with plaster board to form internal walls and

have wood or non-reinforced concrete between floors. Office buildings, on the other hand, often have large open areas (open plan) which are constructed by using moveable office partitions so that the space may be reconfigured easily, and use metal reinforced concrete between floors. Partitions that are formed as part of the building structure are called hard partitions, and partitions that may be moved and which do not span to the ceiling are called soft partitions. Partitions vary widely in their physical and electrical characteristics, making it difficult to apply general models to specific indoor installations.

Partition Losses between Floors:

The losses between floors of a building are determined by the external dimensions and materials of the building, as well as the type of construction used to create the floors and the external surroundings. Even the number of windows in a building and the presence of tinting (which attenuates radio energy) can impact the loss between floors. It can be seen that for all three buildings, the attenuation between one floor of the building is greater than the incremental attenuation caused by each additional floor. After about five or six floor separations, very little additional path loss is experienced.

Log-distance Path Loss Model:

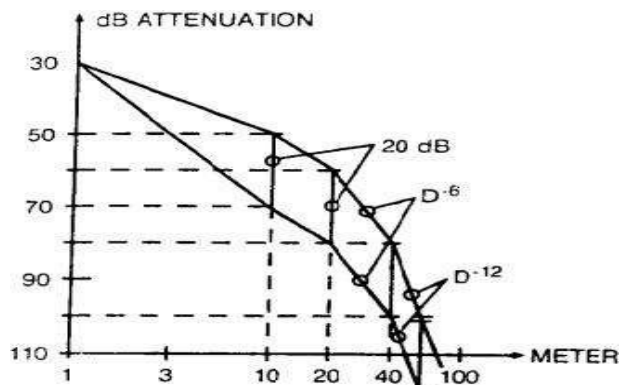
According to this model the received power at distance d is given by,

$$PL (dB) = PL (d_0) + 10n \log\left(\frac{d}{d_0}\right) + X_\sigma$$

The value of n varies with propagation environments. The value of n is 2 for free space. The value of n varies from 4 to 6 for obstruction of building, and 3 to 5 for urban scenarios. The important factor is to select the correct reference distance d_0 . For large cell area it is 1 Km, while for micro-cell system it varies from 10m-1m.

Ericsson multiple breakpoint model:

It was obtained by measurements in a multiple floor office building. It has 4 breakpoints and considers both an upper and lower bound on path loss. It assumes that there is 30dB attenuation at $d_0 = 1m$ which is accurate for $f = 900MHz$ & unity gain antennas.



AttenuationFactorModel:

The attenuation factor model incorporates a special path loss exponent and a floor attenuation factor to provide an estimate of indoor path loss. NSF is the path loss exponent for a same floor measurement and FAF is a floor attenuation factor based on the number of floors between transmitter and receiver.

RayTracingandSiteSpecificModeling

In physics, ray tracing is a method for calculating the path of waves or particles through a system with regions of varying propagation velocity, absorption characteristics, and reflecting surfaces. Under these circumstances, wave fronts may bend, change direction, or reflect off surfaces, complicating analysis. Ray tracing solves the problem by repeatedly advancing idealized narrow beams called rays through the medium by discrete amounts. Simple problems can be analyzed by propagating a few rays using simple arithmetic. More detailed analysis can be performed by using a computer to propagate many rays.

When applied to problems of electromagnetic radiation, ray tracing often relies on approximate solutions to Maxwell's equations that are valid as long as the light waves propagate through and around objects whose dimensions are much greater than the light's wavelength. Ray theory does not describe phenomena such as interference and diffraction, which require wave theory (involving the phase of the wave).

UNIT-III

MobileRadioPropagation:Small–ScaleFadingandMultipath

Small Scale Multipath propagation-Factors influencing small scale fading, Doppler shift. Impulse Response Model of a multipath channel-Relationship between Bandwidth and Received power Small-Scale MultipathMeasurements-DirectRFPulseSystem,SpreadSpectrum SlidingCorrelator Channel Sounding, Frequency Domain Channels Sounding Parameters of MobileMultipathChannels-TimeDispersionParameters,CoherenceBandwidth,DopplerSpread and Coherence Time. Types of Small-Scale Fading-Fading effects Due to Multipath Time Delay Spread, Flat fading, Frequency selective fading.Fading effects Due to Doppler Spread-Fast fading, slow fading Statistical Models for multipath Fading Channels-Clarke's model for flat fading, spectral shape due to Doppler spread in Clarke's model, Simulation of Clarke and Gans Fading Model, Level crossing and fading statistics, Two-ray Rayleigh Fading Model.

SmallScale Multipath propagation:

Multipath in the radio channel creates small-scale fading effects. The three most important effects are:

- **Rapid changes in signal strength over a small travel distance or time interval**
- **Random frequency modulation due to varying Doppler shifts on different multipath signals**
- **Time dispersion (echoes) caused by multipath propagation delays.**

In built-up urban areas, fading occurs because the height of the mobile antennas are well below the height of surrounding structures, so there is no single line-of-sight path to the base station. Even when a line-of-sight exists, multipath still occurs due to reflections from the ground and surrounding structures. The incoming radio waves arrive from different directions with different propagation delays. The signal received by the mobile at any point in space may consist of a large number of plane waves having randomly distributed amplitudes, phases, and angles of arrival. These multipath components combine vectorially at the receiver antenna, and can cause the signal received by the mobile to distort or fade. Even when a mobile receiver is stationary, the received signal may fade due to movement of surrounding objects in the radio channel.

Factorsinfluencingsmallscalefading:

Many physical factors in the radio propagation channel influence small-scale fading. These include the following:

- **Multipath propagation** — The presence of reflecting objects and scatterers in the channel creates a constantly changing environment that dissipates the signal energy in amplitude, phase, and time. These effects result in multiple versions of the transmitted signal that arrive at the receiving antenna, displaced with respect to one another in time and spatial orientation. The random phase and amplitudes of the different multipath components cause fluctuations in signal strength, thereby inducing small-scale fading, signal distortion, or both. Multipath propagation often lengthens the time required for the baseband portion of the signal to reach the receiver which can cause signal smearing due to intersymbol interference.
- **Speed of the mobile** — The relative motion between the base station and the mobile results in random frequency modulation due to different Doppler shifts on each of the multipath components. Doppler shift will be positive or negative depending on whether the mobile receiver is moving toward or away from the base station.

- **Speed of surrounding objects** — If objects in the radio channel are in motion, they induce a time varying Doppler shift on multipath components. If the surrounding objects move at a greater rate than the mobile, then this effect dominates the small-scale fading. Otherwise, motion of surrounding objects may be ignored, and only the speed of the mobile need be considered.
- **The transmission bandwidth of the signal** — If the transmitted radio signal bandwidth is greater than the “bandwidth” of the multipath channel, the received signal will be distorted, but the received signal strength will not fade much over a local area (i.e., the small-scale signal fading will not be significant). As will be shown, the bandwidth of the channel can be quantified by the *coherence bandwidth* which is related to the specific multipath structure of the channel. The coherence bandwidth is a measure of the maximum frequency difference for which signals are still strongly correlated in amplitude. If the transmitted signal has a narrow bandwidth as compared to the channel, the amplitude of the signal will change rapidly, but the signal will not be distorted in time. Thus, the statistics of small-scale signal strength and the likelihood of signal smearing appearing over small-scale distances are very much related to the specific amplitudes and delays of the multipath channel, as well as the bandwidth of the transmitted signal.

Dopplershift:

Consider a mobile moving at a constant velocity v , along a path segment having length d between points X and Y, while it receives signals from a remote source S, as illustrated in Figure 4.1. The difference in path lengths traveled by the wave from source S to the mobile at points X and Y is $\Delta l = d \cos \theta = v \Delta t \cos \theta$, where Δt is the time required for the mobile to travel from X to Y, and θ is assumed to be the same at points X and Y since the source is assumed to be very far away. The phase change in the received signal due to the difference in path lengths is therefore

$$\Delta \phi = \frac{2\pi \Delta l}{\lambda} = \frac{2\pi v \Delta t}{\lambda} \cos \theta \quad (4.1)$$

and hence the apparent change in frequency, or Doppler shift, is given by f_d , where

$$f_d = \frac{1}{2\pi} \cdot \frac{\Delta \phi}{\Delta t} = \frac{v}{\lambda} \cdot \cos \theta \quad (4.2)$$

Equation (4.2) relates the Doppler shift to the mobile velocity and the spatial angle between the direction of motion of the mobile and the direction of arrival of the wave. It can be seen from equation (4.2) that if the mobile is moving toward the direction of arrival of the wave, the Doppler shift is positive (i.e., the apparent received frequency is increased), and if the mobile is moving away from the direction of arrival of the wave, the Doppler shift is negative (i.e. the

apparent received frequency is decreased). As shown in section 4.7.1, multipath components from a CW signal which arrive from different directions contribute to Doppler spreading of the received signal, thus increasing the signal bandwidth.

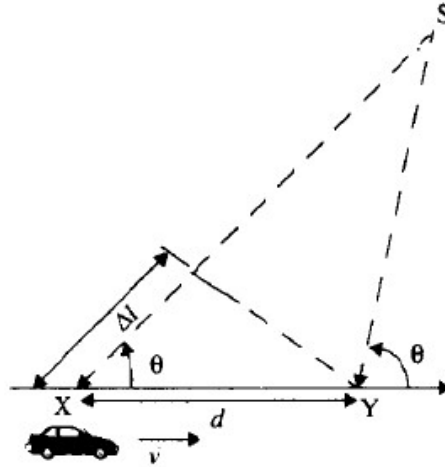


Figure 4.1
Illustration of Doppler effect.

Example 4.1

Consider a transmitter which radiates a sinusoidal carrier frequency of 1850 MHz. For a vehicle moving 60 mph, compute the received carrier frequency if the mobile is moving (a) directly towards the transmitter, (b) directly away from the transmitter, (c) in a direction which is perpendicular to the direction of arrival of the transmitted signal.

Solution to Example 4.1

Given:

$$\text{Carrier frequency } f_c = 1850 \text{ MHz}$$

$$\text{Therefore, wavelength } \lambda = c/f_c = \frac{3 \times 10^8}{1850 \times 10^6} = 0.162 \text{ m}$$

$$\text{Vehicle speed } v = 60 \text{ mph} = 26.82 \text{ m/s}$$

(a) The vehicle is moving directly towards the transmitter.

The Doppler shift in this case is positive and the received frequency is given by equation (4.2)

$$f = f_c + f_d = 1850 \times 10^6 + \frac{26.82}{0.162} = 1850.00016 \text{ MHz}$$

(b) The vehicle is moving directly away from the transmitter.

The Doppler shift in this case is negative and hence the received frequency is given by

$$f = f_c - f_d = 1850 \times 10^6 - \frac{26.82}{0.162} = 1849.999834 \text{ MHz}$$

(c) The vehicle is moving perpendicular to the angle of arrival of the transmitted signal.

In this case, $\theta = 90^\circ$, $\cos\theta = 0$, and there is no Doppler shift.

The received signal frequency is the same as the transmitted frequency of 1850 MHz.

Impulse Response Model of a Multipath Channel

The small-scale variations of a mobile radio signal can be directly related to the impulse response of the mobile radio channel. The impulse response is a wideband channel characterization and contains all information necessary to simulate or analyze any type of radio transmission through the channel. This stems from the fact that a mobile radio channel may be modeled as a linear filter with a time varying impulse response, where the time variation is due to receiver motion in space. The filtering nature of the channel is caused by the summation of amplitudes and delays of the multiple arriving waves at any instant of time. The impulse response is a useful characterization of the channel, since it may be used to predict and compare the performance of many different mobile communication systems and transmission bandwidths for a particular mobile channel condition.

To show that a mobile radio channel may be modeled as a linear filter with a time varying impulse response, consider the case where time variation is due strictly to receiver motion in space. This is shown in Figure 4.2.



Figure 4.2

The mobile radio channel as a function of time and space.

In Figure 4.2, the receiver moves along the ground at some constant velocity v . For a fixed position d , the channel between the transmitter and the receiver can be modeled as a linear time invariant system. However, due to the different multipath waves which have propagation delays which vary over different spatial locations of the receiver, the impulse response of the linear time invariant channel should be a function of the position of the receiver. That is, the channel impulse response can be expressed as $h(d,t)$. Let $x(t)$ represent the transmitted signal, then the received signal $y(d,t)$ at position d can be expressed as a convolution of $x(t)$ with $h(d,t)$.

$$y(d, t) = x(t) \otimes h(d, t) = \int_{-\infty}^{\infty} x(\tau) h(d, t - \tau) d\tau \quad (4.3)$$

For a causal system, $h(d, t) = 0$ for $t < 0$, thus equation (4.3) reduces to

$$y(d, t) = \int_{-\infty}^t x(\tau) h(d, t - \tau) d\tau \quad (4.4)$$

Since the receiver moves along the ground at a constant velocity v , the position of the receiver can be expressed as

$$d = vt \quad (4.5)$$

Substituting (4.5) in (4.4), we obtain

$$y(vt, t) = \int_{-\infty}^t x(\tau) h(vt, t - \tau) d\tau \quad (4.6)$$

Since v is a constant, $y(vt, t)$ is just a function of t . Therefore, equation (4.6) can be expressed as

$$y(t) = \int_{-\infty}^t x(\tau) h(vt, t - \tau) d\tau = x(t) \otimes h(vt, t) = x(t) \otimes h(d, t) \quad (4.7)$$

From equation (4.7) it is clear that the mobile radio channel can be modeled as a linear time varying channel, where the channel changes with time and distance.

Since v may be assumed constant over a short time (or distance) interval, we may let $x(t)$ represent the transmitted bandpass waveform, $y(t)$ the received waveform, and $h(t, \tau)$ the impulse response of the time varying multipath radio channel. The impulse response $h(t, \tau)$ completely characterizes the channel and is a function of both t and τ . The variable t represents the time variations due to motion, whereas τ represents the channel multipath delay for a fixed value of t . One may think of τ as being a vernier adjustment of time. The received signal $y(t)$ can be expressed as a convolution of the transmitted signal $x(t)$ with the channel impulse response (see Figure 4.3a).

$$y(t) = \int_{-\infty}^{\infty} x(\tau) h(t, \tau) d\tau = x(t) \otimes h(t, \tau) \quad (4.8)$$

If the multipath channel is assumed to be a bandlimited bandpass channel, which is reasonable, then $h(t, \tau)$ may be equivalently described by a complex baseband impulse response $h_b(t, \tau)$ with the input and output being the complex envelope representations of the transmitted and received signals, respectively (see Figure 4.3b). That is,

$$r(t) = c(t) \otimes \frac{1}{2} h_b(t, \tau) \quad (4.9)$$

$$\begin{array}{ll}
 x(t) & \blacktriangleright \quad h(t, \tau) = \text{Re} \left\{ h_b(t, \tau) e^{j\omega_c t} \right\} & \blacktriangleright y(t) \\
 & & y(t) = \text{Re} \{ r(t) e^{j\omega_c t} \} \\
 & \text{(a)} & y(t) = x(t) \otimes h(t) \\
 \\
 c(t) & \blacktriangleright \quad \frac{1}{2} h_b(t, \tau) & \blacktriangleright r(t) \\
 & & \frac{1}{2} r(t) = \frac{1}{2} c(t) \otimes \frac{1}{2} h_b(t) \\
 & \text{(b)} &
 \end{array}$$

Figure 4.3

(a) Bandpass channel impulse response model.

(b) Baseband equivalent channel impulse response model.

where $c(t)$ and $r(t)$ are the complex envelopes of $x(t)$ and $y(t)$, defined as

$$x(t) = \text{Re} \{ c(t) \exp(j2\pi f_c t) \} \quad (4.10)$$

$$y(t) = \text{Re} \{ r(t) \exp(j2\pi f_c t) \} \quad (4.11)$$

The factor of $1/2$ in equation (4.9) is due to the properties of the complex envelope, in order to represent the passband radio system at baseband. The low-pass characterization removes the high frequency variations caused by the carrier, making the signal analytically easier to handle. It is shown by Couch

[Cou93] that the average power of a bandpass signal $\overline{x^2(t)}$ is equal to $\frac{1}{2} \overline{|c(t)|^2}$,

where the overbar denotes ensemble average for a stochastic signal, or time average for a deterministic or ergodic stochastic signal.

It is useful to discretize the multipath delay axis τ of the impulse response into equal time delay segments called *excess delay bins*, where each bin has a time delay width equal to $\tau_{i+1} - \tau_i$, where τ_0 is equal to 0, and represents the first arriving signal at the receiver. Letting $i = 0$, it is seen that $\tau_1 - \tau_0$ is equal to the time delay bin width given by $\Delta\tau$. For convention, $\tau_0 = 0$, $\tau_1 = \Delta\tau$, and $\tau_i = i\Delta\tau$, for $i = 0$ to $N - 1$, where N represents the total number of possible equally-spaced multipath components, including the first arriving component. Any number of multipath signals received within the i th bin are represented by a single resolvable multipath component having delay τ_i . This technique of quantizing the delay bins determines the time delay resolution of the channel model, and the useful frequency span of the model can be shown to be $1/(2\Delta\tau)$. That is, the model may be used to analyze transmitted signals having bandwidths which are less than $1/(2\Delta\tau)$. Note that $\tau_0 = 0$ is the excess time delay

of the first arriving multipath component, and neglects the propagation delay between the transmitter and receiver. *Excess delay* is the relative delay of the i th multipath component as compared to the first arriving component and is given by τ_i . The *maximum excess delay* of the channel is given by $N\Delta\tau$.

Since the received signal in a multipath channel consists of a series of attenuated, time-delayed, phase shifted replicas of the transmitted signal, the baseband impulse response of a multipath channel can be expressed as

$$h_b(t, \tau) = \sum_{i=0}^{N-1} a_i(t, \tau) \exp[j(2\pi f_c \tau_i(t) + \phi_i(t, \tau))] \delta(\tau - \tau_i(t)) \quad (4.12)$$

where $a_i(t, \tau)$ and $\tau_i(t)$ are the real amplitudes and excess delays, respectively, of i th multipath component at time t [Tur72]. The phase term $2\pi f_c \tau_i(t) + \phi_i(t, \tau)$ in (4.12) represents the phase shift due to free space propagation of the i th multipath component, plus any additional phase shifts which are encountered in the channel. In general, the phase term is simply represented by a single variable $\theta_i(t, \tau)$ which lumps together all the mechanisms for phase shifts of a single multipath component within the i th excess delay bin. Note that some excess delay bins may have no multipath at some time t and delay τ_i , since $a_i(t, \tau)$ may be zero. In equation (4.12), N is the total possible number of multipath components (bins), and $\delta(\cdot)$ is the unit impulse function which determines the specific multipath bins that have components at time t and excess delays τ_i . Figure 4.4 illustrates an example of different snapshots of $h_b(t, \tau)$, where t varies into the page, and the time delay bins are quantized to widths of $\Delta\tau$.

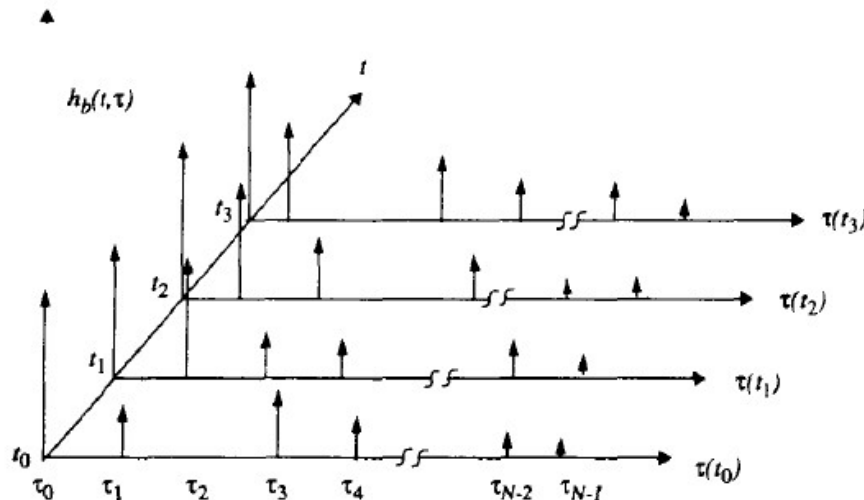


Figure 4.4
An example of the time varying discrete-time impulse response model for a multipath radio channel.

If the channel impulse response is assumed to be time invariant, or is at least wide sense stationary over a small-scale time or distance interval, then the channel impulse response may be simplified as

$$h_b(\tau) = \sum_{i=0}^{N-1} a_i \exp(-j\theta_i) \delta(\tau - \tau_i) \quad (4.13)$$

When measuring or predicting $h_b(\tau)$, a probing pulse $p(t)$ which approximates a delta function is used at the transmitter. That is

$$p(t) \approx \delta(t - \tau) \quad (4.14)$$

is used to sound the channel to determine $h_b(\tau)$.

For small-scale channel modeling, the *power delay profile* of the channel is found by taking the spatial average of $|h_b(t;\tau)|^2$ over a local area. By making several local area measurements of $|h_b(t;\tau)|^2$ in different locations, it is possible to build an ensemble of power delay profiles, each one representing a possible small-scale multipath channel state [Rap91a].

Based on work by Cox [Cox72], [Cox75], if $p(t)$ has a time duration much smaller than the impulse response of the multipath channel, $p(t)$ does not need to be deconvolved from the received signal $r(t)$ in order to determine relative multipath signal strengths. The received power delay profile in a local area is given by

$$P(t;\tau) \approx k|h_b(t;\tau)|^2 \quad (4.15)$$

and many snapshots of $|h_b(t;\tau)|^2$ are typically averaged over a local (small-scale) area to provide a single time-invariant multipath power delay profile $P(\tau)$. The gain k in equation (4.15) relates the transmitted power in the probing pulse $p(t)$ to the total power received in a multipath delay profile.

Relationship Between Bandwidth and Received Power

In actual wireless communication systems, the impulse response of a multipath channel is measured in the field using channel sounding techniques. We now consider two extreme channel sounding cases as a means of demonstrating how the small-scale fading behaves quite differently for two signals with different bandwidths in the identical multipath channel.

Consider a pulsed, transmitted RF signal of the form

$$x(t) = \text{Re} \{ p(t) \exp(j2\pi f_c t) \}$$

where $p(t)$ is a repetitive baseband pulse train with very narrow pulse width T_{bb} and repetition period T_{REP} which is much greater than the maximum measured excess delay τ_{max} in the channel. Now let

$$p(t) = 2\sqrt{\tau_{max}/T_{bb}} \text{ for } 0 \leq t \leq T_{bb}$$

and let $p(t)$ be zero elsewhere for all excess delays of interest. The low pass channel output $r(t)$ closely approximates the impulse response $h_s(t)$ and is given by

$$\begin{aligned} r(t) &= \frac{1}{2} \sum_{i=0}^{N-1} a_i (\exp(-j\theta_i)) \cdot p(t - \tau_i) \\ &= \sum_{i=0}^{N-1} a_i \exp(-j\theta_i) \cdot \sqrt{\frac{\tau_{max}}{T_{bb}}} \text{rect} \left[t - \frac{T_{bb}}{2} - \tau_i \right] \end{aligned} \quad (4.16)$$

To determine the received power at some time t_0 , the power $|r(t_0)|^2$ is measured. The quantity $|r(t_0)|^2$ is called the *instantaneous multipath power delay profile* of the channel, and is equal to the energy received over the time duration of the multipath delay divided by τ_{max} . That is, using equation (4.16)

$$\begin{aligned}
|r(t_0)|^2 &= \frac{1}{\tau_{max}} \int_0^{\tau_{max}} r(t) \times r^*(t) dt & (4.17) \\
&= \frac{1}{\tau_{max}} \int_0^{\tau_{max}} \frac{1}{4} \text{Re} \left\{ \sum_{j=0}^{N-1} \sum_{i=0}^{N-1} a_j(t_0) a_i(t_0) p(t-\tau_j) p(t-\tau_i) \exp(-j(\theta_j-\theta_i)) \right\} dt
\end{aligned}$$

Note that if all the multipath components are resolved by the probe $p(t)$, then $|\tau_j - \tau_i| > T_{bb}$ for all $j \neq i$, and

$$\begin{aligned}
|r(t_0)|^2 &= \frac{1}{\tau_{max}} \int_0^{\tau_{max}} \frac{1}{4} \left(\sum_{k=0}^{N-1} a_k^2(t_0) p^2(t-\tau_k) \right) dt & (4.18) \\
&= \frac{1}{\tau_{max}} \sum_{k=0}^{N-1} a_k^2(t_0) \int_0^{\tau_{max}} \left\{ \sqrt{\frac{\tau_{max}}{T_{bb}}} \text{rect} \left[t - \frac{T_{bb}}{2} - \tau_k \right] \right\}^2 dt \\
&= \sum_{k=0}^{N-1} a_k^2(t_0)
\end{aligned}$$

For a wideband probing signal $p(t)$, T_{bb} is smaller than the delays between multipath components in the channel, and equation (4.18) shows that the total received power is simply related to the sum of the powers in the individual multipath components, and is scaled by the ratio of the probing pulse's width and amplitude, and the maximum observed excess delay of the channel. Assuming that the received power from the multipath components forms a random process where each component has a random amplitude and phase at any time t , the average small-scale received power for the wideband probe is found from equation (4.17) as

$$E_{a,\theta}[P_{WB}] = E_{a,\theta} \left[\sum_{i=0}^{N-1} |a_i \exp(j\theta_i)|^2 \right] \approx \sum_{i=0}^{N-1} \overline{a_i^2} \quad (4.19)$$

In equation (4.19), $E_{a,\theta}[\cdot]$ denotes the ensemble average over all possible values of a_i and θ_i in a local area, and the overbar denotes sample average over a local measurement area which is generally measured using multipath measurement equipment. The striking result of equations (4.18) and (4.19) is that if a transmitted signal is able to resolve the multipaths, then the *small-scale received power is simply the sum of the powers received in each multipath component*. In practice, the amplitudes of individual multipath components do not fluctuate widely in a local area. Thus, the received power of a wideband signal such as $p(t)$ does not fluctuate significantly when a receiver is moved about a local area [Rap89].

Now, instead of a pulse, consider a CW signal which is transmitted into the exact same channel, and let the complex envelope be given by $c(t) = 2$. Then, the instantaneous complex envelope of the received signal is given by the phasor sum

$$r(t) = \sum_{i=0}^{N-1} a_i \exp(j\theta_i(t, \tau)) \quad (4.20)$$

and the instantaneous power is given by

$$|r(t)|^2 = \left| \sum_{i=0}^{N-1} a_i \exp(j\theta_i(t, \tau)) \right|^2 \quad (4.21)$$

As the receiver is moved over a local area, the channel changes, and the received signal strength will vary at a rate governed by the fluctuations of a_i and θ_i . As mentioned earlier, a_i varies little over local areas, but θ_i will vary greatly due to changes in propagation distance over space, resulting in large fluctuations of $r(t)$ as the receiver is moved over small distances (on the order of a wavelength). That is, since $r(t)$ is the phasor sum of the individual multipath components, the instantaneous phases of the multipath components cause the large fluctuations which typifies small-scale fading for CW signals. The average received power over a local area is then given by

$$E_{a,\theta}[P_{CW}] = E_{a,\theta} \left[\left| \sum_{i=0}^{N-1} a_i \exp(j\theta_i) \right|^2 \right] \quad (4.22)$$

$$E_{a,\theta}[P_{CW}] \approx \frac{\left[a_0 e^{j\theta_0} + a_1 e^{j\theta_1} + \dots + a_{N-1} e^{j\theta_{N-1}} \right] \times \left[a_0 e^{-j\theta_0} + a_1 e^{-j\theta_1} + \dots + a_{N-1} e^{-j\theta_{N-1}} \right]}{\quad} \quad (4.23)$$

$$E_{a,\theta}[P_{CW}] \approx \sum_{i=0}^{N-1} \overline{a_i^2} + 2 \sum_{i=0}^{N-1} \sum_{i,j \neq i}^N r_{ij} \overline{\cos(\theta_i - \theta_j)} \quad (4.24)$$

where r_{ij} is the path amplitude correlation coefficient defined to be

$$r_{ij} = E_a[a_i a_j] \quad (4.25)$$

and the overbar denotes time average for CW measurements made by a mobile receiver over the local measurement area [Rap89]. Note that when $\overline{\cos(\theta_i - \theta_j)} = 0$ and/or $r_{ij} = 0$, then the average power for a CW signal is equivalent to the average received power for a wideband signal in a small-scale region. This is seen by comparing equation (4.19) and equation (4.24). This can occur when either the multipath phases are identically and independently distributed (i.i.d uniform) over $[0, 2\pi]$ or when the path amplitudes are uncorrelated. The i.i.d uniform distribution of θ is a valid assumption since multipath components traverse differential path lengths that measure hundreds of wavelengths and are likely to arrive with random phases. If for some reason it is believed that the phases are not independent, the average wideband power and average CW power will still be equal if the paths have uncorrelated amplitudes. However, if the phases of the paths are dependent upon each other, then the amplitudes are likely to be correlated, since the same mechanism which affects the path phases is likely to also affect the amplitudes. This situation is highly unlikely at transmission frequencies used in wireless mobile systems.

Thus it is seen that the *received local ensemble average power of wideband and narrowband signals are equivalent*. When the transmitted signal has a bandwidth much greater than the bandwidth of the channel, then the multipath structure is completely resolved by the received signal at any time, and the received power varies very little since the individual multipath amplitudes do not change rapidly over a local area. However, if the transmitted signal has a very narrow bandwidth (e.g., the baseband signal has a duration greater than the excess delay of the channel), then multipath is not resolved by the received signal, and large signal fluctuations (fading) occur at the receiver due to the phase shifts of the many unresolved multipath components.

Figure 4.5 illustrates actual indoor radio channel measurements made simultaneously with a wideband probing pulse having $T_{bb} = 10$ ns, and a CW transmitter. The carrier frequency was 4 GHz. It can be seen that the CW signal undergoes rapid fades, whereas the wideband measurements change little over the 5λ measurement track. However, the local average received powers of both signals were measured to be virtually identical [Haw91].

Small-Scale Multipath Measurements

Because of the importance of the multipath structure in determining the small-scale fading effects, a number of wideband channel sounding techniques have been developed. These techniques may be classified as *direct pulse measurements*, *spread spectrum sliding correlator measurements*, and *swept frequency measurements*.

1. Direct RF Pulse System

A simple channel sounding approach is the direct RF pulse system (see Figure 4.6). This technique allows engineers to determine rapidly the power delay profile of any channel, as demonstrated by Rappaport and Seidel [Rap89], [Rap90]. Essentially a wide band pulsed bistatic radar, this system transmits a repetitive pulse of width τ_{bb} s, and uses a receiver with a wide bandpass filter ($BW = 2/\tau_{bb}$ Hz). The signal is then amplified, detected with an envelope detector, and displayed and stored on a high speed oscilloscope. This gives an immediate measurement of the square of the channel impulse response convolved with the probing pulse (see equation (4.17)). If the oscilloscope is set on averaging mode, then this system can provide a local average power delay profile. Another attractive aspect of this system is the lack of complexity, since off-the-shelf equipment may be used.

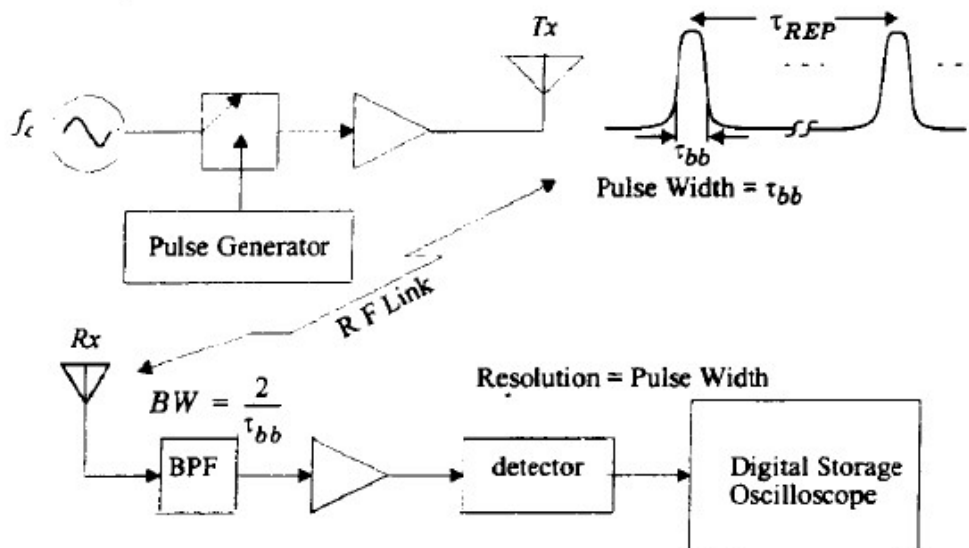


Figure 4.6
Direct RF channel impulse response measurement system.

The minimum resolvable delay between multipath components is equal to the probing pulse width τ_{bb} . The main problem with this system is that it is subject to interference and noise, due to the wide passband filter required for multipath time resolution. Also, the pulse system relies on the ability to trigger the oscilloscope on the first arriving signal. If the first arriving signal is blocked or fades, severe fading occurs, and it is possible the system may not trigger properly. Another disadvantage is that the phases of the individual multipath components are not received, due to the use of an envelope detector. However, use of a coherent detector permits measurement of the multipath phase using this technique.

2. Spread Spectrum Sliding Correlator Channel Sounding

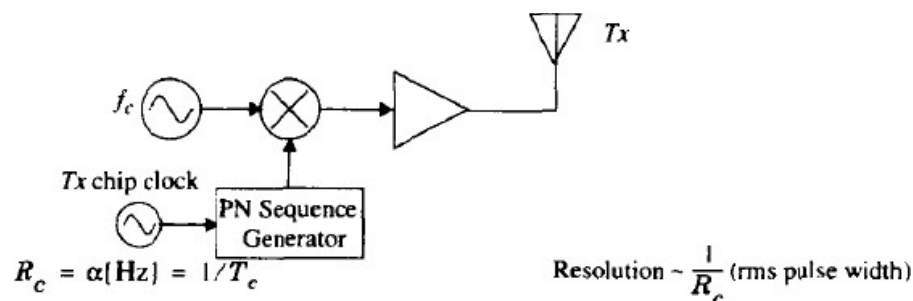
The basic block diagram of a spread spectrum channel sounding system is shown in Figure 4.7. The advantage of a spread spectrum system is that, while the probing signal may be wideband, it is possible to detect the transmitted signal using a narrowband receiver preceded by a wideband mixer, thus improving the dynamic range of the system as compared to the direct RF pulse system.

In a spread spectrum channel sounder, a carrier signal is "spread" over a large bandwidth by mixing it with a binary pseudo-noise (PN) sequence having a chip duration T_c and a chip rate R_c equal to $1/T_c$ Hz. The power spectrum envelope of the transmitted spread spectrum signal is given by [Dix84] as

$$S(f) = \left[\frac{\sin \pi(f - f_c)T_c}{\pi(f - f_c)T_c} \right]^2 \quad (4.26)$$

and the null-to-null bandwidth is

$$BW = 2R_c \quad (4.27)$$



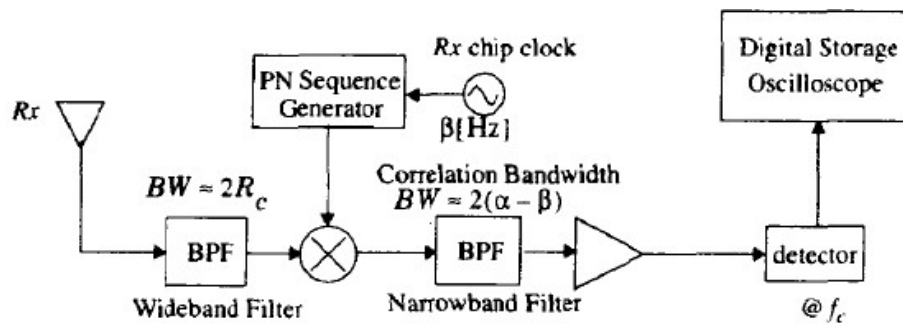


Figure 4.7
Spread spectrum channel impulse response measurement system.

The spread spectrum signal is then received, filtered, and *despread* using a PN sequence generator identical to that used at the transmitter. Although the two PN sequences are identical, the transmitter chip clock is run at a slightly faster rate than the receiver chip clock. Mixing the chip sequences in this fashion implements a *sliding correlator* [Dix84]. When the PN code of the faster chip clock catches up with the PN code of the slower chip clock, the two chip sequences will be virtually identically aligned, giving maximal correlation. When the two sequences are not maximally correlated, mixing the incoming spread spectrum signal with the unsynchronized receiver chip sequence will spread this signal into a bandwidth at least as large as the receiver's reference PN sequence. In this way, the narrowband filter that follows the correlator can reject almost all of the incoming signal power. This is how *processing gain* is realized in a spread spectrum receiver and how it can reject passband interference, unlike the direct RF pulse sounding system.

Processing gain (PG) is given as

$$PG = \frac{2R_c}{R_{bb}} = \frac{2\tau_{bb}}{T_c} = \frac{(S/N)_{out}}{(S/N)_{in}} \quad (4.28)$$

where $\tau_{bb} = 1/R_{bb}$, is the period of the baseband information. For the case of a sliding correlator channel sounder, the baseband information rate is equal to the frequency offset of the PN sequence clocks at the transmitter and receiver.

When the incoming signal is correlated with the receiver sequence, the signal is collapsed back to the original bandwidth (i.e., "despread"), envelope detected, and displayed on an oscilloscope. Since different incoming multipaths will have different time delays, they will maximally correlate with the receiver PN sequence at different times. The energy of these individual paths will pass through the correlator depending on the time delay. Therefore, after envelope detection, the channel impulse response convolved with the pulse shape of a single chip is displayed on the oscilloscope. Cox [Cox72] first used this method to measure channel impulse responses in outdoor suburban environments at 910 MHz. Devasirvatham [Dev86], [Dev90a] successfully used a direct sequence spread spectrum channel sounder to measure time delay spread of multipath components and signal level measurements in office and residential buildings at 850 MHz. Bultitude [Bul89] used this technique for indoor and microcellular channel sounding work, as did Landron [Lan92].

The time resolution ($\Delta\tau$) of multipath components using a spread spectrum system with sliding correlation is

$$\Delta\tau = 2T_c = \frac{2}{R_c} \quad (4.29)$$

In other words, the system can resolve two multipath components as long as they are equal to or greater than $2T_c$ seconds apart. In actuality, multipath components with interarrival times smaller than $2T_c$ can be resolved since the rms pulse width is smaller than the absolute width of the triangular correlation pulse, and is on the order of T_c .

The sliding correlation process gives *equivalent time* measurements that are updated every time the two sequences are maximally correlated. The time between maximal correlations (T) can be calculated from equation (4.30)

$$\Delta T = T_c \gamma l = \frac{\gamma l}{R_c} \quad (4.30)$$

where T_c = chip period (s)
 R_c = chip rate (Hz)
 γ = slide factor (dimensionless)
 l = sequence length (chips)

The slide factor is defined as the ratio between the transmitter chip clock rate and the difference between the transmitter and receiver chip clock rates [Dev86]. Mathematically, this is expressed as

$$\gamma = \frac{\alpha}{\alpha - \beta} \quad (4.31)$$

where α = transmitter chip clock rate (Hz)
 β = receiver chip clock rate (Hz)

For a maximal length PN sequence, the sequence length is

$$l = 2^n - 1 \quad (4.32)$$

where n is the number of shift registers in the sequence generator [Dix84].

Since the incoming spread spectrum signal is mixed with a receiver PN sequence that is slower than the transmitter sequence, the signal is essentially down-converted ("collapsed") to a low-frequency narrowband signal. In other words, the relative rate of the two codes slipping past each other is the rate of information transferred to the oscilloscope. This narrowband signal allows narrowband processing, eliminating much of the passband noise and interference. The processing gain of equation (4.28) is then realized using a narrowband filter ($BW = 2(\alpha - \beta)$).

The equivalent time measurements refer to the relative times of multipath components as they are displayed on the oscilloscope. The observed time scale on the oscilloscope using a sliding correlator is related to the actual propagation time scale by

$$\text{Actual Propagation Time} = \frac{\text{Observed Time}}{\gamma} \quad (4.33)$$

This effect is due to the relative rate of information transfer in the sliding correlator. For example, T_c of equation (4.30) is an observed time measured on an oscilloscope and not actual propagation time. This effect, known as *time dilation*, occurs in the sliding correlator system because the propagation delays are actually expanded in time by the sliding correlator.

Caution must be taken to ensure that the sequence length has a period which is greater than the longest multipath propagation delay. The PN sequence period is

$$\tau_{PNseq} = T_c l \quad (4.34)$$

The sequence period gives an estimate of the maximum unambiguous range of incoming multipath signal components. This range is found by multiplying the speed of light with τ_{PNseq} in equation (4.34).

There are several advantages to the spread spectrum channel sounding system. One of the key spread spectrum modulation characteristics is the ability to reject passband noise, thus improving the coverage range for a given transmitter power. Transmitter and receiver PN sequence synchronization is eliminated by the sliding correlator. Sensitivity is adjustable by changing the sliding factor and the post-correlator filter bandwidth. Also, required transmitter powers can be considerably lower than comparable direct pulse systems due to the inherent "processing gain" of spread spectrum systems.

A disadvantage of the spread spectrum system, as compared to the direct pulse system, is that measurements are not made in real time, but they are compiled as the PN codes slide past one another. Depending on system parameters and measurement objectives, the time required to make power delay profile measurements may be excessive. Another disadvantage of the system described here is that a noncoherent detector is used, so that phases of individual multipath components can not be measured. Even if coherent detection is used, the sweep time of a spread spectrum signal induces delay such that the phases of individual multipath components with different time delays would be measured at substantially different times, during which the channel might change.

3. Frequency Domain Channel Sounding

Because of the dual relationship between time domain and frequency domain techniques, it is possible to measure the channel impulse response in the frequency domain. Figure 4.8 shows a frequency domain channel sounder used for measuring channel impulse responses. A vector network analyzer controls a synthesized frequency sweeper, and an S-parameter test set is used to monitor the frequency response of the channel. The sweeper scans a particular frequency band (centered on the carrier) by stepping through discrete frequencies. The number and spacings of these frequency steps impact the time resolution of the impulse response measurement. For each frequency step, the S-parameter test set transmits a known signal level at port 1 and monitors the received signal level at port 2. These signal levels allow the analyzer to determine the complex response (i.e., transmissivity $S_{21}(\omega)$) of the channel over the measured frequency range. The transmissivity response is a frequency domain representation of the channel impulse response. This response is then converted to the time domain using inverse discrete Fourier transform (IDFT) processing, giving a band-limited version of the impulse response. In theory, this technique works well and indirectly provides amplitude and phase information in the time domain. However, the system requires careful calibration and hardwired synchronization between the transmitter and receiver, making it useful only for very close measurements (e.g., indoor channel sounding). Another limitation with this system is the non-real-time nature of the measurement. For time varying channels, the channel frequency response can change rapidly, giving an erroneous impulse response measurement. To mitigate this effect, fast sweep times are necessary to keep the total swept frequency response measurement interval as short as possible. A faster sweep time can be accomplished by reducing the number of frequency steps, but this sacrifices time resolution and excess delay range in the time domain. The swept frequency system has been used successfully for indoor propagation studies by Pahlavan [Pah95] and Zaghloul, et.al. [Zag91a], [Zag91b].

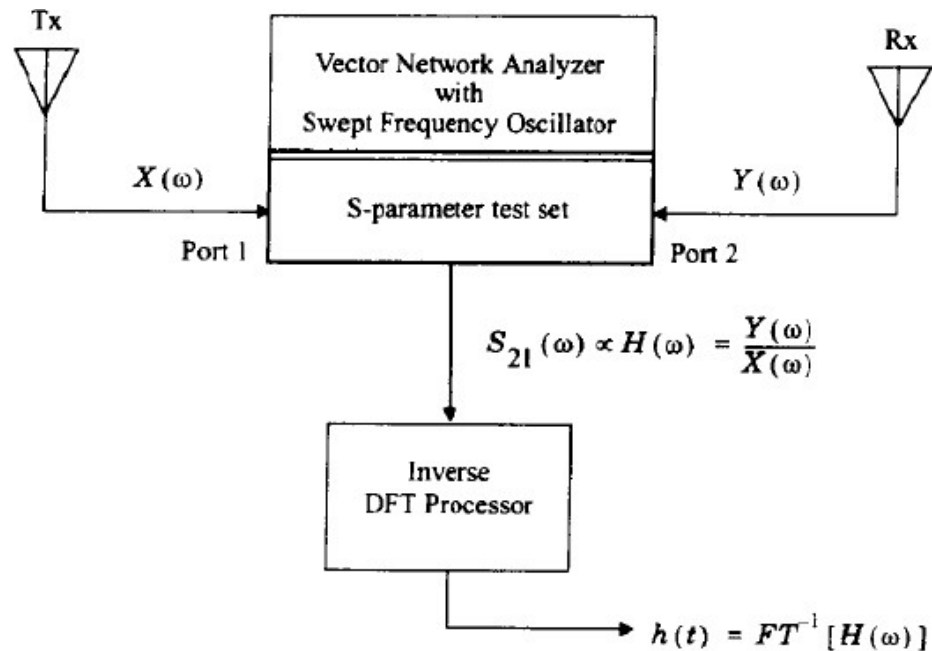
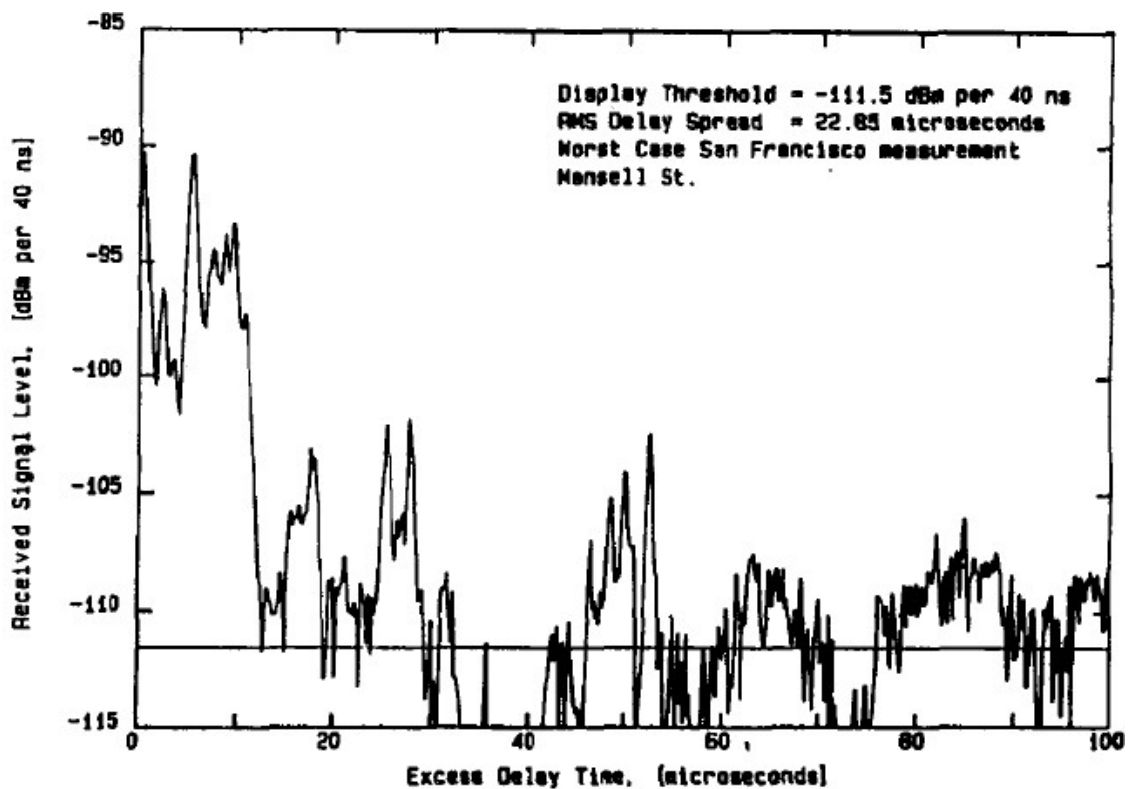


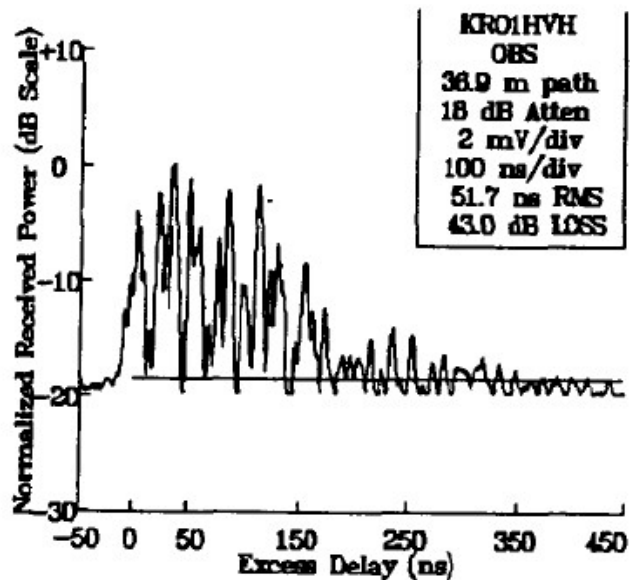
Figure 4.8
Frequency domain channel impulse response measurement system.

Parameters of Mobile Multipath Channels

Many multipath channel parameters are derived from the power delay profile, given by equation (4.18). Power delay profiles are measured using the techniques discussed in Section 4.4 and are generally represented as plots of relative received power as a function of excess delay with respect to a fixed time delay reference. Power delay profiles are found by averaging instantaneous power delay profile measurements over a local area in order to determine an average small-scale power delay profile. Depending on the time resolution of the probing pulse and the type of multipath channels studied, researchers often choose to sample at spatial separations of a quarter of a wavelength and over receiver movements no greater than 6 m in outdoor channels and no greater than 2 m in indoor channels in the 450 MHz - 6 GHz range. This small-scale sampling avoids large-scale averaging bias in the resulting small-scale statistics. Figure 4.9 shows typical power delay profile plots from outdoor and indoor channels, determined from a large number of closely sampled instantaneous profiles.



(a)



(b)

Figure 4.9

Measured multipath power delay profiles

a) From a 900 MHz cellular system in San Francisco [From [Rap90] © IEEE].

b) Inside a grocery store at 4 GHz [From [Haw91] © IEEE].

Time Dispersion Parameters

In order to compare different multipath channels and to develop some general design guidelines for wireless systems, parameters which grossly quantify the multipath channel are used. The *mean excess delay*, *rms delay spread*, and *excess delay spread* (X dB) are multipath channel parameters that can be determined from a power delay profile. The time dispersive properties of wide band multipath channels are most commonly quantified by their mean excess delay ($\bar{\tau}$) and rms delay spread (σ_{τ}). The mean excess delay is the first moment of the power delay profile and is defined to be

$$\bar{\tau} = \frac{\sum_k a_k^2 \tau_k}{\sum_k a_k^2} = \frac{\sum_k P(\tau_k) \tau_k}{\sum_k P(\tau_k)} \quad (4.35)$$

The rms delay spread is the square root of the second central moment of the power delay profile and is defined to be

$$\sigma_{\tau} = \sqrt{\overline{\tau^2} - (\bar{\tau})^2} \quad (4.36)$$

where

$$\overline{\tau^2} = \frac{\sum_k a_k^2 \tau_k^2}{\sum_k a_k^2} = \frac{\sum_k P(\tau_k) \tau_k^2}{\sum_k P(\tau_k)} \quad (4.37)$$

These delays are measured relative to the first detectable signal arriving at the receiver at $\tau_0 = 0$. Equations (4.35) - (4.37) do not rely on the absolute power level of $P(\tau)$, but only the relative amplitudes of the multipath components within $P(\tau)$. Typical values of rms delay spread are on the order of microseconds in outdoor mobile radio channels and on the order of nanoseconds in indoor radio channels. Table 4.1 shows the typical measured values of rms delay spread.

It is important to note that the rms delay spread and mean excess delay are defined from a single power delay profile which is the temporal or spatial average of consecutive impulse response measurements collected and averaged over a local area. Typically, many measurements are made at many local areas in order to determine a statistical range of multipath channel parameters for a mobile communication system over a large-scale area [Rap90].

The *maximum excess delay* (X dB) of the power delay profile is defined to be the time delay during which multipath energy falls to X dB below the maxi-

Table 4.1 Typical Measured Values of RMS Delay Spread

Environment	Frequency (MHz)	RMS Delay Spread (σ_τ)	Notes	Reference
Urban	910	1300 ns avg. 600 ns st. dev. 3500 ns max.	New York City	[Cox75]
Urban	892	10-25 μ s	Worst case San Francisco	[Rap90]
Suburban	910	200-310 ns	Averaged typical case	[Cox72]
Suburban	910	1960-2110 ns	Averaged extreme case	[Cox72]
Indoor	1500	10-50 ns 25 ns median	Office building	[Sal87]
Indoor	850	270 ns max.	Office building	[Dev90a]
Indoor	1900	70-94 ns avg. 1470 ns max.	Three San Francisco buildings	[Sei92a]

mum. In other words, the maximum excess delay is defined as $\tau_X - \tau_0$, where τ_0 is the first arriving signal and τ_X is the maximum delay at which a multipath component is within X dB of the strongest arriving multipath signal (which does not necessarily arrive at τ_0). Figure 4.10 illustrates the computation of the maximum excess delay for multipath components within 10 dB of the maximum. The maximum excess delay (X dB) defines the temporal extent of the multipath that is above a particular threshold. The value of τ_X is sometimes called the *excess delay spread* of a power delay profile, but in all cases must be specified with a threshold that relates the multipath noise floor to the maximum received multipath component.

In practice, values for $\bar{\tau}$, $\bar{\tau}^2$, and σ_τ depend on the choice of noise threshold used to process $P(\tau)$. The noise threshold is used to differentiate between received multipath components and thermal noise. If the noise threshold is set too low, then noise will be processed as multipath, thus giving rise to values of $\bar{\tau}$, $\bar{\tau}^2$, and σ_τ that are artificially high.

It should be noted that the power delay profile and the magnitude frequency response (the spectral response) of a mobile radio channel are related through the Fourier transform. It is therefore possible to obtain an equivalent description of the channel in the frequency domain using its frequency response characteristics. Analogous to the delay spread parameters in the time domain, *coherence bandwidth* is used to characterize the channel in the frequency domain. The rms delay spread and coherence bandwidth are inversely proportional to one another, although their exact relationship is a function of the exact multipath structure.

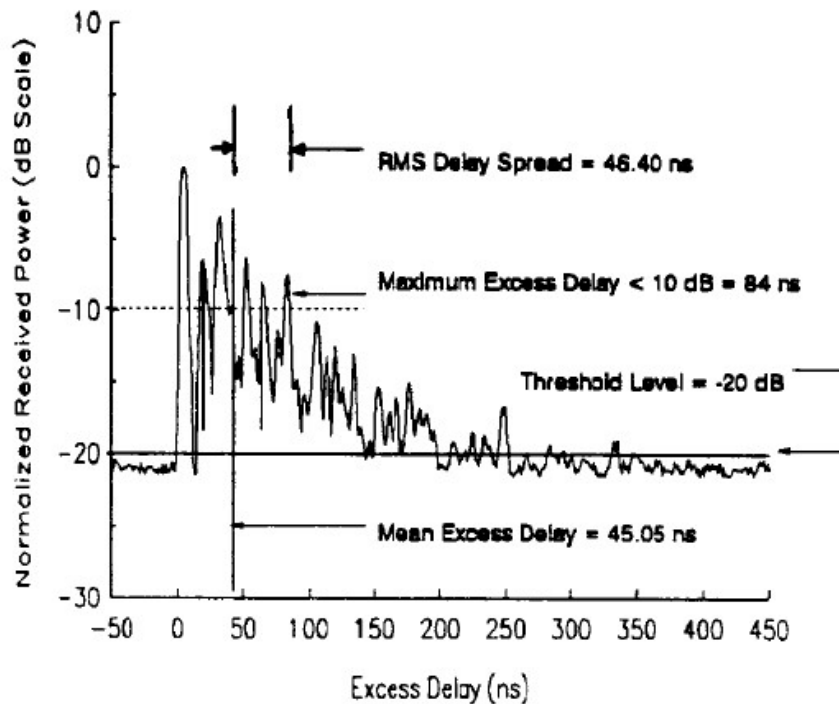


Figure 4.10
 Example of an indoor power delay profile; rms delay spread, mean excess delay, maximum excess delay (10 dB), and threshold level are shown.

Coherence Bandwidth

While the delay spread is a natural phenomenon caused by reflected and scattered propagation paths in the radio channel, the coherence bandwidth, B_c , is a defined relation derived from the rms delay spread. Coherence bandwidth is a statistical measure of the range of frequencies over which the channel can be considered “flat” (i.e., a channel which passes all spectral components with approximately equal gain and linear phase). In other words, coherence bandwidth is the range of frequencies over which two frequency components have a strong potential for amplitude correlation. Two sinusoids with frequency separation greater than B_c are affected quite differently by the channel. If the coherence bandwidth is defined as the bandwidth over which the frequency correlation function is above 0.9, then the coherence bandwidth is approximately [Lee89b]

$$B_c \approx \frac{1}{50\sigma_\tau} \quad (4.38)$$

If the definition is relaxed so that the frequency correlation function is above 0.5, then the coherence bandwidth is approximately

$$B_c \approx \frac{1}{5\sigma_\tau} \quad (4.39)$$

It is important to note that an exact relationship between coherence bandwidth and rms delay spread does not exist, and equations (4.38) and (4.39) are "ball park estimates". In general, spectral analysis techniques and simulation are required to determine the exact impact that time varying multipath has on a particular transmitted signal [Chu87], [Fun93], [Ste94]. For this reason, accurate multipath channel models must be used in the design of specific modems for wireless applications [Rap91a], [Woe94].

Doppler Spread and Coherence Time

Delay spread and coherence bandwidth are parameters which describe the time dispersive nature of the channel in a local area. However, they do not offer information about the time varying nature of the channel caused by either relative motion between the mobile and base station, or by movement of objects in the channel. *Doppler spread* and *coherence time* are parameters which describe the time varying nature of the channel in a small-scale region.

Doppler spread B_D is a measure of the spectral broadening caused by the time rate of change of the mobile radio channel and is defined as the range of frequencies over which the received Doppler spectrum is essentially non-zero. When a pure sinusoidal tone of frequency f_c is transmitted, the received signal spectrum, called the Doppler spectrum, will have components in the range $f_c - f_d$ to $f_c + f_d$, where f_d is the Doppler shift. The amount of spectral broadening depends on f_d which is a function of the relative velocity of the mobile, and the angle θ between the direction of motion of the mobile and direction of arrival of the scattered waves. *If the baseband signal bandwidth is much greater than B_D , the effects of Doppler spread are negligible at the receiver. This is a slow fading channel.*

Coherence time T_c is the time domain dual of Doppler spread and is used to characterize the time varying nature of the frequency dispersiveness of the channel in the time domain. The Doppler spread and coherence time are inversely proportional to one another. That is,

$$T_c \approx \frac{1}{f_m} \quad (4.40.a)$$

Coherence time is actually a statistical measure of the time duration over which the channel impulse response is essentially invariant, and quantifies the similarity of the channel response at different times. In other words, coherence time is the time duration over which two received signals have a strong potential for amplitude correlation. If the reciprocal bandwidth of the baseband signal is greater than the coherence time of the channel, then the channel will change during the transmission of the baseband message, thus causing distortion at the receiver. If the coherence time is defined as the time over which the time correlation function is above 0.5, then the coherence time is approximately [Ste94]

$$T_c \approx \frac{9}{16\pi f_m} \quad (4.40.b)$$

where f_m is the maximum Doppler shift given by $f_m = v/\lambda$. In practice, (4.40.a) suggests a time duration during which a Rayleigh fading signal may fluctuate

wildly, and (4.40.b) is often too restrictive. A popular rule of thumb for modern digital communications is to define the coherence time as the geometric mean of equations (4.40.a) and (4.40.b). That is,

$$T_C = \sqrt{\frac{9}{16\pi f_m^2}} = \frac{0.423}{f_m} \quad (4.40.c)$$

The definition of coherence time implies that two signals arriving with a time separation greater than T_C are affected differently by the channel. For example, for a vehicle traveling 60 mph using a 900 MHz carrier, a conservative value of T_C can be shown to be 2.22 ms from (4.40.b). If a digital transmission system is used, then as long as the symbol rate is greater than $1/T_C = 454$ bps, the channel will not cause distortion due to motion (however, distortion could result from multipath time delay spread, depending on the channel impulse response). Using the practical formula of (4.40.c), $T_C = 6.77$ ms and the symbol rate must exceed 150 bits/s in order to avoid distortion due to frequency dispersion.

Example 4.5

Determine the proper spatial sampling interval required to make small-scale propagation measurements which assume that consecutive samples are highly correlated in time. How many samples will be required over 10 m travel distance if $f_c = 1900$ MHz and $v = 50$ m/s. How long would it take to make these measurements, assuming they could be made in real time from a moving vehicle? What is the Doppler spread B_D for the channel?

Solution to Example 4.5

For correlation, ensure that the time between samples is equal to $T_C/2$, and use the smallest value of T_C for conservative design.

Using equation (4.40.b)

$$T_C \approx \frac{9}{16\pi f_m} = \frac{9\lambda}{16\pi v} = \frac{9c}{16\pi v f_c} = \frac{9 \times 3 \times 10^8}{16 \times 3.14 \times 50 \times 1900 \times 10^6}$$

$$T_C = 565 \mu\text{s}$$

Taking time samples at less than half T_C , at $282.5 \mu\text{s}$ corresponds to a spatial sampling interval of

$$\Delta x = \frac{v T_C}{2} = \frac{50 \times 565 \mu\text{s}}{2} = 0.014125 \text{ m} = 1.41 \text{ cm}$$

Therefore, the number of samples required over a 10 m travel distance is

$$N_x = \frac{10}{\Delta x} = \frac{10}{0.014125} = 708 \text{ samples}$$

The time taken to make this measurement is equal to $\frac{10 \text{ m}}{50 \text{ m/s}} = 0.2 \text{ s}$

The Doppler spread is

$$B_D = f_m = \frac{v f_c}{c} = \frac{50 \times 1900 \times 10^6}{3 \times 10^8} = 316.66 \text{ Hz}$$

Types of Small-Scale Fading

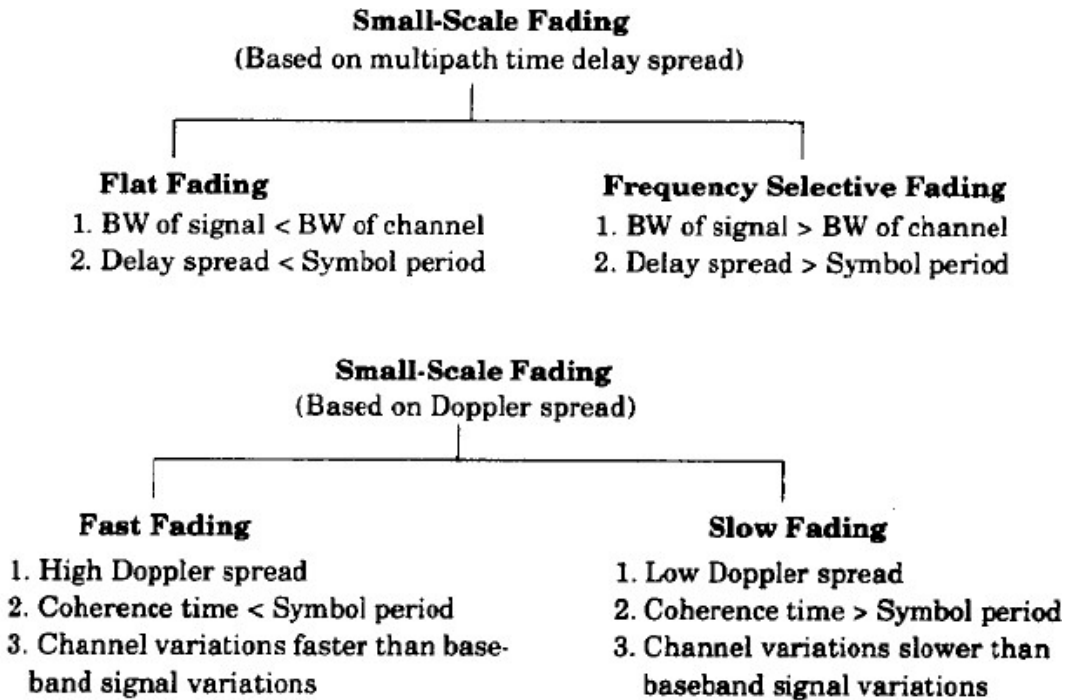


Figure 4.11
Types of small-scale fading.

Fading Effects Due to Multipath Time Delay Spread

Time dispersion due to multipath causes the transmitted signal to undergo either flat or frequency selective fading.

1. Flat fading

If the mobile radio channel has a constant gain and linear phase response over a bandwidth which is greater than the bandwidth of the transmitted signal, then the received signal will undergo *flat fading*. This type of fading is historically the most common type of fading described in the technical literature. In flat fading, the multipath structure of the channel is such that the spectral characteristics of the transmitted signal are preserved at the receiver. However the strength of the received signal changes with time, due to fluctuations in the gain of the channel caused by multipath. The characteristics of a flat fading channel are illustrated in Figure 4.12.

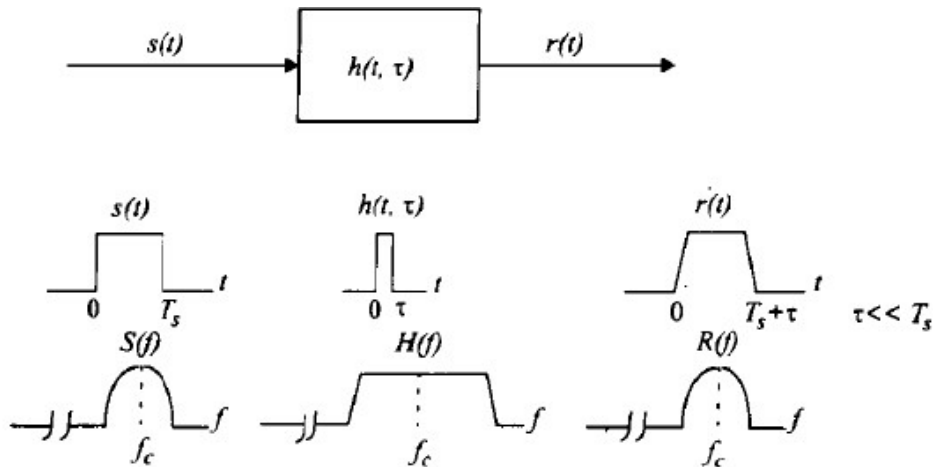


Figure 4.12
Flat fading channel characteristics.

It can be seen from Figure 4.12 that if the channel gain changes over time, a change of amplitude occurs in the received signal. Over time, the received signal $r(t)$ varies in gain, but the spectrum of the transmission is preserved. In a flat fading channel, the reciprocal bandwidth of the transmitted signal is much larger than the multipath time delay spread of the channel, and $h_b(t, \tau)$ can be approximated as having no excess delay (i.e., a single delta function with $\tau = 0$). Flat fading channels are also known as *amplitude varying channels* and are sometimes referred to as *narrowband channels*, since the bandwidth of the applied signal is *narrow* as compared to the channel flat fading bandwidth. Typical flat fading channels cause deep fades, and thus may require 20 or 30 dB more transmitter power to achieve low bit error rates during times of deep fades as compared to systems operating over non-fading channels. The distribution of the instantaneous gain of flat fading channels is important for designing radio links, and the most common amplitude distribution is the Rayleigh distribution. The Rayleigh flat fading channel model assumes that the channel induces an amplitude which varies in time according to the Rayleigh distribution.

To summarize, a signal undergoes flat fading if

$$B_S \ll B_C \quad (4.41)$$

and

$$T_S \gg \sigma_\tau \quad (4.42)$$

where T_S is the reciprocal bandwidth (e.g., symbol period) and B_S is the bandwidth, respectively, of the transmitted modulation, and σ_τ and B_C are the rms delay spread and coherence bandwidth, respectively, of the channel.

2. Frequency Selective Fading

If the channel possesses a constant-gain and linear phase response over a bandwidth that is smaller than the bandwidth of transmitted signal, then the channel creates *frequency selective fading* on the received signal. Under such conditions the channel impulse response has a multipath delay spread which is greater than the reciprocal bandwidth of the transmitted message waveform. When this occurs, the received signal includes multiple versions of the transmitted waveform which are attenuated (faded) and delayed in time, and hence the received signal is distorted. Frequency selective fading is due to time dispersion of the transmitted symbols within the channel. Thus the channel induces *intersymbol interference* (ISI). Viewed in the frequency domain, certain frequency components in the received signal spectrum have greater gains than others.

Frequency selective fading channels are much more difficult to model than flat fading channels since each multipath signal must be modeled and the channel must be considered to be a linear filter. It is for this reason that wideband multipath measurements are made, and models are developed from these measurements. When analyzing mobile communication systems, statistical impulse response models such as the 2-ray Rayleigh fading model (which considers the impulse response to be made up of two delta functions which independently fade and have sufficient time delay between them to induce frequency selective fading upon the applied signal), or computer generated or measured impulse responses, are generally used for analyzing frequency selective small-scale fading. Figure 4.13 illustrates the characteristics of a frequency selective fading channel.

For frequency selective fading, the spectrum $S(f)$ of the transmitted signal has a bandwidth which is greater than the coherence bandwidth B_C of the channel. Viewed in the frequency domain, the channel becomes frequency selective, where the gain is different for different frequency components. Frequency selec-

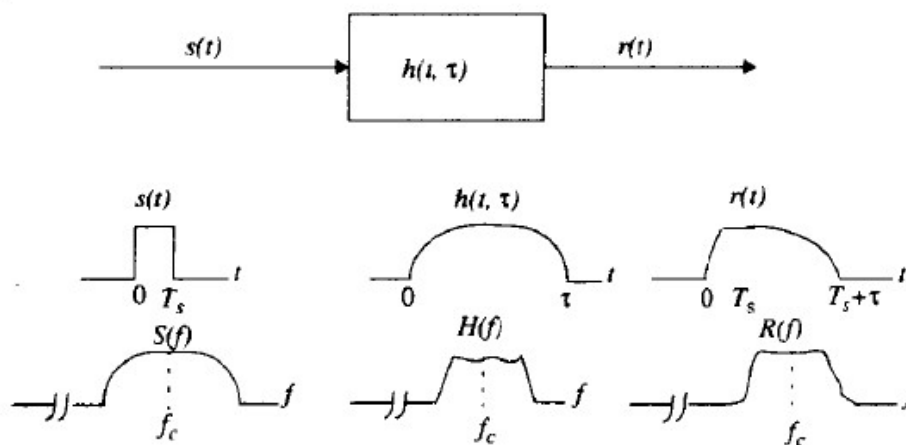


Figure 4.13
Frequency selective fading channel characteristics.

tive fading is caused by multipath delays which approach or exceed the symbol period of the transmitted symbol. Frequency selective fading channels are also known as *wideband channels* since the bandwidth of the signal $s(t)$ is wider than the bandwidth of the channel impulse response. As time varies, the channel varies in gain and phase across the spectrum of $s(t)$, resulting in time varying distortion in the received signal $r(t)$. To summarize, a signal undergoes frequency selective fading if

$$B_S > B_C \quad (4.43)$$

and

$$T_S < \sigma_\tau \quad (4.44)$$

A common rule of thumb is that a channel is frequency selective if $T_S \leq 10\sigma_\tau$, although this is dependent on the specific type of modulation used. Chapter 5 presents simulation results which illustrate the impact of time delay spread on bit error rate (BER).

Fading Effects Due to Doppler Spread

1. Fast Fading

Depending on how rapidly the transmitted baseband signal changes as compared to the rate of change of the channel, a channel may be classified either as a *fast fading* or *slow fading* channel. In a *fast fading channel*, the channel impulse response changes rapidly within the symbol duration. That is, the coherence time of the channel is smaller than the symbol period of the transmitted signal. This causes frequency dispersion (also called time selective fading) due to Doppler spreading, which leads to signal distortion. Viewed in the frequency domain, signal distortion due to fast fading increases with increasing Doppler spread relative to the bandwidth of the transmitted signal. Therefore, a signal undergoes fast fading if

$$T_S > T_C \quad (4.45)$$

and

$$B_S < B_D \quad (4.46)$$

It should be noted that when a channel is specified as a fast or slow fading channel, it does not specify whether the channel is flat fading or frequency selective in nature. Fast fading only deals with the rate of change of the channel due to motion. In the case of the flat fading channel, we can approximate the impulse response to be simply a delta function (no time delay). Hence, a *flat fading, fast fading* channel is a channel in which the amplitude of the delta function varies faster than the rate of change of the transmitted baseband signal. In the case of a *frequency selective, fast fading* channel, the amplitudes, phases, and time delays of any one of the multipath components vary faster than the rate of change of the transmitted signal. In practice, fast fading only occurs for very low data rates.

2. Slow Fading

In a *slow fading channel*, the channel impulse response changes at a rate much slower than the transmitted baseband signal $s(t)$. In this case, the channel may be assumed to be static over one or several reciprocal bandwidth intervals. In the frequency domain, this implies that the Doppler spread of the channel is much less than the bandwidth of the baseband signal. Therefore, a signal undergoes slow fading if

$$T_S \ll T_C \quad (4.47)$$

and

$$B_S \gg B_D \quad (4.48)$$

It should be clear that the velocity of the mobile (or velocity of objects in the channel) and the baseband signaling determines whether a signal undergoes fast fading or slow fading.

The relation between the various multipath parameters and the type of fading experienced by the signal are summarized in Figure 4.14. Over the years, some authors have confused the terms fast and slow fading with the terms large-scale and small-scale fading. It should be emphasized that fast and slow fading deal with the relationship between the time rate of change in the channel and the transmitted signal, and not with propagation path loss models.

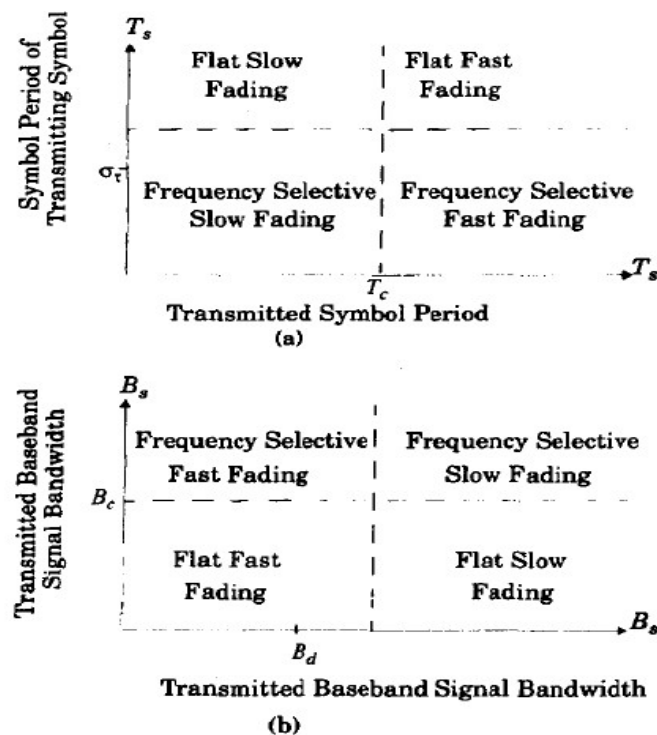


Figure 4.14
Matrix illustrating type of fading experienced by a signal as a function of
(a) symbol period
(b) baseband signal bandwidth.

Statistical Models for Multipath Fading Channels:

Several multipath models have been suggested to explain the observed statistical nature of a mobile channel. The first model presented by Ossana [Oss64] was based on interference of waves incident and reflected from the flat sides of randomly located buildings. Although Ossana's model [Oss64] predicts flat fading power spectra that were in agreement with measurements in suburban areas, it assumes the existence of a direct path between the transmitter and receiver, and is limited to a restricted range of reflection angles. Ossana's model is therefore rather inflexible and inappropriate for urban areas where the direct path is almost always blocked by buildings or other obstacles. Clarke's model [Cla68] is based on scattering and is widely used.

.1 Clarke's Model for Flat Fading

Clarke [Cla68] developed a model where the statistical characteristics of the electromagnetic fields of the received signal at the mobile are deduced from scattering. The model assumes a fixed transmitter with a vertically polarized antenna. The field incident on the mobile antenna is assumed to be comprised of N azimuthal plane waves with arbitrary carrier phases, arbitrary azimuthal angles of arrival, and each wave having equal average amplitude. It should be noted that the equal average amplitude assumption is based on the fact that in the absence of a direct line-of-sight path, the scattered components arriving at a receiver will experience similar attenuation over small-scale distances.

Figure 4.19 shows a diagram of plane waves incident on a mobile traveling at a velocity v , in the x -direction. The angle of arrival is measured in the x - y plane with respect to the direction of motion. Every wave that is incident on the mobile undergoes a Doppler shift due to the motion of the receiver and arrives at the receiver at the same time. That is, no excess delay due to multipath is assumed for any of the waves (flat fading assumption). For the n th wave arriving at an angle α_n to the x -axis, the Doppler shift in Hertz is given by

$$f_n = \frac{v}{\lambda} \cos \alpha_n \quad (4.57)$$

where λ is the wavelength of the incident wave.

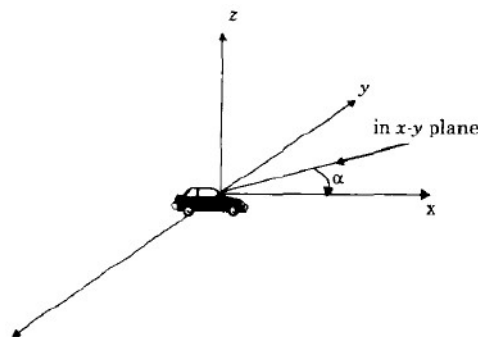


Figure 4.19
Illustrating plane waves arriving at random angles.

The vertically polarized plane waves arriving at the mobile have E and H field components given by

$$E_z = E_0 \sum_{n=1}^N C_n \cos(2\pi f_c t + \theta_n) \quad (4.58)$$

$$H_x = -\frac{E_0}{\eta} \sum_{n=1}^N C_n \sin \alpha_n \cos(2\pi f_c t + \theta_n) \quad (4.59)$$

$$H_y = -\frac{E_0}{\eta} \sum_{n=1}^N C_n \cos \alpha_n \cos(2\pi f_c t + \theta_n) \quad (4.60)$$

where E_0 is the real amplitude of local average E-field (assumed constant), C_n is a real random variable representing the amplitude of individual waves, η is the intrinsic impedance of free space (377Ω), and f_c is the carrier frequency. The random phase of the n th arriving component θ_n is given by

$$\theta_n = 2\pi f_n t + \phi_n \quad (4.61)$$

The amplitudes of the E-and H-field are normalized such that the ensemble average of the C_n 's is given by

$$\sum_{n=1}^N \overline{C_n^2} = 1 \quad (4.62)$$

Since the Doppler shift is very small when compared to the carrier frequency, the three field components may be modeled as narrow band random processes. The three components E_z , H_x , and H_y can be approximated as Gaussian random variables if N is sufficiently large. The phase angles are assumed to have a uniform probability density function (pdf) on the interval $(0, 2\pi]$. Based on the analysis by Rice [Ric48] the E-field can be expressed in an in-phase and quadrature form

$$E_z = T_c(t) \cos(2\pi f_c t) - T_s(t) \sin(2\pi f_c t) \quad (4.63)$$

where

$$T_c(t) = E_0 \sum_{n=1}^N C_n \cos(2\pi f_n t + \phi_n) \quad (4.64)$$

and

$$T_s(t) = E_0 \sum_{n=1}^N C_n \sin(2\pi f_n t + \phi_n) \quad (4.65)$$

Both $T_c(t)$ and $T_s(t)$ are Gaussian random processes which are denoted as T_c and T_s , respectively, at any time t . T_c and T_s are uncorrelated zero-mean Gaussian random variables with an equal variance given by

$$\overline{T_c^2} = \overline{T_s^2} = \overline{|E_z|^2} = E_0^2/2 \quad (4.66)$$

where the overbar denotes the ensemble average.

The envelope of the received E-field, $E_z(t)$, is given by

$$|E_z(t)| = \sqrt{T_c^2(t) + T_s^2(t)} = r(t) \quad (4.67)$$

Since T_c and T_s are Gaussian random variables, it can be shown through a Jacobean transformation [Pap91] that the random received signal envelope r has a Rayleigh distribution given by

$$p(r) = \begin{cases} \frac{r}{\sigma^2} \exp\left(-\frac{r^2}{2\sigma^2}\right) & 0 \leq r \leq \infty \\ 0 & r < 0 \end{cases} \quad (4.68)$$

where $\sigma^2 = E_0^2/2$

1.1 Spectral Shape Due to Doppler Spread in Clarke's Model

Gans [Gan72] developed a spectrum analysis for Clarke's model. Let $p(\alpha)d\alpha$ denote the fraction of the total incoming power within $d\alpha$ of the angle α , and let A denote the average received power with respect to an isotropic antenna. As $N \rightarrow \infty$, $p(\alpha)d\alpha$ approaches a continuous, rather than a discrete, distribution. If $G(\alpha)$ is the azimuthal gain pattern of the mobile antenna as a function of the angle of arrival, the total received power can be expressed as

$$P_r = \int_0^{2\pi} AG(\alpha)p(\alpha)d\alpha \quad (4.69)$$

where $AG(\alpha)p(\alpha)d\alpha$ is the differential variation of received power with angle. If the scattered signal is a CW signal of frequency f_c , then the instantaneous frequency of the received signal component arriving at an angle α is obtained using equation (4.57)

$$f(\alpha) = f = \frac{v}{\lambda} \cos(\alpha) + f_c = f_m \cos \alpha + f_c \quad (4.70)$$

where f_m is the maximum Doppler shift. It should be noted that $f(\alpha)$ is an even function of α , (i.e., $f(\alpha) = f(-\alpha)$).

If $S(f)$ is the power spectrum of the received signal, the differential variation of received power with frequency is given by

$$S(f)|df \quad (4.71)$$

Equating the differential variation of received power with frequency to the differential variation in received power with angle, we have

$$S(f)|df| = A[p(\alpha)G(\alpha) + p(-\alpha)G(-\alpha)]|d\alpha| \quad (4.72)$$

Differentiating equation (4.70), and rearranging the terms, we have

$$|df| = |d\alpha| \sin \alpha f_m \quad (4.73)$$

Using equation (4.70), α can be expressed as a function of f as

$$\alpha = \cos^{-1} \left[\frac{f-f_c}{f_m} \right] \quad (4.74)$$

This implies that

$$\sin \alpha = \sqrt{1 - \left(\frac{f-f_c}{f_m} \right)^2} \quad (4.75)$$

Substituting equation (4.73) and (4.75) into both sides of (4.72), the power spectral density $S(f)$ can be expressed as

$$S(f) = \frac{A[p(\alpha)G(\alpha) + p(-\alpha)G(-\alpha)]}{f_m \sqrt{1 - \left(\frac{f-f_c}{f_m} \right)^2}} \quad (4.76)$$

where

$$S(f) = 0, \quad |f-f_c| > f_m \quad (4.77)$$

The spectrum is centered on the carrier frequency and is zero outside the limits of $f_c \pm f_m$. Each of the arriving waves has its own carrier frequency (due to its direction of arrival) which is slightly offset from the center frequency. For the case of a vertical $\lambda/4$ antenna ($G(\alpha) = 1.5$), and a uniform distribution $p(\alpha) = 1/2\pi$ over 0 to 2π , the output spectrum is given by (4.76) as

$$S_{E_t}(f) = \frac{1.5}{\pi f_m \sqrt{1 - \left(\frac{f-f_c}{f_m} \right)^2}} \quad (4.78)$$

In equation (4.78) the power spectral density at $f = f_c \pm f_m$ is infinite, i.e., Doppler components arriving at exactly 0° and 180° have an infinite power spectral density. This is not a problem since α is continuously distributed and the probability of components arriving at exactly these angles is zero.

Figure 4.20 shows the power spectral density of the resulting RF signal due to Doppler fading. Smith [Smi75] demonstrated an easy way to simulate Clarke's model using a computer simulation as described Section 4.7.2.

After envelope detection of the Doppler-shifted signal, the resulting baseband spectrum has a maximum frequency of $2f_m$. It can be shown [Jak74] that the electric field produces a baseband power spectral density given by

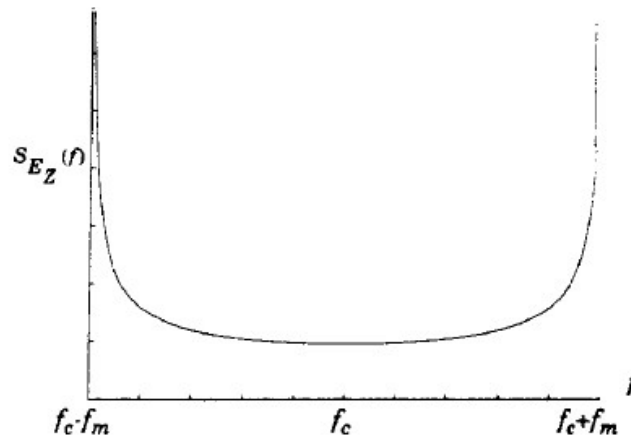


Figure 4.20
Doppler power spectrum for an unmodulated CW carrier [From [Gan72]] © IEEE].

$$S_{bbE_z}(f) = \frac{1}{8\pi f_m} K \left[\sqrt{1 - \left(\frac{f}{2f_m}\right)^2} \right] \quad (4.79)$$

where $K[\cdot]$ is the complete elliptical integral of the first kind. Equation (4.79) is not intuitive and is a result of the temporal correlation of the received signal when passed through a nonlinear envelope detector. Figure 4.21 illustrates the baseband spectrum of the received signal after envelope detection.

The spectral shape of the Doppler spread determines the time domain fading waveform and dictates the temporal correlation and fade slope behaviors. Rayleigh fading simulators must use a fading spectrum such as equation (4.78) in order to produce realistic fading waveforms that have proper time correlation.

1.2 Simulation of Clarke and Gans Fading Model

It is often useful to simulate multipath fading channels in hardware or software. A popular simulation method uses the concept of in-phase and quadrature modulation paths to produce a simulated signal with spectral and temporal characteristics very close to measured data.

As shown in Figure 4.22, two independent Gaussian low pass noise sources are used to produce in-phase and quadrature fading branches. Each Gaussian source may be formed by summing two independent Gaussian random variables which are orthogonal (i.e., $g = a + jb$, where a and b are real Gaussian random variables and g is complex Gaussian). By using the spectral filter defined by equation (4.78) to shape the random signals in the frequency domain, accurate time domain waveforms of Doppler fading can be produced by using an inverse fast Fourier transform (IFFT) at the last stage of the simulator.

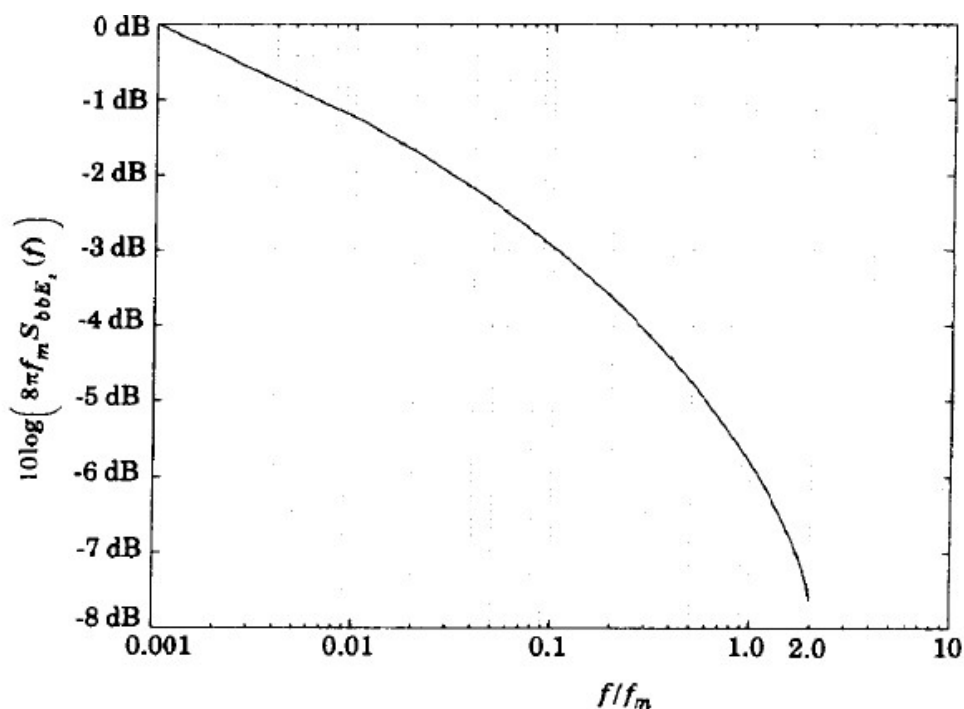


Figure 4.21
Baseband power spectral density of a CW Doppler signal after envelope detection.

Smith [Smi75] demonstrated a simple computer program that implements Figure 4.22(b). His method uses a complex Gaussian random number generator (noise source) to produce a baseband line spectrum with complex weights in the positive frequency band. The maximum frequency component of the line spectrum is f_m . Using the property of real signals, the negative frequency components are constructed by simply conjugating the complex Gaussian values obtained for the positive frequencies. Note that the IFFT of this signal is a purely real Gaussian random process in the time domain which is used in one of the quadrature arms shown in Figure 4.22. The random valued line spectrum is then multiplied with a discrete frequency representation of $\sqrt{S_{E_r}(f)}$ having the same number of points as the noise source. To handle the case where equation (4.78) approaches infinity at the passband edge, Smith truncated the value of $S_{E_r}(f_m)$ by computing the slope of the function at the sample frequency just prior to the passband edge and extended the slope to the passband edge. Simulations using the architecture in Figure 4.22 are usually implemented in the frequency domain using complex Gaussian line spectra to take advantage of easy implementation of equation (4.78). This, in turn, implies that the low pass Gaus-

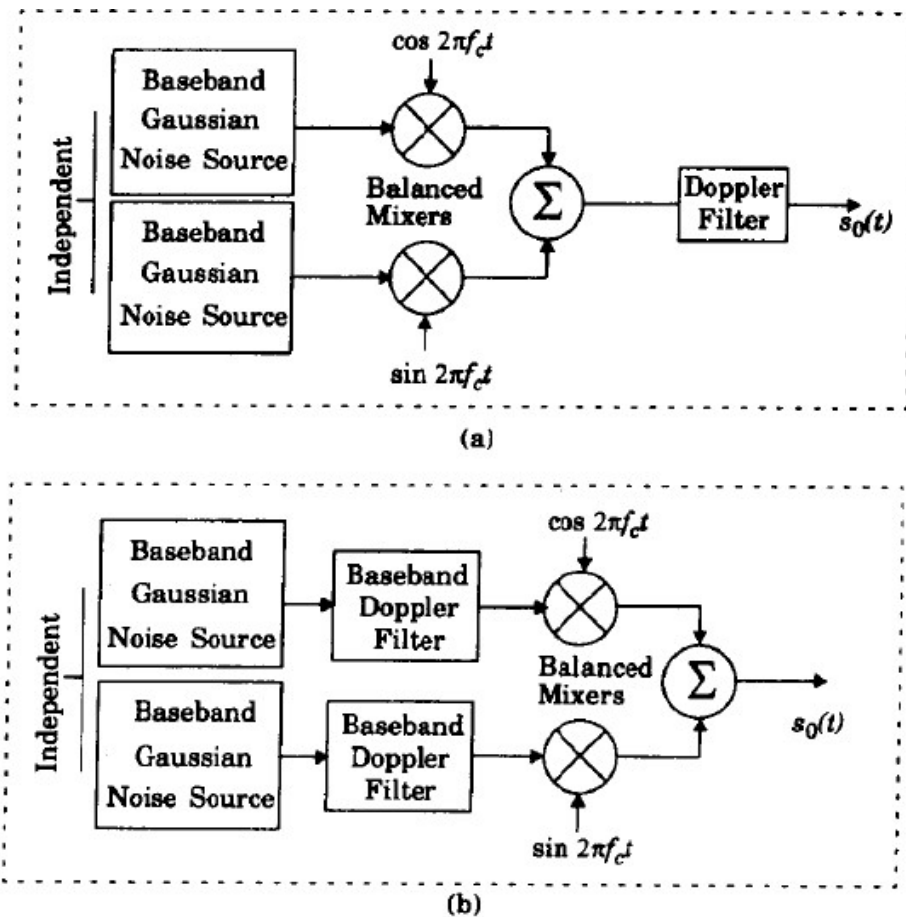


Figure 4.22

Simulator using quadrature amplitude modulation with (a) RF Doppler filter and (b) baseband Doppler filter.

sian noise components are actually a series of frequency components (line spectrum from $-f_m$ to f_m), which are equally spaced and each have a complex Gaussian weight. Smith's simulation methodology is shown in Figure 4.23.

To implement the simulator shown in Figure 4.23, the following steps are used:

- (1) Specify the number of frequency domain points (N) used to represent $\sqrt{S_{E_i}(f)}$ and the maximum Doppler frequency shift (f_m). The value used for N is usually a power of 2.
- (2) Compute the frequency spacing between adjacent spectral lines as $\Delta f = 2f_m / (N - 1)$. This defines the time duration of a fading waveform, $T = 1 / \Delta f$.
- (3) Generate complex Gaussian random variables for each of the $N/2$ positive frequency components of the noise source.
- (4) Construct the negative frequency components of the noise source by conju-

gating positive frequency values and assigning these at negative frequency values.

- (5) Multiply the in-phase and quadrature noise sources by the fading spectrum $\sqrt{S_{E_i}(f)}$.
- (6) Perform an IFFT on the resulting frequency domain signals from the in-phase and quadrature arms to get two N -length time series, and add the squares of each signal point in time to create an N -point time series like under the radical of equation (4.67).
- (7) Take the square root of the sum obtained in step 6 to obtain an N point time series of a simulated Rayleigh fading signal with the proper Doppler spread and time correlation.

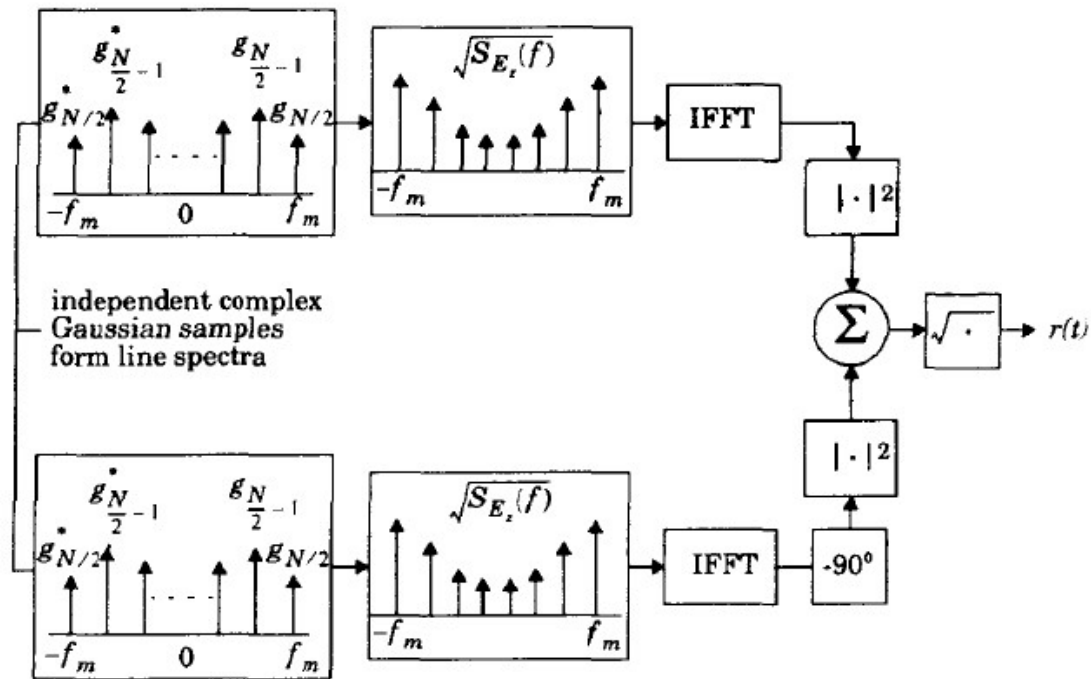


Figure 4.23

Frequency domain implementation of a Rayleigh fading simulator at baseband

Several Rayleigh fading simulators may be used in conjunction with variable gains and time delays to produce frequency selective fading effects. This is shown in Figure 4.24.

By making a single frequency component dominant in amplitude within $\sqrt{S_{E_i}(f)}$, the fading is changed from Rayleigh to Ricean. This can be used to alter the probability distributions of the individual multipath components in the simulator of Figure 4.24.

To determine the impact of flat fading on an applied signal $s(t)$, one merely needs to multiply the applied signal by $r(t)$, the output of the fading simulator. To determine the impact of more than one multipath component, a convolution must be performed as shown in Figure 4.24.

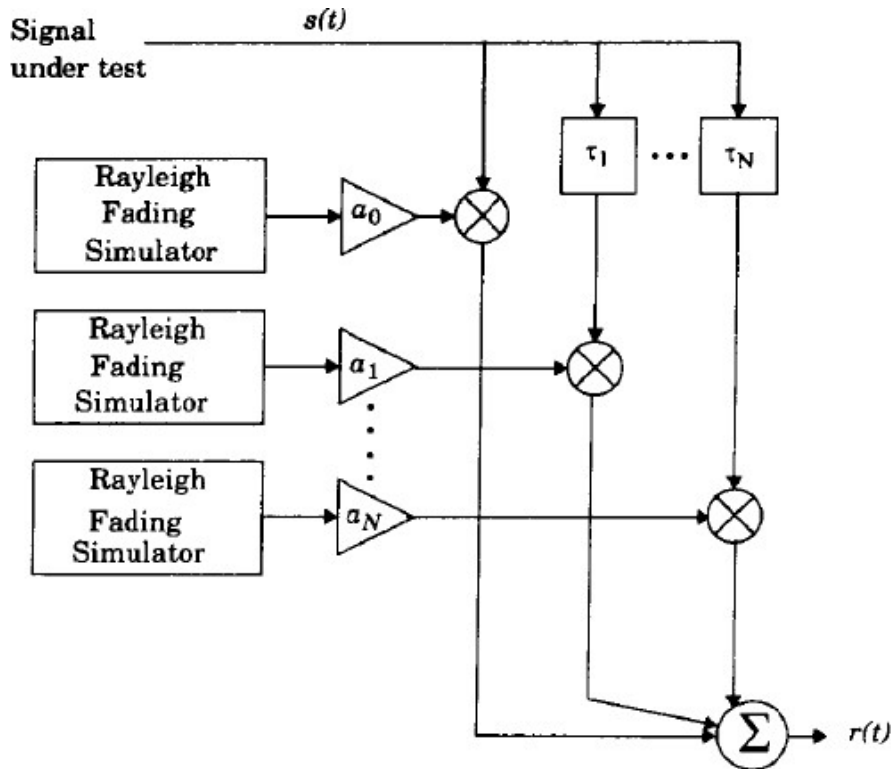


Figure 4.24

A signal may be applied to a Rayleigh fading simulator to determine performance in a wide range of channel conditions. Both flat and frequency selective fading conditions may be simulated, depending on gain and time delay settings.

1.3 Level Crossing and Fading Statistics

Rice computed joint statistics for a mathematical problem which is similar to Clarke's fading model [Cla68], and thereby provided simple expressions for computing the average number of level crossing and the duration of fades. The *level crossing rate* (LCR) and *average fade duration* of a Rayleigh fading signal are two important statistics which are useful for designing error control codes and diversity schemes to be used in mobile communication systems, since it becomes possible to relate the time rate of change of the received signal to the signal level and velocity of the mobile.

The *level crossing rate* (LCR) is defined as the expected rate at which the Rayleigh fading envelope, normalized to the local rms signal level, crosses a specified level in a positive-going direction. The number of level crossings per second is given by

$$N_R = \int_0^{\infty} \dot{r} p(R, \dot{r}) d\dot{r} = \sqrt{2\pi} f_m \rho e^{-\rho^2} \quad (4.80)$$

where \dot{r} is the time derivative of $r(t)$ (i.e., the slope), $p(R, \dot{r})$ is the joint density function of r and \dot{r} at $r = R$, f_m is the maximum Doppler frequency and $\rho = R/R_{rms}$ is the value of the specified level R , normalized to the local rms amplitude of the fading envelope [Jak74]. Equation (4.80) gives the value of N_R , the average number of level crossings per second at specified R . The level crossing rate is a function of the mobile speed as is apparent from the presence of f_m in equation (4.80). There are few crossings at both high and low levels, with the maximum rate occurring at $\rho = 1/\sqrt{2}$, (i.e., at a level 3 dB below the rms level). The signal envelope experiences very deep fades only occasionally, but shallow fades are frequent.

Example 4.6

For a Rayleigh fading signal, compute the positive-going level crossing rate for $\rho = 1$, when the maximum Doppler frequency (f_m) is 20 Hz. What is the maximum velocity of the mobile for this Doppler frequency if the carrier frequency is 900 MHz?

Solution to Example 4.6

Using equation (4.80), the number of zero level crossings is

$$N_R = \sqrt{2\pi} (20) (1) e^{-1} = 18.44 \text{ crossings per second}$$

The maximum velocity of the mobile can be obtained using the Doppler relation, $f_{d,max} = v/\lambda$.

Therefore velocity of the mobile at $f_m = 20$ Hz is

$$v = f_m \lambda = 20 \text{ Hz} (1/3 \text{ m}) = 6.66 \text{ m/s} = 24 \text{ km/hr}$$

The *average fade duration* is defined as the average period of time for which the received signal is below a specified level R . For a Rayleigh fading signal, this is given by

$$\bar{\tau} = \frac{1}{N_R} Pr[r \leq R] \quad (4.81)$$

where $Pr[r \leq R]$ is the probability that the received signal r is less than R and is given by

$$Pr[r \leq R] = \frac{1}{T} \sum_i \tau_i \quad (4.82)$$

where τ_i is the duration of the fade and T is the observation interval of the fading signal. The probability that the received signal r is less than the threshold R is found from the Rayleigh distribution as

$$Pr_r[r \leq R] = \int_0^R p(r) dr = 1 - \exp(-\rho^2) \quad (4.83)$$

where $p(r)$ is the pdf of a Rayleigh distribution. Thus, using equations (4.80), (4.81), and (4.83), the average fade duration as a function of ρ and f_m can be expressed as

$$\bar{\tau} = \frac{e^{\rho^2} - 1}{\rho f_m \sqrt{2\pi}} \quad (4.84)$$

1.4 Two-ray Rayleigh Fading Model

Clarke's model and the statistics for Rayleigh fading are for flat fading conditions and do not consider multipath time delay. In modern mobile communication systems with high data rates, it has become necessary to model the effects of multipath delay spread as well as fading. A commonly used multipath model is an independent Rayleigh fading 2-ray model (which is a specific implementation of the generic fading simulator shown in Figure 4.24). Figure 4.25 shows a block diagram of the 2-ray independent Rayleigh fading channel model. The impulse response of the model is represented as

$$h_b(t) = \alpha_1 \exp(j\phi_1) \delta(t) + \alpha_2 \exp(j\phi_2) \delta(t - \tau) \quad (4.85)$$

where α_1 and α_2 are independent and Rayleigh distributed, ϕ_1 and ϕ_2 are independent and uniformly distributed over $[0, 2\pi]$, and τ is the time delay between the two rays. By setting α_2 equal to zero, the special case of a flat Rayleigh fading channel is obtained as

$$h_b(t) = \alpha_1 \exp(j\phi_1) \delta(t) \quad (4.86)$$

By varying τ , it is possible to create a wide range of frequency selective fading effects. The proper time correlation properties of the Rayleigh random variables α_1 and α_2 are guaranteed by generating two independent waveforms, each produced from the inverse Fourier transform of the spectrum described in Section 4.7.2.

Multiple Access Techniques

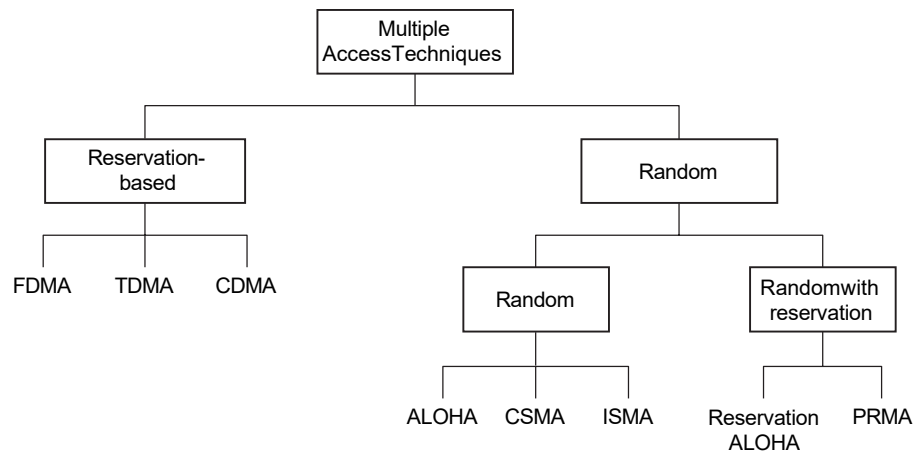
6.1 Introduction

In Chapter 5 we discussed that cellular systems divide a geographic region into cells where mobile units in each cell communicate with the cell's base station. The goal in the design of a cellular system is to be able to handle as many calls as possible in a given bandwidth with the specified blocking probability (reliability).

Multiplexing deals with the division of the resources to create multiple channels. Multiplexing can create channels in frequency, time, etc., and the corresponding terms are then frequency division multiplexing (FDM), time division multiplexing (TDM), etc. [1,3]. Since the amount of spectrum available is limited, we need to find ways to allow multiple users to share the available spectrum simultaneously. Shared access is used to implement a multiple access scheme when access by many users to a channel is required [13,14,15]. For example, one can create multiple channels using TDM, but each of these channels can be accessed by a group of users using the ALOHA multiple access scheme [8,9]. The multiple access schemes can be either reservation-based or random.

Multiple access schemes allow many users to share the radio spectrum. Sharing the bandwidth efficiently among users is one of the main objectives of multiple access schemes [16,17]. The variability of wireless channels presents both challenges and opportunities in designing multiple access communications systems. Multiple access strategy has an impact on robustness and interference levels generated in other cells. Therefore, multiple access schemes are designed to maintain orthogonality and reduce interference effects [10].

Multiple access schemes can be classified as *reservation-based* multiple access (e.g., FDMA, TDMA, CDMA) [4,5] and *random* multiple access (e.g., ALOHA, CSMA) (see Figure 6.1) [9,23]. If data traffic is continuous and a small transmission delay is required (for example in voice communication) reservation-based multiple access is used. The family of reservation-based multiple access includes frequency division multiple access (FDMA), time division multiple access (TDMA), and code division multiple access (CDMA) [6,7,12,21,22]. In many wireless systems for voice communication, the control channel is based on random multiple access and the communication channel is based on FDMA, TDMA, or CDMA. The reservation-based multiple access technique has a disadvantage in that once the channel is assigned, it remains idle if the user has nothing to transmit,



ISMA : Idle Signal Casting Multiple
Access PRMA: Packet Reservation Multiple Access

Figure 6.1 Multiple access schemes.

while other users may have data waiting to be transmitted. This problem is critical when data generation is random and has a high peak-rate to average-rate ratio. In this situation, random multiple access is more efficient, because a communication channel is shared by many users and users transmit their data in a random or partially coordinated fashion. ALOHA and carrier sense multiple access (CSMA) are examples of random multiple access [8]. If the data arrives in a random manner, and the data length is large, then random multiple access combined with a reservation protocol will perform better than both random- and reservation-based schemes.

We first focus on the reservation-based multiple access schemes including narrowband channelized and wideband nonchannelized systems for wireless communications. We discuss access technologies—FDMA, TDMA, and CDMA. We examine FDMA and TDMA from a capacity, performance, and spectral efficiency viewpoint. As networks have evolved, the demand for higher capacities has encouraged researchers and system designers to examine access schemes that are even more spectrally efficient than TDMA. Therefore, we also examine the CDMA system. Work in standards bodies around the world indicates that the 3G/4G wireless systems are evolving to wideband CDMA (WCDMA) to achieve high efficiencies and high access data rates. The later part of the chapter is devoted to the discussion of random multiple access schemes.

6.2 Narrowband Channelized Systems

Traditional architectures for analog and digital wireless systems are channelized [6,11]. In a channelized system, the total spectrum is divided into a large number of

relatively narrow radio channels that are defined by carrier frequency. Each radio channel consists of a pair of frequencies. The frequency used for transmission from the base station to the mobile station is called the *forward channel* (downlink channel) and the frequency used for transmission from the mobile station to the base station is called the *reverse channel* (uplink channel). A user is assigned both frequencies for the duration of the call. The forward and reverse channels are assigned widely separated frequencies to keep the interference between transmission and reception to a minimum.

A narrowband channelized system demands precise control of output frequencies for an individual transmitter. In this case, the transmission by a given mobile station occurs within the specified narrow bandwidth to avoid interference with adjacent channels. The tightness of bandwidth limitations plays a dominant role in the evaluation and selection of modulation technique. It also influences the design of transmitter and receiver elements, particularly the filters which can greatly affect the cost of a mobile station.

A critical issue with regulators and operators around the world is how efficiently the radio spectrum is being used. Regulatory bodies want to encourage competition for cellular services. Thus for a given availability of bandwidth, more operators can be licensed. For a particular operator, a more efficient technology can support more users within the assigned spectrum and thus increase profits.

When we examine efficiencies of various technologies, we find that each system has made different trade-offs in determining the optimum method for access. Some of the parameters that are used in the trade-off are bandwidth per user, guard bands between channels, frequency reuse among different cells in the system, the signal-to-noise and signal-to-interference ratio, the methods of channel and speech coding, and the complexity of the system.

The first-generation analog cellular systems showed signs of capacity saturation in major urban areas, even with a modest total user population. A major capacity increase was needed to meet future demand. Several digital techniques were deployed to solve the capacity problem of analog cellular systems. There are two basic types of systems whereby a fixed spectrum resource is partitioned and shared among different users [13,16]. The channels are created by dividing the total system bandwidth into frequency channels through the use of FDM and then further dividing each frequency channel into time channels through the use of TDM. Most systems use a combination of FDMA and TDMA.

6.2.1 Frequency Division Duplex (FDD) and Time Division Duplex (TDD) System

Many cellular systems (such as AMP, GSM, DAMP, etc.) use *frequency division duplex (FDD)* in which the transmitter and receiver operate simultaneously on different frequencies. Separation is provided between the downlink and uplink channels to avoid the transmitter causing self-interference to its receiver. Other

precautions are also needed to prevent self interference, such as the use of two antennas, or alternatively one antenna with a duplexer (a special design of RF filters protecting the receiver from the transmit frequency). A duplexer adds weight, size, and cost to a radio transceiver and can limit the minimum size of a subscriber unit.

A cellular system can be designed to use one frequency band by using *time division duplex (TDD)*. In TDD a bidirectional flow of information is achieved using the simplex-type scheme by automatically alternating in time the direction of transmission on a single frequency. At best TDD can only provide a quasi-simultaneous bidirectional flow, since one direction must be off while the other is using the frequency. However, with a high enough transmission rate on the channel, the off time is not noticeable during conversations, and with a digital speech system, the only effect is a very short delay.

The amount of spectrum required for both FDD and TDD is the same. The difference lies in the use of two bands of spectrum separated by the required bandwidth for FDD, whereas TDD requires only one band of frequencies but twice the bandwidth. It may be easier to find a single band of unassigned frequencies than finding two bands separated by the required bandwidth.

With TDD systems, the transmit time slot and the receiver time slot of the subscriber unit occur at different times. With the use of a simple RF switch in the subscriber unit, the antenna can be connected to the transmitter when a transmit burst is required (thus disconnecting the receiver from the antenna) and to the receiver for the incoming signal. The RF switch thus performs the function of the duplexer, but is less complex, smaller in size, and less costly. TDD uses a burst mode scheme like TDMA and therefore also does not require a duplexer. Since the bandwidth of a TDD channel is twice that of a transmitter and receiver in an FDD system, RF filters in all the transmitters and receivers for TDD systems must be designed to cover twice the bandwidth of FDD system filters.

6.2.2 Frequency Division Multiple Access

The FDMA is the simplest scheme used to provide multiple access. It separates different users by assigning a different carrier frequency (see Figure 6.2). Multiple users are isolated using bandpass filters. In FDMA, signals from various users are assigned different frequencies, just as in an analog system. Frequency guard bands are provided between adjacent signal spectrums to minimize crosstalk between adjacent channels.

The advantages and disadvantages of FDMA with respect to TDMA or CDMA are:

Advantages

1. Capacity can be increased by reducing the information bit rate and using an efficient digital speech coding scheme (See Chapter 8) [20].

- 2. Technological advances required for implementation are simple. A system can be configured so that improvements in terms of a lower bit rate speech coding could be easily incorporated.
- 3. Hardware simplicity, because multiple users are isolated by employing simple bandpass filters.

Disadvantages

- 1. The system architecture based on FDMA was implemented in first-generation analog systems such as advanced mobile phone system (AMPS) or total access communication system (TACS). The improvement in capacity depends on operation at a reduced signal-to-interference (S/I) ratio. But the narrowband digital approach gives only limited advantages in this regard so that modest capacity improvements could be expected from the allocated spectrum.
- 2. The maximum bit-rate per channel is fixed and small, inhibiting the flexibility in bit-rate capability that may be a requirement for computer file transfer in some applications in the future.

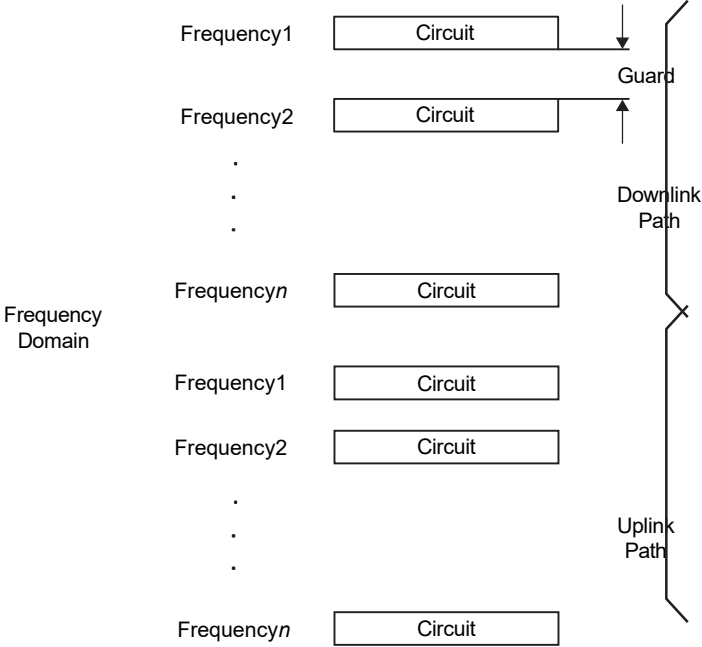


Figure 6.2 FDMA/FDD channel architecture.

3. Inefficient use of spectrum, in FDMA if a channel is not in use, it remains idle and cannot be used to enhance the system capacity.
4. Crosstalk arising from adjacent channel interference is produced by non-linear effects.

6.2.3 Time Division Multiple Access

In a TDMA system, each user uses the whole channel bandwidth for a fraction of time (see Figure 6.3) compared to an FDMA system where a single user occupies the channel bandwidth for the entire duration (see Figure 6.2) [2]. In a TDMA system, time is divided into equal time intervals, called *slots*. User data is transmitted in the slots. Several slots make up a frame. Guard times are used between each user's transmission to minimize crosstalk between channels (see Figure 6.4).

Each user is assigned a frequency and a time slot to transmit data. The data is transmitted via a radio-carrier from a base station to several active mobiles in the downlink. In the reverse direction (uplink), transmission from mobiles to base stations is time-sequenced and synchronized on a common frequency for TDMA. The preamble carries the address and synchronization information that both the base station and mobile stations use for identification.

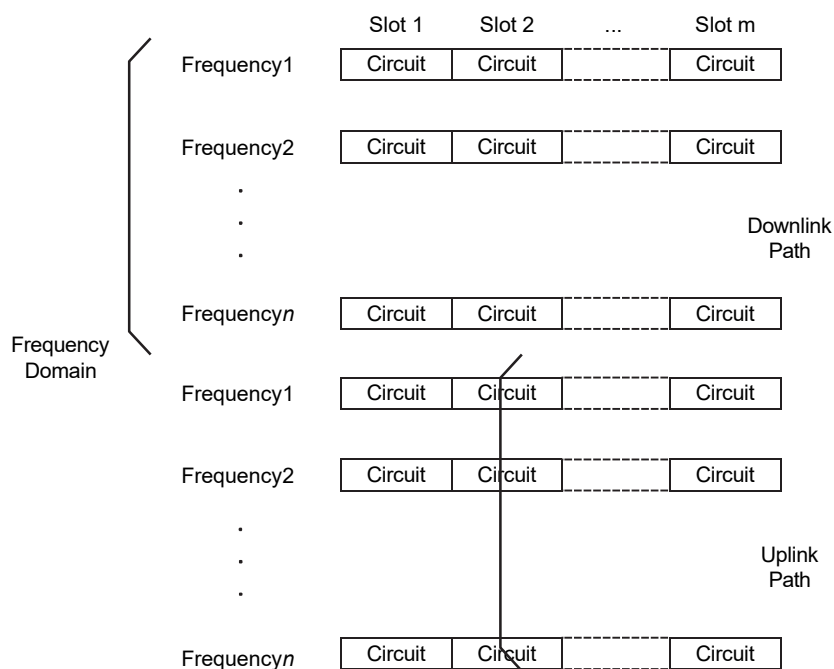


Figure 6.3 TDMA/FDD channel architecture.

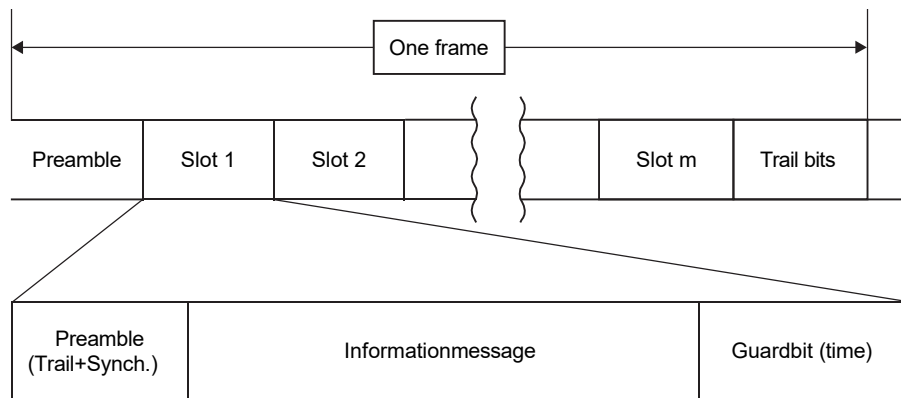


Figure 6.4 TDMA frame.

In a TDMA system, the user can use multiple slots to support a wide range of bit rates by selecting the lowest multiplexing rate or multiple of it. This enables supporting a variety of voice coding techniques at different bit rates with different voice qualities. Similarly, data communications customers could use the same kinds of schemes, choosing and paying for the digital data rate as required. This would allow customers to request and pay for a bandwidth on demand.

Depending on the data rate used and the number of slots per frame, a TDMA system can use the entire bandwidth of the system or can employ an FDD scheme. The resultant multiplexing is a mixture of frequency division and time division. The entire frequency band is divided into a number of duplex channels (about 350 to 400 kHz). These channels are deployed in a frequency-reuse pattern, in which radio-port frequencies are assigned using an autonomous adaptive frequency assignment algorithm. Each channel is configured in a TDM mode for the downlink direction and a TDMA mode for the uplink direction.

The advantages and disadvantages of TDMA are:

Advantages

1. TDMA permits a flexible bit rate, not only for multiples of the basic single channel rate but also submultiples for low bit rate broadcast-type traffic.
2. TDMA offers the opportunity for frame-by-frame monitoring of signal strength/bit error rate to enable either mobile or base stations to initiate and execute handoffs.
3. TDMA, when used exclusively and not with FDMA, utilizes bandwidth more efficiently because no frequency guard band is required between channels.

4. TDMA transmits each signal with sufficient guard time between time slots to accommodate time inaccuracies because of clock instability, delay spread, transmission delay because of propagation distance, and the tails of signal pulse because of transient responses.

Disadvantages

1. For mobiles and particularly for hand-sets, TDMA on the uplink demands high peak power in transmit mode, that shortens battery life.
2. TDMA requires a substantial amount of signal processing for matched filtering and correlation detection for synchronizing with a time slot.
3. TDMA requires synchronization. If the time slot synchronization is lost, the channels may collide with each other.
4. One complicating feature in a TDMA system is that the propagation time for a signal from a mobile station to a base station varies with its distance to the base station.

6.3 Spectral Efficiency

An efficient use of the spectrum is the most desirable feature of a mobile communications system. To realize this, a number of techniques have been proposed or already implemented. Some of these techniques used to improve spectral efficiency are reducing the channel bandwidth, information compression (low-rate speech coding), variable bitrate codec (see Chapter 8), improved channel assignment algorithms (dynamic channel assignment), and so on [11,17,19]. Spectral efficiency of a mobile communications system shows how efficiently the spectrum is used by the system. Spectral efficiency of a mobile communication system depends on the choice of a multiple access scheme. The measure of spectral efficiency enables one to estimate the capacity of a mobile communications system.

The overall spectral efficiency of a mobile communications system can be estimated by knowing the *modulation* and the *multiple access* spectral efficiencies separately [16].

6.3.1 Spectral Efficiency of Modulation

The spectral efficiency with respect to modulation is defined as [16]:

$$h_m = \frac{(\text{Total Number of Channels Available in the System})}{(\text{Bandwidth})(\text{Total Coverage Area})} \quad (6.1a)$$

$$h_m = \frac{B_w \times N_c}{B_c \times \frac{N}{w} \times \frac{A}{c \times c}} \quad (6.1b)$$

$$h_m = \frac{1}{B_c \times N \times A_c} \text{ Channels/MHz/km}^2 \quad (6.1c)$$

where:

h_m = modulation efficiency (channels/MHz/km²)

B_w = bandwidth of the system (MHz)

B_c = channel spacing (MHz)

N_c = total number of cells in the covered area

N = frequency reuse factor of the system (or cluster size)

A_c = area covered by a cell (km²)

Equation 6.1c indicates that the spectral efficiency of modulation does not depend on the bandwidth of the system. It only depends on the channel spacing, the cell area, and the frequency reuse factor, N . By reducing the channel spacing, the spectral efficiency of modulation for the system is increased, provided the cell area (A_c) and reuse factor (N) remain unchanged. If a modulation scheme can be designed to reduce N then more channels are available in a cell and efficiency is improved.

Another definition of spectral efficiency of modulation is Erlangs/MHz/km²

$$h_m = \frac{\text{Maximum Total Traffic Carried by System}}{(\text{System Bandwidth})(\text{Total Coverage Area})} \quad (6.2a)$$

$$h_m = \frac{\text{Total Traffic Carried by } \left(\frac{B_w/B_c}{N}\right) \text{ Channels}}{B_w A_c} \quad (6.2b)$$

By introducing the trunking efficiency factor, h_t in Equation 6.2b (< 1 , it is a function of the blocking probability and number of available channels per cell), the total traffic carried by the system is given as:

$$h_m = \frac{h_t \left(\frac{B_w/B_c}{N}\right)}{B_w A_c} \quad (6.2c)$$

$$h_m = \frac{h_t}{B_c N A_c} \quad (6.2d)$$

where:

h_t is a function of the blocking probability and the total number of available channels per cell $\left[\frac{B_w/B_c}{N}\right]$

Based on Equation 6.2d we can conclude:

1. The voice quality will depend on the frequency reuse factor, N , which is a function of the signal-to-interference (S/I) ratio of the modulation scheme used in the mobile communications system (see Chapter 5).
2. The relationship between system bandwidth, B_w , and the amount of traffic carried by the system is nonlinear, i.e., for a given percentage increase in B_w , the increase in the traffic carried by the system is more than the increase in B_w .
3. From the average traffic per user ($Erlang/user$) during the busy hour and $Erlang/MHz/km^2$, the capacity of the system in terms of $users/km^2/MHz$ can be obtained.
4. The spectral efficiency of modulation depends on the blocking probability.

Example 6.1

In the GSM800 digital channelized cellular system, the one-way bandwidth of the system is 12.5 MHz. The RF channel spacing is 200 kHz. Eight users share each RF channel and three channels per cell are used for control channels. Calculate the spectral efficiency of modulation (for a dense metropolitan area with small cells) using the following parameters:

- Area of a cell = $8 km^2$
- Total coverage area = $4000 km^2$
- Average number of calls per user during the busy hour = 1.2
- Average holding time of a call = 100 seconds
- Call blocking probability = 2%
- Frequency reuse factor = 4

Solution

$$\text{Number of 200 kHz RF channels} = \frac{12.5 \times 1000}{200} = 62$$

$$\text{Number of traffic channels} = 62 \times 8 = 496$$

$$\text{Number of signaling channels per cell} = 3$$

$$\text{Number of traffic channels per cell} = \frac{496}{4} = 121$$

$$\text{Number of cells} = \frac{4000}{8} = 500$$

With 2% blocking for an omnidirectional case, the total traffic carried by 121 channels (using Erlang-B tables) = $108.4(1.0 - 0.02) = 106.2$ Erlangs/cell or 13.28 Erlangs/ km^2

$$\begin{aligned} \text{Number of calls per hour per cell} &= \frac{106.2 \times 3600}{100} = 3823, \text{ calls/hour/km}^2 = \\ &= \frac{3823}{8} = 477.9 \text{ calls/hour/km}^2 \\ \text{Max. number of users/cell/hour} &= \frac{3823}{1.2} = 3186, \text{ users/hour/channel} = \frac{3186}{121} \end{aligned}$$

26.33

$$h_m = \frac{(\text{Erlangs per cell}) \times \text{no. of cells}}{B_w \times \text{Coverage Area}} = \frac{106.2 \times 500}{12.5 \times 4000} = 1.06 \text{ Erlangs/MHz/km}^2$$

6.3.2 Multiple Access Spectral Efficiency

Multiple access spectral efficiency is defined as the ratio of the total time or frequency dedicated for traffic transmission to the total time or frequency available to the system. Thus, the multiple access spectral efficiency is a dimensionless number with an upper limit of unity.

In FDMA, users share the radio spectrum in the frequency domain. In FDMA, the multiple access efficiency is reduced because of guard bands between channels and also because of signaling channels. In TDMA, the efficiency is reduced because of guard time and synchronization sequence.

FDMA Spectral Efficiency

For FDMA, multiple access spectral efficiency is given as:

$$h_a = \frac{B_c N_T}{B_w} \leq 1 \quad (6.3)$$

where:

h_a = multiple access spectral efficiency

N_T = total number of traffic channels in the covered area

B_c = channel spacing

B_w = system bandwidth

Example 6.2

In a first-generation AMPS system where there are 395 channels of 30 kHz each in a bandwidth of 12.5 MHz, what is the multiple access spectral efficiency for FDMA?

Solution

$$h_a = \frac{30 \times 395}{12.5 \times 1000} = 0.948$$

TDMA Spectral Efficiency

TDMA can operate as wideband or narrowband. In the wideband TDMA, the entire spectrum is used by each individual user. For the wideband TDMA, multiple access spectral efficiency is given as:

$$h_a = \frac{\tau M_t}{T_f} \quad (6.4)$$

where:

τ = duration of a time slot that carries data

T_f = frame duration

M_t = number of time slots per frame

In Equation 6.4 it is assumed that the total available bandwidth is shared by all users. For the narrowband TDMA schemes, the total band is divided into a number of sub-bands, each using the TDMA technique. For the narrowband TDMA system, frequency domain efficiency is not unity as the individual user channel does not use the whole frequency band available to the system. The multiple access spectral efficiency of the narrowband TDMA system is given as:

$$h_a = \left(\frac{\tau M_t}{T_f} \right) \left(\frac{B_u N_u}{B_w} \right) \quad (6.5)$$

where:

B_u = bandwidth of an individual user during his or her time slot

N_u = number of users sharing the same time slot in the system, but having access to different frequency sub-bands

6.3.3 Overall Spectral Efficiency of FDMA and TDMA Systems

The overall spectral efficiency, h , of a mobile communications system is obtained by considering both the modulation and multiple access spectral efficiencies

$$h = h_m h_a \quad (6.6)$$

Example 6.3

In the North American Narrowband TDMA cellular system, the one-way bandwidth of the system is 12.5 MHz. The channel spacing is 30 kHz and the total number of voice channels in the system is 395. The frame duration is 40 ms, with six time slots per frame. The system has an individual user data rate of 16.2 kbps in which the speech with error protection has a rate of 13 kbps. Calculate the multiple access spectral efficiency of the TDMA system.

Solution

The timeslot duration that carries data: $\tau = \binom{13}{16.2} \left(\frac{40}{6}\right) = 5.35 \text{ ms}$

$T_f = 40 \text{ ms}, M_t = 6, N_u = 395, B_u = 30 \text{ kHz}, \text{ and } B_w = 12.5 \text{ MHz}$

$$h_a = \frac{5.35 \times 6 \times 30 \times 395}{40 \times 12500} = 0.76$$

The overhead portion of the frame = $1.0 - 0.76 = 24\%$

Capacity and Frame Efficiency of a TDMA System**Cell Capacity**

The cell capacity is defined as the maximum number of users that can be supported simultaneously in each cell.

The capacity of a TDMA system is given by [16]:

$$N_u = \frac{h_b \mu \times B_w}{v_f \overline{RN}} \quad (6.7)$$

where:

N_u = number of channels (mobile users) per cell

h_b = bandwidth efficiency factor (< 1.0)

μ = bit efficiency (= 2 bit/symbol for QPSK, = 1 bit/symbol for GMSK as used in GSM)

v_f = voice activity factor (equal to one for TDMA)

B_w = one-way bandwidth of the system

R = information (bit rate plus overhead) per user

N = frequency reuse factor

$$\text{Spectral efficiency } h = \frac{N_u \times R}{B_w} \text{ bit/sec/Hz} \quad (6.8)$$

Example 6.4

Calculate the capacity and spectral efficiency of a TDMA system using the following parameters: bandwidth efficiency factor $h_b = 0.9$, bit efficiency (with QPSK) $\mu = 2$, voice activity factor $v_f = 1.0$, one-way system bandwidth $B_w = 12.5 \text{ MHz}$, information bit rate $R = 16.2 \text{ kbps}$, and frequency reuse factor $N = 19$.

Solution

$$N_u = \frac{0.9 \times 2}{1.0} \times \frac{12.5 \times 10^6}{16.2 \times 10^3 \times 19}$$

$N = 73.1$ (say 73 mobile users per cell)

$$\text{Spectral efficiency } h = \frac{73 \times 16.2}{12.5 \times 1000} = 0.094 \text{ bit/sec/Hz}$$

Efficiency of a TDMA Frame

We refer to Figure 6.4 that shows a TDMA frame. The number of overhead bits per frame is:

$$b_0 = N_r b_r + N_t b_p + (N_t + N_r) b_g \quad (6.9)$$

where:

N_r = number of reference bursts per frame

N_t = number of traffic bursts (slots) per frame

b_r = number of overhead bits per reference burst

b_p = number of overhead bits per preamble per slot

b_g = number of equivalent bits in each guard time interval

The total number of bits per frame is:

$$b_T = T_f \times R_{rf} \quad (6.10a)$$

where:

T_f = frame duration

R_{rf} = bitrate of the RF channel

$$\text{Frame efficiency } h = (1 - b_0/b_T) \times 100\% \quad (6.10b)$$

It is desirable to maintain the efficiency of the frame as high as possible.

The number of bits per data channel (user) per frame is $b_c = R T_f$, where R = bitrate of each channel (user).

$$\text{No. of channels/frame } N_{CF} = \frac{(\text{Total Data Bits})/(\text{frame})}{(\text{Bits per Channel})/(\text{frame})}$$

$$N_{CF} = \frac{h R_{rf} T_f}{R T_f} \quad (6.11a)$$

$$N_{CF} = \frac{h R_{rf}}{R} \quad (6.11b)$$

Equation 6.11b indicates the number of times slots per frame.

Example 6.5

Consider the GSM TDMA system with the following parameters:

$$N_r = 2$$

$$N_t = 24 \text{ frames of } 120 \text{ ms each with eight time slots per frame}$$

$$b_r = 148 \text{ bits in each of 8 time slots}$$

$$b_p = 34 \text{ bits in each of 8 time slots}$$

$$b_g = 8.25 \text{ bits in each of 8 time slots}$$

$$T_f = 120 \text{ ms}$$

$$R_{rf} = 270.8333333 \text{ kbps}$$

$$R = 22.8 \text{ kbps}$$

Calculate the frame efficiency and the number of channels per frame.

Solution

$$b_0 = 2 \times (8 \times 148) + 24 \times (8 \times 34) + 8 \times 8.25 = 10,612 \text{ bits per frame}$$

$$b_T = 120 \times 10^{-3} \times 270.8333333 \times 10^3 = 32,500 \text{ bits per frame}$$

$$h = 1 - \left(\frac{10612}{32500} \right) \times 100 = 67.35\%$$

$$\text{Number of channels/frame} = \frac{0.6735 \times 270.8333333}{22.8} = 8$$

The last calculation, with an answer of 8 channels, confirms that our calculation of efficiency is correct.

6.4 WidebandSystems

In wideband systems, the entire system bandwidth is made available to each user, and is many times larger than the bandwidth required to transmit information. Such systems are known as *spread spectrum* (SS) systems. There are two fundamental types of spread spectrum systems: (1) direct sequence spread spectrum (DSSS) and (2) frequency hopping spread spectrum (FHSS) [3,26].

In a DSSS system, the bandwidth of the baseband information carrying signals from a different user is spread by different codes with a bandwidth much larger than that of the baseband signals (see Chapter 11 for details). The spreading codes used for different users are orthogonal or nearly orthogonal to each other. In the DSSS, the spectrum of the transmitted signal is much wider than the spectrum associated with the information rate. At the receiver, the same code is used for despreading to recover the baseband signal from the target user while suppressing the transmissions from all other users (see Figure 6.5).

One of the advantages of the DSSS system is that the transmission bandwidth exceeds the coherence bandwidth (see Chapter 3). The received signal, after despreading (see Chapter 11 for details), resolves into multiple signals with different time delays. A rake receiver (see Chapter 11) can be used to recover the multiple time

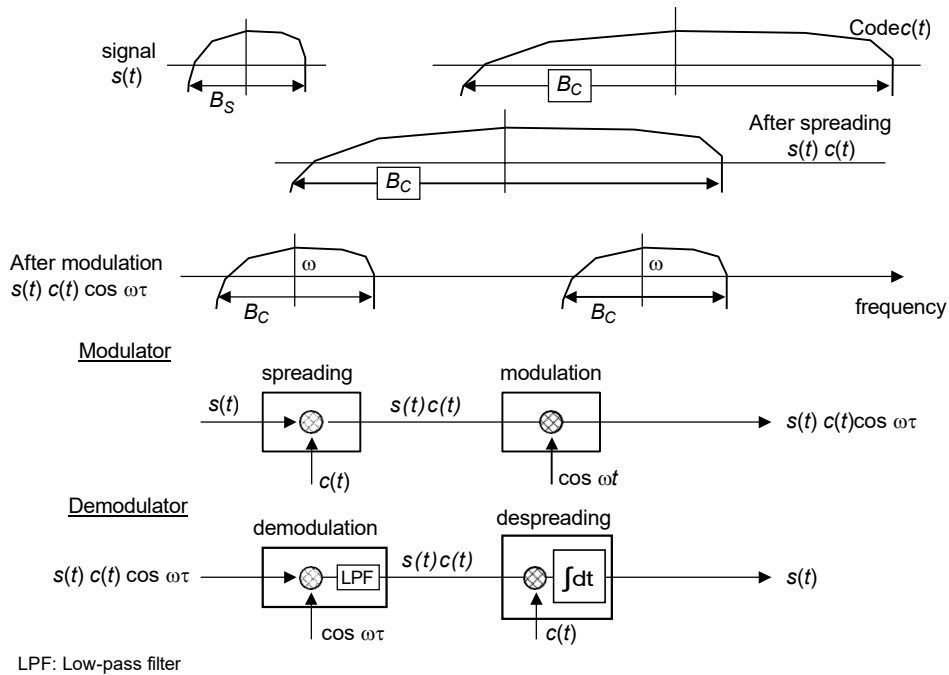


Figure 6.5 Direct sequence spread spectrum.

delayed signals and combine them into one signal to provide time diversity with a lower frequency of deep fades. Thus, the DSSS system provides an inherent robustness against mobile-channel degradations. Another potential benefit of a DSSS system is the greater resistance to interference effects in a frequency reuse situation. Also, there may be no hard limit on the number of mobile users who can simultaneously gain access. The capacity of a DSSS system depends upon the desired value of E_b/I_0 instead of resources (frequencies or time slots) as in FDMA or TDMA systems.

Frequency hopping (FH) is the periodic changing of the frequency or the frequency set associated with transmission (see Figure 6.6). If the modulation is M -ary frequency-shift-keying (MFSK) (see Chapter 9 for details), two or more frequencies are in the set that change at each hop. For other modulations, a single center or carrier frequency is changed at each hop.

An FH signal may be considered a sequence of modulated pulses with pseudo-random carrier frequencies. This set of possible carrier frequencies is called the hop set. Hopping occurs over a frequency band that includes a number of frequency channels. The bandwidth of a frequency channel is called the *instantaneous bandwidth* (B_I). The bandwidth of the frequency band over which the hopping occurs is called the total *hopping bandwidth* (B_H). The time duration between hops is called the *hop duration* or hopping period (T_H).

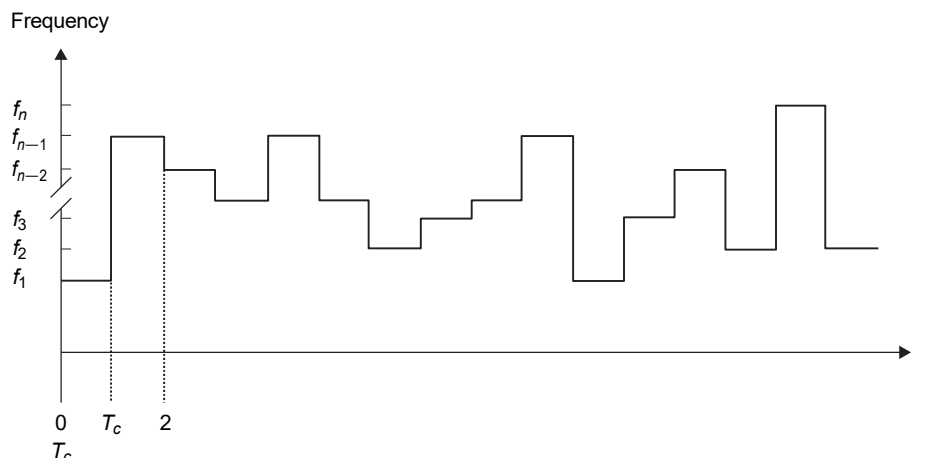


Figure 6.6 Frequency hopping spread spectrum system.

Frequency hopping can be classified as fast or slow. *Fast frequency hopping* occurs if there is frequency hop for each transmitted symbol. Thus, fast frequency hopping implies that the hopping rate equals or exceeds the information symbol rate. *Slow frequency hopping* occurs if two or more symbols are retransmitted in the time interval between frequency hops.

Frequency hopping allows communicators to hop out of frequency channels with interference or to hop out of fades. To exploit this capability, error-correcting codes, appropriate interleaving, and disjoint frequency channels are nearly always used. A frequency synthesizer is required for frequency hopping systems to convert a stable reference frequency into the various frequency of hop set.

Frequency hopping communicators do not often operate in isolation. Instead, they are usually elements of a network of frequency hopping systems that create mutual multiple-access interference. This network is called a frequency-hopping multiple-access (FHMA) network.

If the hoppers of an FHMA network all use the same M frequency channels, but coordinate their frequency transitions and their hopping sequence, then the multiple-access interference for a lightly loaded system can be greatly reduced compared to a non-hopped system. For the number of hopped signals (M_h) less than the number of channels (N_c), a coordinated hopping pattern can eliminate interference. As the number of hopped signals increases beyond N_c , then the interference will increase in proportion to the ratio of the number of signals to the number of channels. In the absence of fading or multipath interference, since there is no interference suppression system in frequency hopping, for a high channel loading the performance of a frequency hopping system is no better than a non-hopped system. Frequency hopping systems are best for light channel loadings in the presence of conventional non-hopped systems.

When fading or multipath interference is present, the frequency hopping system has better error performance than a non-hopped system. If the transmitter hops to a channel in a fade, the errors are limited in duration since the system will shortly hop to a new frequency where the fade may not be as deep.

Network coordination for frequency hopping systems are simpler to implement than that for DSSS systems because the timing alignments must be within a fraction of a hop duration, rather than a fraction of a sequence chip (narrow pulse). In general, frequency hopping systems reject interference by trying to avoid it, whereas DSSS systems reject interference by spreading it. The interleaving and error-correcting codes that are effective with frequency hopping systems are also effective with DSSS systems.

The major problems with frequency hopping systems with increasing hopping rates are the cost of the frequency synthesizer increases and its reliability decreases, and synchronization becomes more difficult.

In theory, a wideband system can be overlaid on existing, fully loaded, narrowband channelized systems (as an example, the IS-95 CDMA system overlays on existing AMPS [FDMA]). Thus, it may be possible to create a wideband network right on top of the narrowband cellular system using the same spectrum.

6.5 Comparison of FDMA, TDMA, and DS-SS-CDMA

The DSSS approach is the basis to implementation of the direct sequence code division multiple access (DS-SS-CDMA) technique introduced by Qualcomm. The DS-SS-CDMA has been used in commercial applications of mobile communications. The primary advantage of DS-SS-CDMA is its ability to tolerate a fair amount of interfering signals compared to FDMA and TDMA that typically cannot tolerate any such interference (Figure 6.7). As a result of the interference tolerance of CDMA, the problems of frequency band assignment and adjacent cell interference are greatly simplified. Also, flexibility in system design and deployment are significantly improved since interference to others is not a problem. On the other hand, FDMA and TDMA radios must be carefully assigned a frequency or time slot to assure that there is no interference with other similar radios. Therefore, sophisticated filtering and guard band protection is needed with FDMA and TDMA technologies. With DS-SS-CDMA, adjacent microcells share the same frequencies whereas with FDMA/TDMA it is not feasible for adjacent microcells to share the same frequencies because of interference. In both FDMA and TDMA systems, a time-consuming frequency planning task is required whenever a network changes, whereas no such frequency planning is needed for a CDMA network since each cell uses the same frequencies.

Capacity improvements with DS-SS-CDMA also result from voice activity patterns during two-way conversations, (i.e., times when a party is not talking) that cannot be cost-effectively exploited in FDMA or TDMA systems. DS-SS-CDMA radios can, therefore, accommodate more mobile users than FDMA/TDMA radios

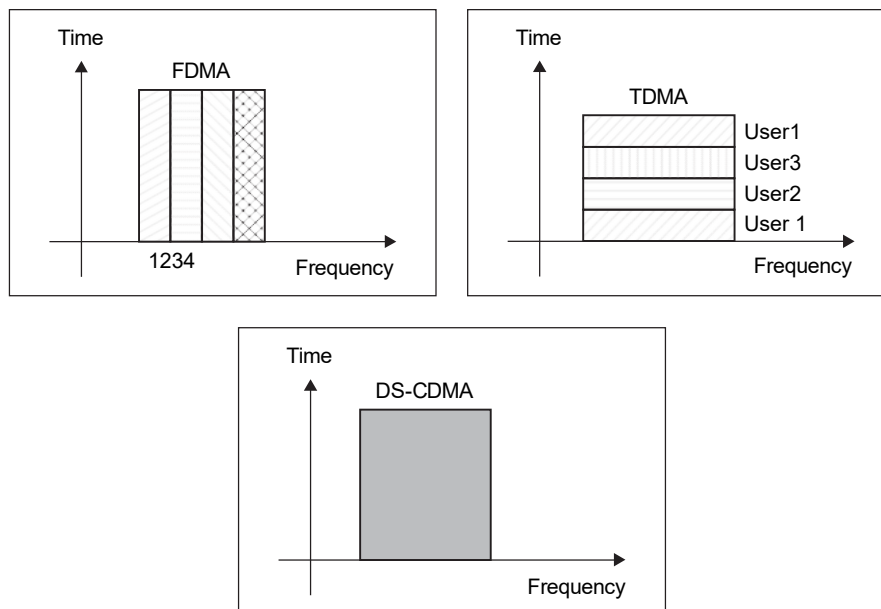


Figure 6.7 Comparison of multiple access methods.

on the same bandwidth. Further capacity gains for FDMA, TDMA, and CDMA can also result from antenna technology advancement by using directional antennas that allow the microcell area to be divided into sectors. Table 6.1 provides a summary of access technologies used for various wireless systems.

Table 6.1 Access technologies for wireless system.

System	Access technology	Mode of operation	Frame rate (kbps)
North American IS-54 (Dual Mode)	TDMA/FDD FDMA/FDD	Digital/ Analog FM	48.6 —
North American IS-95 (Dual Mode)	DS-SS-CDMA/FDD FDMA/FDD	Digital/ Analog FM	1228.8 —
North American IS-136	TDMA/FDD	Digital	48.6
GSM (used all over world)	TDMA/FDD	Digital	270.833
European CT-2 Cordless	FDMA/TDD	Digital	72.0
DECT Cordless	TDMA/TDD	Digital	1152.0

6.6 Capacity of a DS-CDMA System

The capacity of a DS-CDMA system depends on the processing gain, G_p (a ratio of spreading bandwidth, B_w , and information rate, R), the bit energy-to-interference ratio, E_b/I_0 , the voice duty cycle, v_f , the DS-CDMA omnidirectional frequency reuse efficiency, h_f , and the number of sectors, G , in the cell-site antenna.

The received signal power at the cell from a mobile is $S = R \times E_b$. The signal-to-interference ratio is

$$\frac{S}{I} = \frac{R \times E_b}{B_w \times I_0} \quad (6.12)$$

where:

E_b = energy per bit

I_0 = interference density

In a cell with N_u mobile transmitters, the number of effective interferers is $N_u - 1$ because each mobile is an interferer to all other mobiles. This is valid regardless of how the mobiles are distributed within the cell since automatic power control (APC) is used in the mobiles. The APC operates such that the received power at the cell from each mobile is the same as for every other mobile in the cell, regardless of the distance from the center of the cell. APC conserves battery power in the mobiles, minimizes interference to other users, and helps overcome fading.

In a hexagonal cell structure, because of interference from each tier, the S/I ratio is given as (see Chapter 5):

$$\frac{S}{I} = \frac{1}{(N_u - 1) \times [1 + 6 \times k_1 + 12 \times k_2 + 18 \times k_3 + \dots]} \quad (6.13)$$

where:

N_u = number of mobile users in the band, B_w

$k_i, i = 1, 2, 3, \dots$ = the interference contribution from all terminals in individual cells in tiers 1, 2, 3, etc., relative to the interference from the center cell. This loss contribution is a function of both the path loss to the center cell and the power reduction because of power control to an interfering mobile's own cell center.

If we define a frequency reuse efficiency, h_f , as in Equation 6.14a, then E_b/I_0 is given by Equation 6.15.

$$h_f = \frac{1}{[1 + 6 \times k_1 + 12 \times k_2 + 18 \times k_3 + \dots]} \quad (6.14a)$$

$$\frac{S}{I} = \frac{h_f}{(N_u - 1)} \quad (6.14b)$$

$$\frac{E_b}{I_0} = \frac{B_w}{R} \times \frac{h_f}{(N_u - 1)} \quad (6.15)$$

This equation does not include the effect of background thermal and spurious noise (i.e., p) in the spreading bandwidth B_w . Including this as an additive degradation term in the denominator results in a bit energy-to-interference ratio of:

$$\frac{E_b}{I_0} = \frac{B_w}{R} \times \frac{h_f}{(N_u - 1) + p/S} \quad (6.16)$$

Note that from Equation 6.16 the capacity of the DS-CDMA system is reduced by p/S which is the ratio of background thermal plus spurious noise to power level.

For a fixed $G_p = B_w/R$, one way to increase the capacity of the DS-CDMA system is to reduce the required E_b/I_0 , which depends upon the modulation and coding scheme. By using a powerful coding scheme, the E_b/I_0 ratio can be reduced, but this increases system complexity. Also, it is not possible to reduce the E_b/I_0 ratio indefinitely. The only other way to increase the system capacity is to reduce the interference. Two approaches are used: one is based on the natural behavior of human speech and the other is based on the application of the sectorized antennas. From experimental studies it has been found that typically in a full duplex 2-way voice conversation, the duty cycle of each voice is, on the average, less than 40%. Thus, for the remaining period of time the interference induced by the speaker can be eliminated. Since the channel is shared among all the users, noise induced in the desired channel is reduced due to the silent interval of other interfering channels. It is not cost-effective to exploit the voice activity in the FDMA or TDMA system because of the time delay associated with reassigning the channel resource during the speech pauses. If we define v_f as the voice activity factor (< 1), then Equation 6.16 can be written as:

$$\frac{E_b}{I_0} = \frac{h_f}{v_f} \times \frac{B_w}{R} \times \frac{1}{(N_u - 1) + p/S} \quad (6.17a)$$

$$(N_u - 1) + \frac{p}{S} \equiv \left[\frac{h_f}{v_f} \times \left[\frac{B_w}{R} \right] \times \frac{1}{E_b} \right] I_0 \quad (6.17b)$$

The equation to determine the capacity of a DS-SS-CDMA system should also include additional parameters to reflect the bandwidth efficiency factor, the capacity degradation factor due to imperfect power control, and the number of sectors in the cell-site antenna. Equation 6.17b is augmented by these additional factors to provide the following equation for DS-SS-CDMA capacity at one cell:

$$N_u = \frac{h_f h_b c_d \lambda}{v_f} \times \frac{B_w}{R \times (E_b/I_0)} + \frac{1 - \rho}{S} \quad (6.18a)$$

Equation 6.18a can be rewritten as Equation 6.18b by neglecting the last two terms.

$$N_u = \frac{h_f h_b c_d \lambda}{v_f} \times \frac{B_w}{R \times (E_b/I_0)} \quad (6.18b)$$

where:

h_f = frequency reuse efficiency < 1

h_b = bandwidth efficiency factor < 1

c_d = capacity degradation factor to account for imperfect APC < 1

v_f = voice activity factor < 1

B_w = one-way bandwidth of the system
 R = information bit rate plus overhead

E_b = energy per bit of the desired signal

E_b/I_0 = desired energy-to-interference ratio (dependent on quality of service)

λ = efficiency of sector-antenna in cell (< G , number of sectors in the cell-site antenna)

For digital voice transmission, E_b/I_0 is the required value for a bit error rate (BER) of about 10^{-3} or better, and h_f depends on the quality of the diversity. Under the most optimistic assumption, $h_f < 0.5$. The voice activity factor, v_f is usually assumed to be less than or equal to 0.6. E_b/I_0 for a BER of 10^{-3} can be as high as 63 (18 dB) if no coding is used and as low as 5 (7 dB) for a system using a powerful coding scheme. The capacity degradation factor, c_d will depend on the implementation but will always be less than 1.

Example 6.6

Calculate the capacity and spectral efficiency of the DS-SS-CDMA system with an omnidirectional cell using the following data:

- bandwidth efficiency $h_b = 0.9$
- frequency reuse efficiency $h_f = 0.45$

- capacity degradation factor $c_d = 0.8$
- voice activity factor $v_f = 0.4$
- information rate $R = 16.2$ kbps
- $E_b/I_0 = 7$ dB
- one-way system bandwidth $B_w = 12.5$ MHz

Neglect other sources of interference.

Solution

$$E_b/I_0 = 5.02 \text{ (7 dB)}$$

$$N_u = \frac{0.45 \times 0.9 \times 0.8 \times 1 \times 12.5 \times 10^6}{0.4 \times 16.2 \times 10^3 \times 5.02}$$

$$N_u = 124.5 \text{ (say 125)}$$

$$\text{Spectral efficiency, } \eta = \frac{125 \times 16.2}{12.5 \times 10^3} = 0.162 \text{ bits/sec/Hz}$$

In these calculations, an omnidirectional antenna is assumed. If a three-sector antenna (i.e., $G = 3$) is used at a cell site with $\lambda = 2.6$, the capacity will be increased to 325 mobile users per cell, and spectral efficiency will be 0.421 bits/sec/Hz.

6.7 Comparison of DS-CDMA vs. TDMA System Capacity

Using Equations 6.7 and 6.18b with $v_f = 1$ (no voice activity) for TDMA and $\lambda = 1.0$ (omnidirectional cell) for DS-CDMA the ratio of the cell capacity for the DS-CDMA and TDMA systems is given as:

$$\frac{N_{\text{CDMA}}}{N_{\text{TDMA}}} = \frac{c_d N h_f \times 1}{E_b/I_0} \times \frac{1}{v_{f_{\text{cdma}}}} \times \frac{1}{\mu} \times \frac{R_{\text{TDMA}}}{R_{\text{CDMA}}} \quad (6.19)$$

Example 6.7

Using the data given in Examples 6.4 and 6.6, compare the capacity of the DS-CDMA and TDMA omnidirectional cell.

Solution

$$\frac{N_{\text{CDMA}}}{N_{\text{TDMA}}} = \frac{0.8 \times 19 \times 0.45 \times 1}{5.02} \times \frac{1}{0.4} \times \frac{1}{2} \times \frac{16.2}{16.2} = 1.703$$

6.8 Frequency Hopping Spread Spectrum with M-ary Frequency Shift Keying

The FHSS system uses M-ary frequency shift keying modulation (MFSK) and involves the hopping of the carrier frequency in a random manner. It uses MFSK, in which $b = \log_2 M$ information bits determine which one of M frequencies is to be used [19]. The portion of the M-ary signal set is shifted pseudo-randomly by the frequency synthesizer over a hopping bandwidth, B_{ss} . A typical block diagram is shown in Figure 6.8.

In a conventional MFSK system, the data symbol is modulated on a carrier whose frequency is pseudo-randomly determined. The frequency synthesizer produces a transmission tone based on simultaneous dictations of the pseudo-noise (PN) code (see Chapter 11) and the data. At each frequency hop time a PN generator feeds the frequency synthesizer a frequency word (a sequence of L chips), which dictates one of $2L$ symbol-set positions. The FH bandwidth, B_{ss} , and the minimum frequency spacing between consecutive hop positions, Δf , dictate the minimum number of chips required in the frequency word.

Example 6.8

A hopping bandwidth, B_{ss} , of 600 MHz and a frequency step size, Δf , of 400 Hz are used. What is the minimum number of PN chips that are required for each frequency word?

Solution

$$\text{Number of tones contained in } B_{ss} = \frac{B_{ss}}{\Delta f} = \frac{600 \times 10^6}{400} = 1.5 \times 10^6$$

$$\text{Minimum number of chips required} = \lceil \log_2(1.5 \times 10^6) \rceil = 20 \text{ chips}$$

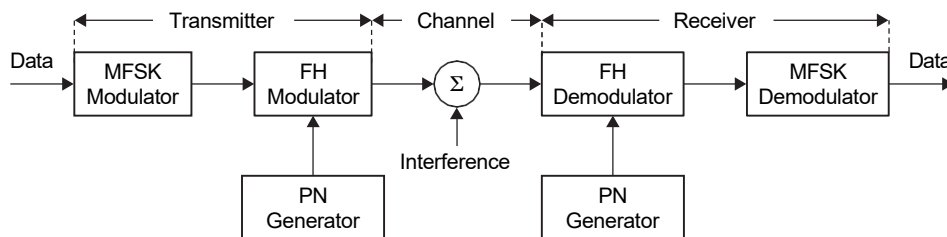


Figure 6.8 Frequency hopping using MFSK.

6.9 Orthogonal Frequency Division Multiplexing (OFDM)

In this section we briefly introduce OFDM. For more details readers should refer to [19]. OFDM uses three transmission principles, multirate, multisymbol, and multicarrier. OFDM is similar to frequency division multiplexing (FDM). OFDM distributes the data over a large number of carriers that are spaced apart at precise frequencies. The spacing provides the orthogonality in this technique, which prevents the demodulator from seeing frequencies other than their own.

Multiple Input, Multiple Output-OFDM (MIMO-OFDM) uses multiple antennas to transmit and receive radio signals. MIMO-OFDM allows service providers to deploy a broadband wireless access system that has non-line-of-sight (NLOS) functionality. MIMO-OFDM takes advantage of the multipath properties of the environment using base station antennas that do not have LOS. The MIMO-OFDM system uses multiple antennas to simultaneously transmit data in small pieces to the receiver, which can process the data flow and put it back together. This process, called *spatial multiplexing*, proportionally boosts the data transmission speed by a factor equal to the number of transmitting antennas. In addition, since all data is transmitted both in the same frequency band and with separate spatial signatures, this technique utilizes spectrum efficiently. VOFDM (vector OFDM) uses the concept of MIMO technology.

We consider a data stream operating at R bps and an available bandwidth of $N\Delta f$ centered at f_c . The entire bandwidth could be used to transmit a data stream, in which case the bit duration would be $1/R$. By splitting the data stream into N substreams using a serial-to-parallel converter, each substream has a data rate of R/N and is transmitted on a separate subcarrier, with spacing between adjacent subcarriers of Δf (see Figure 6.9). The bit duration is N/R . The advantage of OFDM is that on a multiple channel the multipath is reduced relative to the symbol interval by a ratio of $1/N$ and thus imposes less distortion in each modulated symbol. OFDM overcomes inter-symbol interference (ISI) in a multipath environment. ISI has a greater impact at higher data rates because the distance between bits or symbols is smaller. With OFDM, the data rate is reduced by a factor of N , which increases the symbol duration by a factor of N . Thus, if the symbol duration is T_s for the source stream, the duration of OFDM signals is NT_s . This significantly reduces the effect of ISI. As a design criterion, N is selected so that NT_s is significantly greater than τ_{rms} (rms delay spread) of the channel. With the use of OFDM, it may not be necessary to deploy an equalizer. OFDM is an ideal solution for broadband communications, because increasing the data rate is simply a matter of increasing the number of subcarriers. To avoid overlap between consecutive symbols, a time guard is enforced between the transmissions of two OFDM pulses that will reduce the effective data rate. Also, some subcarriers are devoted to synchronization of signal, and some are reserved for redundancy.

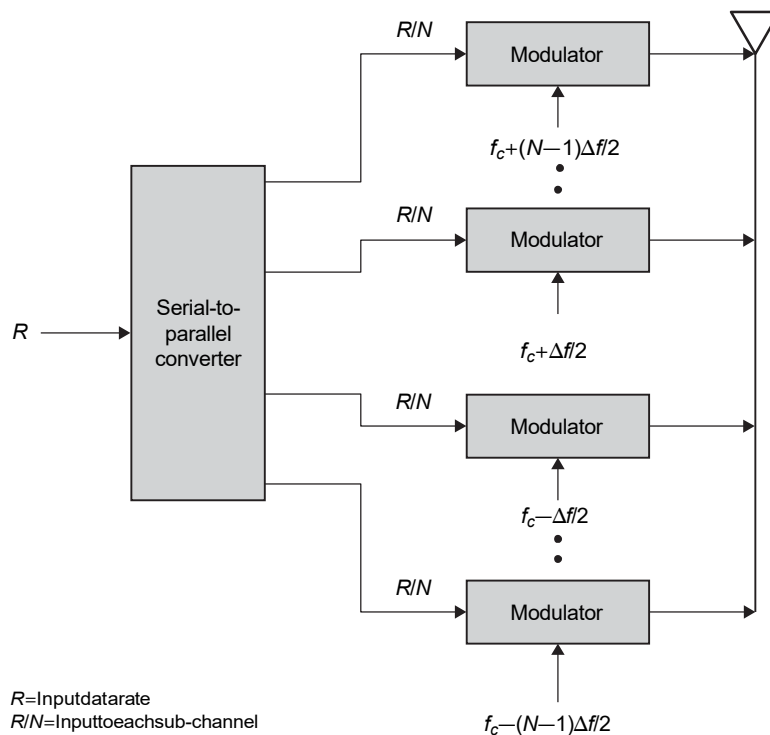


Figure 6.9 Orthogonal frequency division multiplexing (OFDM).

The most important feature of OFDM is the orthogonal relationship between the subcarrier signals. Orthogonality allows the OFDM subcarriers to overlap each other without interference. OFDM uses FH to create a spread spectrum system. FH has several advantages over DSSS, for example, no near-far problem, easier synchronization, less complex receivers, and so on.

In the OFDM the input information sequence is first converted into parallel data sequences and each serial/parallel converter output is multiplied with spreading code. Data from all subcarriers is modulated in baseband by inverse fast Fourier transform (IFFT) and converted back into serial data. The guard interval is inserted between symbols to avoid ISI caused by multipath fading and finally the signal is transmitted after RF up-conversion. At the receiver, after down-conversion, the m -subcarrier component corresponding to the received data is first coherently detected with FFT and then multiplied with gain to combine the energy of the received signals scattered in the frequency domain (see Figure 6.10). Wireless Local Area Networks (WLAN) development is ongoing for wireless point-to-point and point-to-multipoint configurations using OFDM technology.

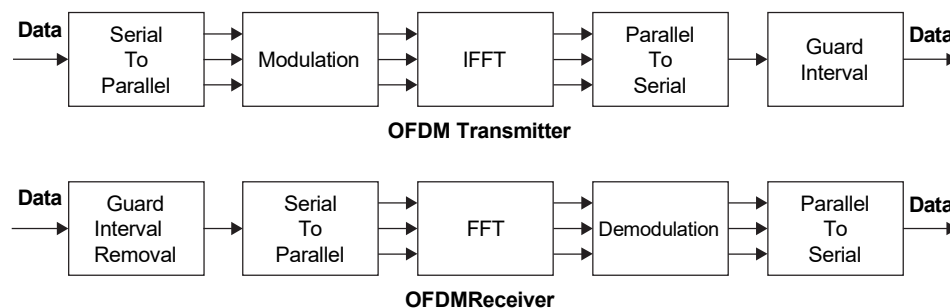


Figure 6.10 IEEE 802.11a Transmit and Receive OFDM.

In a supplement to the IEEE 802.11 standard, the IEEE 802.11 working group published IEEE 802.11a, which outlines the use of OFDM in the 5.8 GHz band.

The basic principle of operation is to divide a high-speed binary signal to be transmitted into a number of lower data rate subcarriers. There are 48 data subcarriers and 4 pilot subcarriers for a total of 52 subcarriers. Each lower data rate bit stream is used to modulate a separate subcarrier from one of the channels in the 5 GHz band. Prior to transmission the data is encoded using convolutional code (see Chapter 8) of rate, $R=1/2$ and bit interleaved for the desired data rate. Each bit is then mapped into a complex number according to the modulation type and subdivided in 48 data subcarriers and 4 pilot subcarriers. The subcarriers are combined using an IFFT and transmitted. At the receiver, the carrier is converted back to a multicarrier lower data rate form using FFT. The lower data subcarriers are recombined to form a high rate data unit.

6.10 Multicarrier DS-SS-CDMA (MC-SS-CDMA)

Future wireless systems such as a fourth-generation (4G) system will need flexibility to provide subscribers with a variety of services such as voice, data, images, and video. Because these services have widely differing data rates and traffic profiles, future generation systems will have to accommodate a wide variety of data rates. DS-SS-CDMA has proven very successful for large-scale cellular voice systems, but there are concerns whether DS-SS-CDMA will be well-suited to non-voice traffic. The DS-SS-CDMA system suffers inter-symbol interference (ISI) and multi-user interference (MUI) caused by multipath propagation, leading to a high loss of performance.

With OFDM, the time dispersive channel is seen in the frequency domain as a set of parallel independent flat subchannels and can be equalized at a low complexity. There are potential benefits to combining OFDM and DS-SS-CDMA. Basically the frequency-selective channel is first equalized in the frequency domain using the

OFDM modulation technique. DS-CDMA is applied on top of the equalized channel, keeping the orthogonality properties of spreading codes. The combination of OFDM and DS-CDMA is used in MC-DS-CDMA. MC-DS-CDMA [4,5,12,25] marries the best of the OFDM and DS-CDMA world and, consequently, it can ensure good performance in severe multipath conditions. MC-DS-CDMA can achieve very large average throughput. To further enhance the spectral efficiency of the system, some form of adaptive modulation can be used.

Basically, three main design exist in the literature, namely, MC-CDMA, MC-DS-CDMA, and multi-tone (MT)-CDMA. In MC-CDMA, the spreading code is applied across a number of orthogonal subcarriers in the frequency domain. In MC-DS-CDMA, the data stream is first divided into a number of substreams. Each substream is spread in time through a spreading code and then transmitted over one of a set of orthogonal subcarriers. In MT-CDMA the system undergoes similar operations as MC-DS-CDMA except that the different subcarriers are not orthogonal after spreading. This allows higher spectral efficiencies and longer spreading codes; however, different substreams interfere with one another. The MC-DS-CDMA transmitter spreads the original data stream over different orthogonal subcarriers using a given spreading code in the frequency domain.

6.11 Random Access Methods

So far we have discussed the reservation-based schemes, now we focus on random-access schemes [8]. When each user has a steady flow of information to transmit (for example, a data file transfer or a facsimile transmission), reservation-based access methods are useful as they make an efficient use of communication resources. However, when the information to be transmitted is bursty in nature, the reservation-based access methods result in wasting communication resources. Furthermore, in a cellular system where subscribers are recharged based on a channel connection time, the reservation-based access methods may be too expensive to transmit short messages. Random-access protocols provide flexible and efficient methods for managing a channel access to transmit short messages. The random-access methods give freedom for each user to gain access to the network whenever the user has information to send. Because of this freedom, these schemes result in contention among users accessing the network. Contention may cause collisions and may require retransmission of the information. The commonly used random-access protocols are pure ALOHA, slotted-ALOHA, and CSMA/CD. In the following section we briefly describe details of each of these protocols and provide the necessary throughput expressions.

6.11.1 Pure ALOHA

In the pure ALOHA [18,23] scheme, each user transmits information whenever the user has information to send. A user sends information in packets. After

sending a packet, the user waits a length of time equal to the round-trip delay for an acknowledgment (ACK) of the packet from the receiver. If no ACK is received, the packet is assumed to be lost in a collision and it is retransmitted with a randomly selected delay to avoid repeated collisions.* The normalized throughput S (average new packet arrival rate divided by the maximum packet throughput) of the pure ALOHA protocol is given as:

$$S = Ge^{-2G} \quad (6.20)$$

where G = normalized offered traffic load

From Equation 6.20 it should be noted that the maximum throughput occurs at traffic load $G = 0.5$ and is $S = 1/2e$. This is about 0.184. Thus, the best channel utilization with the pure ALOHA protocol is only 18.4%.

6.11.2 Slotted ALOHA

In the slotted-ALOHA [23] system, the transmission time is divided into time slots. Each time slot is made exactly equal to packet transmission time. Users are synchronized to the time slots, so that whenever a user has a packet to send, the packet is held and transmitted in the next time slot. With the synchronized time slots scheme, the interval of a possible collision for any packet is reduced to one packet time from two packet times, as in the pure ALOHA scheme. The normalized throughput S for the slotted-ALOHA protocol is given as:

$$S = Ge^{-G} \quad (6.21)$$

where G = normalized offered traffic load

The maximum throughput for the slotted ALOHA occurs at $G = 1.0$ (Equation 6.21) and it is equal to $1/e$ or about 0.368. This implies that at the maximum throughput, 36.8% of the time slots carry successfully transmitted packets. The best channel utilization with the slotted ALOHA protocol is 36.8%—twice the pure ALOHA protocol.

*It should be noted that the protocol on CDMA access channels as implemented in TIA IS-95-A is based upon the pure ALOHA approach. The mobile station randomizes its attempt for sending a message on the access channel and may retry if an acknowledgment is not received from the base station. For further details, one should reference Section 6.6.3.1.1.1 of TIA IS-95-A.

6.11.3 Carrier Sense Multiple Access (CSMA)

The carrier sense multiple access (CSMA) [8,18] protocols have been widely used in both wired and wireless LANs. These protocols provide enhancements over the pure and slotted ALOHA protocols. The enhancements are achieved through the use of the additional capability at each user station to sense the transmissions of other user stations. The carrier sense information is used to minimize the length of collision intervals. For carrier sensing to be effective, propagation delays must be less than packet transmission times. Two general classes of CSMA protocols are nonpersistent and p-persistent.

- **Nonpersistent CSMA:** A user station does not sense the channel continuously while it is busy. Instead, after sensing the busy condition, it waits for a randomly selected interval of time before sensing again. The algorithm works as follows: if the channel is found to be idle, the packet is transmitted; or if the channel is sensed busy, the user station backs off to reschedule the packet to a later time. After backing off, the channel is sensed again, and the algorithm is repeated again.
- **p-persistent CSMA:** The slot length is typically selected to be the maximum propagation delay. When a station has information to transmit, it senses the channel. If the channel is found to be idle, it transmits with probability p . With probability $q = 1 - p$, the user station postpones its action to the next slot, where it senses the channel again. If that slot is idle, the station transmits with probability p or postpones again with probability q . The procedure is repeated until either the frame has been transmitted or the channel is found to be busy. If the station initially senses the channel to be busy, it simply waits one slot and applies the above procedure.
- **1-persistent CSMA:** 1-persistent CSMA is the simplest form of the p-persistent CSMA. It signifies the transmission strategy, which is to transmit with probability 1 as soon as the channel becomes idle. After sending the packet, the user station waits for an ACK, and if it is not received within a specified amount of time, the user station waits for a random amount of time, and then resumes listening to the channel. When the channel is again found to be idle, the packet is retransmitted immediately.

For more details, the reader should refer to [18].

The throughput expressions for the CSMA protocols are:

- **Unslotted nonpersistent CSMA**

$$S = \frac{G e^{-aG}}{aG(1+2a) + e^{-aG}} \quad (6.22)$$

- Slotted nonpersistent CSMA

$$S = \frac{aGe^{-aG}}{1 - e^{-aG} + a} \quad (6.23)$$

- Unslotted 1-persistent CSMA

$$S = \frac{G[1 + G + aG(1 + G + (aG)/2)]e^{-G(1+2a)}}{G(1+2a) - (1 - e^{-aG}) + (1 + aG)e^{-G(1+a)}} \quad (6.24)$$

- Slotted 1-persistent CSMA

$$S = \frac{Ge^{-G(1+a)}[1 + a - e^{-aG}]}{(1 + a)(1 - e^{-aG}) + aeG(1+a)} \quad (6.25)$$

where:

S = normalized throughput

G = normalized offered traffic load

$a = \tau / T_p$

τ = maximum propagation delay

T_p = packet transmission time

Example 6.9

We consider a WLAN installation in which the maximum propagation delay is 0.4 sec. The WLAN operates at a data rate of 10 Mbps, and packets have 400 bits. Calculate the normalized throughput with: (1) an unslotted nonpersistent, (2) a slotted persistent, and (3) a slotted 1-persistent CSMA protocol.

Solution

$$T_p = \frac{400}{10} = 40 \mu\text{s}$$

$$a = \frac{\tau}{T_p} = \frac{0.4}{40} = 0.01$$

$$G = \frac{40 \times 10^{-6} \times 10 \times 10^6}{400} = 1$$

- Slotted nonpersistent:

$$S = \frac{0.01 \times 1 \times e^{-0.01}}{1 - e^{-0.01} + 0.01} = 0.495$$

- Unslotted nonpersistent:

$$S = \frac{1 \times e^{-0.01}}{(1 + 0.02) + e^{-0.01}} = 0.493$$

- Slotted 1-persistent:

$$S = \frac{e^{-1.01}(1 + 0.01 - e^{-0.01})}{(1 + 0.01)(1 - e^{-0.01}) + 0.01e^{-1.01}} = 0.531$$

6.11.4 Carrier Sense Multiple Access with Collision Detection

A considerable performance improvement in the basic CSMA protocol can be achieved by means of the *carrier sense multiple access with collision detection* (CSMA/CD) technique. The CSMA/CD protocols are essentially the same as those for CSMA with addition of the collision-detection feature. Similar to CSMA protocols, there are nonpersistent, 1-persistent, and p-persistent CSMA/CD protocols. More details about CSMA/CD protocols can be found in [27].

When a CSMA/CD station senses that a collision has occurred, it immediately stops transmitting its packets and sends a brief jamming signal to notify all stations of this collision. Collisions are detected by monitoring the analog waveform directly from the channel. When signals from two or more stations are presented simultaneously, the composite waveform is distorted from that of a single station. This is manifested in the form of larger than normal voltage amplitude on the cable. In the Ethernet the collision is recognized by the transmitting station, which goes into a retransmission phase based on an exponential random backoff algorithm.

The normalized throughput for unslotted nonpersistent and slotted nonpersistent CSMA/CD is given as:

Unslotted nonpersistent CSMA/CD

$$S = \frac{Ge^{-aG}}{Ge^{-aG} + bG(1 - e^{-aG}) + 2aG(1 - e^{-aG}) + (2 - e^{-aG})} \quad (6.26)$$

where b = jamming signal length

Slotted nonpersistent CSMA/CD

$$S = \frac{aGe^{-aG}}{aGe^{-aG} + b(1 - e^{-aG}) + a(2 - e^{-aG})} \quad (6.27)$$

While these collision detection mechanisms are good ideas on a wired local area network (LAN), they cannot be used on a wireless local area network (WLAN) environment for two main reasons:

- Implementing a collision detection mechanism would require the implementation of a full duplex radio capable of transmitting and receiving at the same time—an approach that would increase the cost significantly.
- In a wireless environment we cannot assume that all stations hear each other (which is the basic assumption of the collision detection scheme), and the fact that a station wants to transmit and sense the medium as free does not necessarily mean that the medium is free around the receiver area.

6.11.5 Carrier Sense Multiple Access with Collision Avoidance (CSMA/CA)

IEEE 802.11 uses a protocol known as *carrier sense multiple access with collision avoidance* (CSMA/CA) or distributed coordination function (DCF). CSMA/CA attempts to avoid collisions by using *explicit packet acknowledgment* (ACK), which means an ACK packet is sent by the receiving station to confirm that the data packet arrived intact.

The CSMA/CA protocol works as follows. A station wishing to transmit senses the medium, if the medium is busy (i.e., some other station is transmitting) then the station defers its transmission to a later time. If no activity is detected, the station waits an additional, randomly selected period of time and then transmits if the medium is still free. If the packet is received intact, the receiving station issues an ACK frame that, once successfully received by the sender, completes the process. If the ACK frame is not detected by the sending station, either because the original data packet was not received intact or the ACK was not received intact, a collision is assumed to have occurred and the data packet is transmitted again after waiting another random amount of time. The CSMA/CA provides a way to share access over the medium. This explicit ACK mechanism also handles interference and other radio-related problems very effectively. However, it does add some overhead to 802.11 that 802.3 does not have, so that an 802.11 WLAN will always have slower performance than the equivalent Ethernet LAN (802.3).

The CSMA/CA protocol is very effective when the medium is not heavily loaded since it allows stations to transmit with minimum delay. But there is always a chance of stations simultaneously sensing the medium as being free

and transmitting at the same time, causing a collision. These collisions must be identified, so that the media access control (MAC) layer can retransmit the packet by itself and not by the upper layers, which would cause significant delay. In particular, the hidden node and exposed node problems should be addressed by MAC. Both of them give rise to many performance problems including throughput degradation, unfair throughput distribution, and throughput instability (see Chapter 18 for details).

The IEEE 802.11 uses a collision avoidance (CA) mechanism together with a positive ACK. The MAC layer of a station wishing to transmit senses the medium. If the medium is free for a specified time (called *distributed inter-frame space* (DIFS)), then the station is able to transmit the packet; if the medium is busy (or becomes busy during the DIFS interval) the station defers using the *exponential backoff algorithm*.

This scheme implies that, except in cases of very high network congestion, no packets will be lost, because retransmission occurs each time a packet is not acknowledged. This entails that all packets sent will reach their destination in sequence.

The IEEE 802.11 MAC layer provides cyclic redundancy check (CRC) *checksum* and *packet fragmentation*. Each packet has a CRC checksum calculated and attached to ensure that the data was not corrupted in transmit. Packet fragmentation is used to segment large packets into smaller units when sent over the medium. This is useful in very congested environments or when interference is a factor, since large packets have a better chance of being corrupted. This technique reduces the need for retransmission in many cases and improves overall wireless network performance. The MAC layer is responsible for reassembling fragments received, rendering the process transparent to higher-level protocols. The following are some of the reasons it is preferable to use smaller packets in a WLAN environment.

- Due to higher BER of a radio link, the probability of a packet getting corrupted increases with packet size.
- In case of packet corruption (either due to collision or interference), the smaller the packet, the less overhead it needs to retransmit.

A simple *stop-and-wait* algorithm is used at the MAC sublayer. In this mechanism the transmitting station is not allowed to transmit a new fragment until one of the following happens:

- Receives an ACK for the fragment, or
- Decides that the fragment was retransmitted too many times and drops the whole frame.

Exponential backoff scheme is used to resolve contention problems among different stations wishing to transmit data at the same time. When a station goes into the backoff state, it waits an additional, randomly selected number of time slots known as a contention window (in 802.11b a slot has a $20\mu\text{s}$ duration and the random number must be greater than 0 and smaller than a maximum value referred to as a contention window (CW)). During the wait, the station continues sensing the medium to check whether it remains free or if another transmission begins. At the end of its window, if the medium is still free the station can send its frame. If during the window another station begins transmitting data, the backoff counter is frozen and counting down starts again as the channel returns to the idle state.

There is a problem related to the CW dimension. With a small CW, if many stations attempt to transmit data at the same time it is possible that some of them may have the same backoff interval. This means that there will continuously be collisions, with serious effects on network performance. On the other hand, with a large CW, if a few stations wish to transmit data they will likely have long back-off delays resulting in degradation of network performance. The solution is to use an exponentially growing CW size. It starts from a small value ($CW_{\min}=31$) and doubles after each collision, until it reaches the maximum value CW_{\max} ($CW_{\max}=1023$). In 802.11 the backoff algorithm is executed in three cases:

1. When the station senses the medium busy before the first transmission of a packet
2. Before each retransmission
3. After a successful transmission

This is necessary to avoid a single host wanting to transmit a large quantity of data and occupying the channel for too long, denying access to all other stations. The backoff mechanism is not used when the station decides to transmit

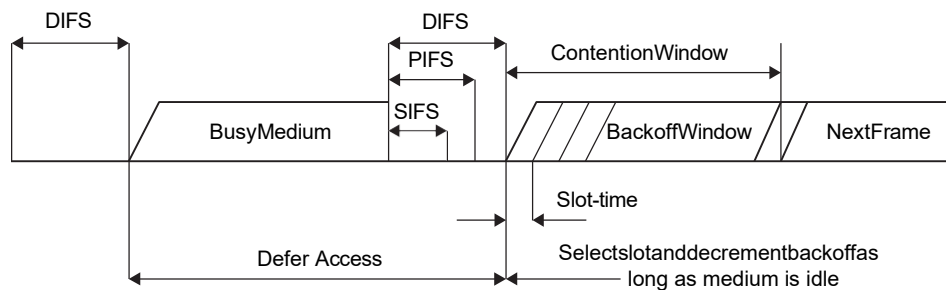


Figure 6.11 CSMA/CA in IEEE 802.11b.

a new packet after an idle period and the medium has been free for more than a distributed inter-frame space (DIFS) (see Figure 6.11).

To support time-bounded services, the IEEE 802.11 standard defines the point coordinate function (PCF) to let stations have priority access to the wireless medium, coordinated by a station called point coordinate (PC). The PCF has higher priority than DIFS, because it may start transmissions after a shorter duration than DIFS; this time space is called PCF interframe space (PIFS), which is $25\mu\text{s}$ for IEEE 802.11 and larger than SIFS.

$$\text{The transmission time for a data frame} = \left(\text{PLCP} + \frac{D}{R} \right) \mu\text{s}$$

where:

PLCP = the time required to transmit the physical layer convergence protocol (PLCP)

D = the frame size

R = the channel bitrate

$$\text{CSMA/CA packet transmission time} = \text{BO} + \text{DIFS} + 2\text{PLCP} + \frac{D}{R} + \text{SIFS} + \frac{A}{R} \mu\text{s}$$

where:

A = the ACK frame size

BO = the backoff time

DIFS = the distributed inter-frame space

SIFS = the short inter-frame space

PIFS = the point coordinate interframe space

6.12 Idle Signal Casting Multiple Access

In the CSMA scheme, each terminal must be able to detect the transmissions of all other terminals. However, not all packets transmitted from different terminals can be sensed, or terminals may be hidden from each other by buildings or some other obstacles. This is known as the *hidden terminal problem*, which severely degrades the throughput of the CSMA. The idle signal casting multiple access (ISMA) system transmits an idle/busy signal from the base station to indicate the presence or absence of another terminal's transmission. The ISMA and CSMA are basically the same. In the CSMA, each terminal must listen to all other terminals, whereas in the ISMA, each terminal is informed from the base station of the other terminals' transmission. Similar to CSMA, there are non-persistent ISMA and 1-persistent ISMA.

6.13 Packet Reservation Multiple Access

Packet reservation multiple access (PRMA) allows a variety of information sources to share the same communication channel and obtains a statistical multiplexing

effect. In PRMA, time is divided into frames, each of which consists of a fixed number of time slots. For voice terminals, voice activity detection is adopted. The voice signal comprises a sequence of talk spurts. At the beginning of a talk spurt, the terminal transmits the first packet based on slotted ALOHA. Once the packet is transmitted successfully, that terminal is allowed to use the same time slots in the succeeding frames (reservation is made). The reservation is kept until the end of the talk spurt. The status “reserved” or “unreserved” of each slot is broadcast from the base station.

6.14 Error Control Schemes for Link Layer

Error control schemes for the link layer are used to improve the performance of mobile communication systems [8]. Several automatic repeat request (ARQ) schemes are used. At the physical layer of wireless mobile communication systems, error detection and correction techniques such as forward error correction (FEC) schemes are used. For some of the data services, higher layer protocols use ARQ schemes to enable retransmission of any data frames in which an error is detected. The ARQ schemes are classified as follows [23]:

- **Stop and Wait:** The sender transmits the first packet numbered 0 after storing a copy of that packet. The sender then waits for an ACK numbered 0, (ACK0) of that packet. If the ACK0 does not arrive before a time-out, the sender transmits another copy of the first packet. If the ACK0 arrives before a time-out, the sender discards the copy of the first packet and is ready to transmit the next packet, which it numbers 1. The sender repeats the previous steps, with numbers 0 and 1 interchanged. The advantages of the Stop and Wait protocol are its simplicity and its small buffer requirements. The sender needs to keep only the copy of the packet that it last transmitted, and the receiver does not need to buffer packets at the data link layer. The main disadvantage of the Stop and Wait protocol is that it does not use the communication link very efficiently.

The total time taken to transmit a packet and to prepare for transmitting the next one is

$$T = T_p + 2T_{\text{prop}} + 2T_{\text{proc}} + T_a \quad (6.28)$$

where:

T = total time for transmission time

T_p = transmission time for a packet

T_{prop} = propagation time of a packet or an ACK

T_{proc} = processing time for a packet or an ACK

T_a = transmission time for an ACK

The protocol efficiency without any error is:

$$h(0) = \frac{T_p}{T} \quad (6.29)$$

If p is the probability that a packet or its ACK is corrupted by transmission errors, and a successful transmission of a packet and its ACK takes T seconds and occurs with probability $1 - p$, the protocol efficiency for full duplex (FD) is given as:

$$h_{FD} = \frac{(1-p)T_p}{(1-p)T + pT_p} \quad (6.30)$$

- **Selective Repeat Protocol (SRP):** In case of the SRP, only the selected packets are retransmitted. The data link layer in the receiver delivers exactly one copy of every packet in the correct order. The data link layer in the receiver may get the packets in the wrong order from the physical layer. This occurs, for example, when transmission errors corrupt the first packet and not the second one. The second packet arrives correctly at the receiver before the first. The data link layer in the receiver uses a buffer to store the packets that arrive out of order. Once the data link layer in the receiver has a consecutive group of packets in its buffer, it can deliver them to the network layer. The sender also uses a buffer to store copies of the unacknowledged packets. The number of the packets which can be held in the sender/receiver buffer is a design parameter.

Let W = the number of packets which the sender and receiver buffers can each hold and $SRP =$ number of packets in modulo $2W$. The protocol efficiency without any error and with a packet error probability of p is given as:

$$h(0) = \min \left\{ \frac{WT_p}{T}, 1 \right\} \quad (6.31)$$

For very large W , the protocol efficiency is

$$h(p) = 1 - p \quad (6.32)$$

where:

$WT_p =$ time-out

$$h(p) = \frac{2 + p(W-1)}{2 + p(3W-1)} \quad (6.33)$$

SRP is very efficient, but it requires buffering packets at both the sender and the receiver.

- **Go-Back-N (GBN):** The Go-Back-N protocol allows the sender to have multiple unacknowledged packets without the receiver having to store packets. This is done by not allowing the receiver to accept packets that are out of order. When a time-out timer expires for a packet, the transmitter resends that packet and all subsequent packets. The Go-Back-N protocol improves on the efficiency of the Stop and Wait protocol, but it is less efficient than SRP. The protocol efficiency for full duplex is given as:

$$h_{FD} = \frac{1}{1 + \left(\frac{p}{1-p}\right)W} \quad (6.34)$$

- **Window-control Operation Based on Reception Memory (WORM) ARQ:** In digital cellular systems, bursty errors occur by multipath fading, shadowing, and handoffs. The typical bit-error rate fluctuates from 10^{-1} to 10^{-6} . Therefore, the conventional ARQ schemes do not operate well in a digital cellular system. WORM ARQ has been suggested for control of dynamic error characteristics. It is a hybrid scheme that combines SRP and GBN protocol. GBN protocol is chosen in these severe error conditions whereas SRP is selected in the normal error condition.
- **Variable Window and Frame Size GBN and SRP [24]:** Since wireless systems have bursty error characteristics, the error control schemes should have a dynamic adaptation to a bursty channel environment. The SRP and GBN with variable window and frame size have been proposed to improve error control in wireless systems. Table 6.2 provides the window and frame size for different BER. If the error rate increases, the window and frame size are decreased. In the case of the error rate being small, the window and frame size are increased. The optimum threshold values of BER, window and frame size were obtained through computer simulation.

Table 6.2 Bit-error rate versus window and size.

Bit-error rate (BER)	Window size (W)	Frame size (bits)
$BER \leq 10^{-4}$	32	172
$10^{-4} < BER < 10^{-3}$	8	80
$10^{-3} < BER < 10^{-2}$	4	40
$10^{-2} < BER$	2	16

Example 6.10

We consider a WLAN in which the maximum propagation delay is 4 sec. The WLAN operates at a data rate of 10 Mbps. The data and ACK packets are of 400 and 20 bits, respectively. The processing time for a data or ACK packet is 1 sec. If the probability p that a data packet or its ACK can be corrupted during transmission is 0.01, find the data link protocol efficiency with (1) Stop and Wait protocol—full duplex, (2) SRP with window size $W=8$, and (3) Go-Back-N protocol with window size $W=8$.

Solution

$$T_p = \frac{400}{10} = 40 \mu\text{s}$$

$$T_a = \frac{20}{10} = 2 \mu\text{s}$$

$$T_{\text{prop}} = 4 \mu\text{s}$$

$$T_{\text{proc}} = 1 \mu\text{s}$$

$$T = 40 + 2 \times 4 + 2 \times 1 + 2 = 52 \mu\text{s}$$

Stop and Wait:

$$h = \frac{(1-0.01) \times 40}{(1-0.01) \times 52 + 0.01 \times 40} = 0.763$$

SRP:

$$h = \frac{2 + 0.01(8-1)}{2 + 0.01(24-1)} = 0.954$$

GBN:

$$h = \frac{1}{1 + 8 \left(\frac{0.01}{1-0.01} \right)} = 0.925$$

6.15 Summary

The chapter described the access technologies used in wireless communications including reservation-based multiple access and random multiple access. FDMA,

added CDMA technology to become 3G systems. Cordless telephone systems started with CT0 and CT1, became digital with CT2, and ended in Europe in the fully digital standard **DECT**. This standard has even been chosen as one of the candidates for a 3G system (IMT-FT).

While the number of different systems might be confusing, there are some “natural” development paths. Most network providers offering GSM service today will deploy UMTS, while cdmaOne users will choose cdma2000 for simpler migration. The reasons for this are quite simple. With the introduction of GPRS in GSM networks, the core of the network was already enhanced in a way that it can be directly used for UMTS with the radio technologies **UTRA FDD** and **UTRA TDD**. A similar path for evolution exists for **TD-SCDMA**, the Chinese proposal for a 3G system (which has been integrated into UTRA TDD). With some simplification it can be said that UMTS mainly adds a new radio interface but relies in its initial phase on the same core network as GSM/GPRS. Also for cdmaOne the evolution to cdma2000 technologies is quite natural, as the new standard is backward compatible and can reuse frequencies. Cdma2000 1x still uses the same 1.25 MHz channels as cdmaOne does, but offers data rates of up to 153 kbit/s. The **cdma2000 3x** standard uses three 1.25 MHz channels to fit into ITU’s frequency scheme for 3G. However, this standard is not pushed as much as the following enhancements of cdma2000 1x. These enhancements are:

- **cdma2000 1x EV-DO** (evolution-data optimized, also known as high data rate (HDR), some call it data only) promising peak data rates of 2.4 Mbit/s using a second 1.25 MHz channel; and
- **cdma2000 1x EV-DV** (evolution-data and voice) aiming at 1.2 Mbit/s for mobile and 5.2 Mbit/s for stationary users.

Cdma2000 1x EV-DO was the first version of cdma2000 accepted by the ITU as 3G system. More information about the technologies and acronyms used in the diagram is provided in the following sections.

4.1 GSM

GSM is the most successful digital mobile telecommunication system in the world today. It is used by over 800 million people in more than 190 countries. In the early 1980s, Europe had numerous coexisting analog mobile phone systems, which were often based on similar standards (e.g., NMT 450), but ran on slightly different carrier frequencies. To avoid this situation for a second generation fully digital system, the **groupe spéciale mobile (GSM)** was founded in 1982. This system was soon named the **global system for mobile communications (GSM)**, with the specification process lying in the hands of ETSI (ETSI, 2002), (GSM Association, 2002). In the context of UMTS and the creation of 3GPP (Third generation partnership project, 3GPP, 2002a) the whole development process of GSM was transferred to 3GPP and further development is combined with 3G development. 3GPP assigned new numbers to all GSM stan-

dards. However, to remain consistent with most of the GSM literature, this GSM section stays with the original numbering (see 3GPP, 2002a, for conversion). Section 4.4 will present the ongoing joint specification process in more detail.

The primary goal of GSM was to provide a mobile phone system that allows users to roam throughout Europe and provides voice services compatible to ISDN and other PSTN systems. The specification for the initial system already covers more than 5,000 pages; new services, in particular data services, now add even more specification details. Readers familiar with the ISDN reference model will recognize many similar acronyms, reference points, and interfaces. GSM standardization aims at adopting as much as possible.

GSM is a typical second generation system, replacing the first generation analog systems, but not offering the high worldwide data rates that the third generation systems, such as UMTS, are promising. GSM has initially been deployed in Europe using 890–915 MHz for uplinks and 935–960 MHz for downlinks – this system is now also called **GSM 900** to distinguish it from the later versions. These versions comprise GSM at 1800 MHz (1710–1785 MHz uplink, 1805–1880 MHz downlink), also called **DCS (digital cellular system) 1800**, and the GSM system mainly used in the US at 1900 MHz (1850–1910 MHz uplink, 1930–1990 MHz downlink), also called **PCS (personal communications service) 1900**. Two more versions of GSM exist. **GSM 400** is a proposal to deploy GSM at 450.4–457.6/478.8–486 MHz for uplinks and 460.4–467.6/488.8–496 MHz for downlinks. This system could replace analog systems in sparsely populated areas.

A GSM system that has been introduced in several European countries for railroad systems is **GSM-Rail** (GSM-R, 2002), (ETSI, 2002). This system does not only use separate frequencies but offers many additional services which are unavailable using the public GSM system. GSM-R offers 19 exclusive channels for railroad operators for voice and data traffic (see section 4.1.3 for more information about channels). Special features of this system are, e.g., emergency calls with acknowledgements, voice group call service (VGCS), voice broadcast service (VBS). These so-called advanced speech call items (ASCI) resemble features typically available in trunked radio systems only (see section 4.3). Calls are prioritized: high priority calls pre-empt low priority calls. Calls have very short set-up times: emergency calls less than 2 s, group calls less than 5 s. Calls can be directed for example, to all users at a certain location, all users with a certain function, or all users within a certain number space. However, the most sophisticated use of GSM-R is the control of trains, switches, gates, and signals.

Trains going not faster than 160 km/h can control all gates, switches, and signals themselves. If the train goes faster than 160 km/h (many trains are already capable of going faster than 300 km/h) GSM-R can still be used to maintain control. The following section describes the architecture, services, and protocols of GSM that are common to all three major solutions, **GSM 900**, **GSM 1800**, and **GSM 1900**.

GSM has mainly been designed for this and voice services and this still constitutes the main use of GSM systems. However, one can foresee that many future applications for mobile communications will be data driven. The relationship of data to voice traffic will shift more and more towards data.

4.1.1 Mobile services

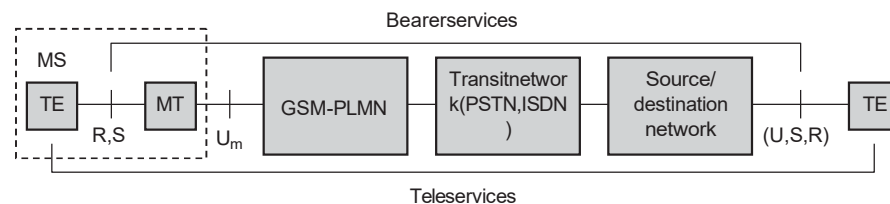
GSM permits the integration of different voice and data services and the interworking with existing networks. Services make a network interesting for customers. GSM has defined three different categories of services: bearer, tele, and supplementary services. These are described in the following subsections. Figure 4.3 shows a reference model for GSM services. A **mobile station MS** is connected to the **GSM public land mobile network (PLMN)** via the U_m interface. (GSM-PLMN is the infrastructure needed for the GSM network.) This network is connected to transit networks, e.g., **integrated services digital network (ISDN)** or traditional **public switched telephone network (PSTN)**. There might be an additional network, the source/destination network, before another **terminal TE** is connected. **Bearer services** now comprise all services that enable the transparent transmission of data between the interfaces to the network, i.e., in case of the mobile station, and a similar interface for the other terminal (e.g., S_0 for ISDN terminals). Interfaces like U, S, and R in case of ISDN have not been defined for all networks, so it depends on the specific network which interface is used as a reference for the transparent transmission of data. In the classical GSM model, bearer services are connection-oriented and circuit- or packet-switched. These services only need the lower three layers of the ISO/OSI reference model.

Within the mobile station MS, the **mobile termination (MT)** performs all network-specific tasks (TDMA, FDMA, coding etc.) and offers an interface for data transmission (S) to the terminal TE which can then be network independent. Depending on the capabilities of TE, further interfaces may be needed, such as R, according to the ISDN reference model (Halsall, 1996). **Tele services** are application specific and may thus need all seven layers of the ISO/OSI reference model. These services are specified end-to-end, i.e., from one terminal TE to another.

4.1.1.1 Bearer services

GSM specifies different mechanisms for data transmission, the original GSM allowing for data rates of up to 9600 bit/s for non-voice services. Bearer services permit transparent and non-transparent, synchronous or asynchronous data transmission. **Transparent bearer services** only use the functions of the physical layer (layer 1) to transmit data. Data transmission has a constant delay and throughput if no transmission errors occur. The only mechanism to increase

Figure 4.3
Bearer and tele services reference model



transmission quality is the use of **forward error correction (FEC)**, which codes redundancy into the data stream and helps to reconstruct the original data in case of transmission errors. Depending on the FEC, data rates of 2.4, 4.8, or 9.6 kbit/s are possible. Transparent bearer services do not try to recover lost data in case of, for example, shadowing or interruptions due to handover.

Non-transparent bearer services use protocols of layers two and three to implement error correction and flow control. These services use the transparent bearer services, adding **radio link protocol (RLP)**. This protocol comprises mechanisms of **high-level data link control (HDLC)**, (Halsall, 1996) and special selective-reject mechanisms to trigger retransmission of erroneous data. The achieved bit error rate is less than 10^{-7} , but now throughput and delay may vary depending on transmission quality.

Using transparent and non-transparent services, GSM specifies several bearer services for interworking with PSTN, ISDN, and packet-switched public data networks (PSPDN) like X.25, which is available worldwide. Data transmission can be full-duplex, synchronous with data rates of 1.2, 2.4, 4.8, and 9.6 kbit/s or full-duplex, asynchronous from 300 to 9,600 bit/s (ETSI, 1991a). Clearly, these relatively low data rates reflect the assumption that data services will only constitute some small percentage of the overall traffic. While this is still true of GSM networks today, the relation of data and voice services is changing, with data becoming more and more important. This development is also reflected in the new data services (see section 4.1.8).

4.1.1.2 Teleservices

GSM mainly focuses on voice-oriented tele services. These comprise encrypted voice transmission, message services, and basic data communication with terminals as known from the PSTN or ISDN (e.g., fax). However, as the main service is **telephony**, the primary goal of GSM was the provision of high-quality digital voice transmission, offering at least the typical bandwidth of 3.1 kHz of analog phone systems. Special codecs (coder/decoder) are used for voice transmission, while other codecs are used for the transmission of analog data for communication with traditional computer modems used in, e.g., fax machines.

Another service offered by GSM is the **emergency number**. The same number can be used throughout Europe. This service is mandatory for all providers and free of charge. This connection also has the highest priority, possibly pre-empting other connections, and will automatically be set up with the closest emergency center.

A useful service for very simple message transfer is the **short message service (SMS)**, which offers transmission of messages of up to 160 characters. SMS messages do not use the standard data channels of GSM but exploit unused capacity in the signalling channels (see section 4.1.3.1). Sending and receiving of SMS is possible during data or voice transmission. SMS was in the GSM standard from the beginning; however, almost no one used it until millions of young people discovered this service in the mid-nineties as a fun service. SMS

can be used for “serious” applications such as displaying road conditions, e-mail headers or stock quotes, but it can also transfer logos, ringtones, horoscopes and love letters. Today more than 30 billion short messages are transferred worldwide per month! SMS is big business today, not only for the network operators, but also for many content providers. It should be noted that SMS is typically the only way to reach a mobile phone from within the network. Thus, SMS is used for updating mobile phone software or for implementing so-called push services (see chapter 10).

The successor of SMS, the **enhanced message service (EMS)**, offers a larger message size (e.g., 760 characters, concatenating several SMSs), formatted text, and the transmission of animated pictures, small images and ring tones in a standardized way (some vendors offered similar proprietary features before). EMS never really took off as the **multimedia message service (MMS)** was available. (Nokia never liked EMS but pushed Smart Messaging, a proprietary system.) MMS offers the transmission of larger pictures (GIF, JPG, WBMP), short video clips etc. and comes with mobile phones that integrate small cameras. MMS is further discussed in the context of WAP in chapter 10.

Another non-voice tele service is **group 3 fax**, which is available worldwide. In this service, fax data is transmitted as digital data over the analog telephone network according to the ITU-T standards T.4 and T.30 using modems. Typically, a transparent fax service is used, i.e., fax data and fax signaling is transmitted using a transparent bearer service. Lower transmission quality causes an automatic adaptation of the bearer service to lower data rates and higher redundancy for better FEC.

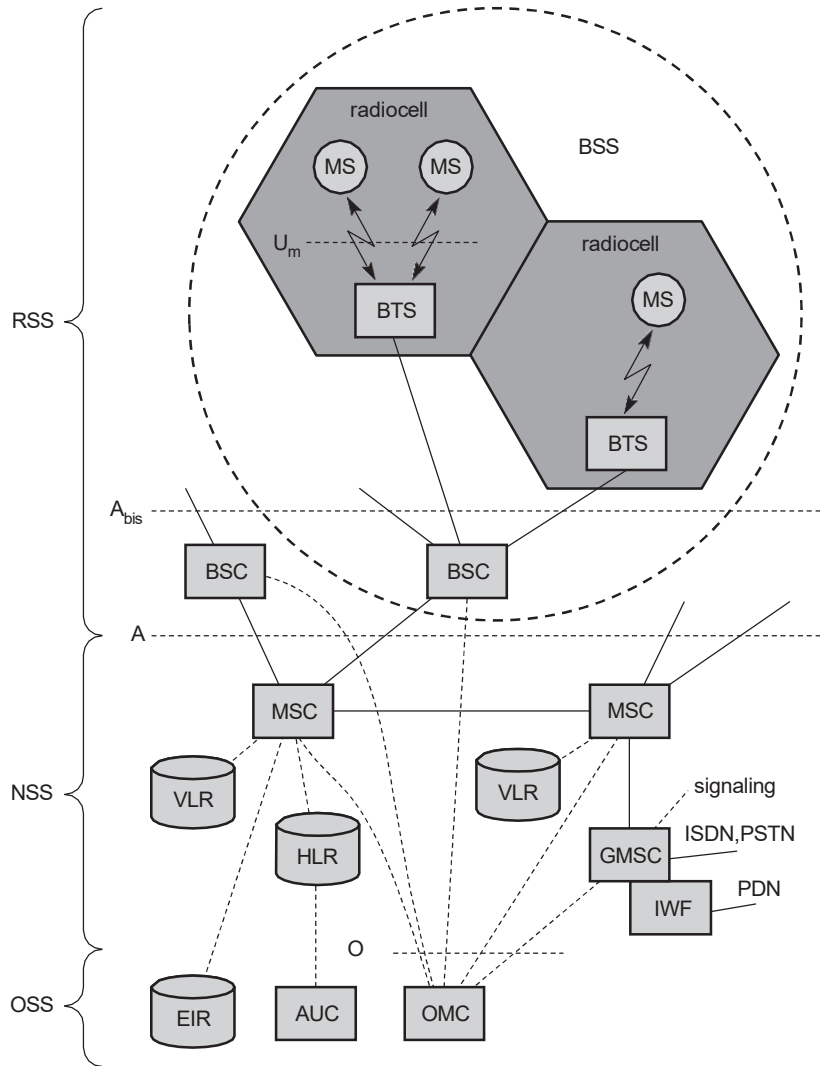
4.1.1.3 Supplementary services

In addition to tele and bearer services, GSM providers can offer **supplementary services**. Similar to ISDN networks, these services offer various enhancements for the standard telephony service, and may vary from provider to provider. Typical services are **user identification**, **call redirection**, or **forwarding** of ongoing calls. Standard ISDN features such as **closed user groups** and **multi-party** communication may be available. Closed user groups are of special interest to companies because they allow, for example, a company-specific GSM sub-network, to which only members of the group have access.

4.1.2 System architecture

As with all systems in the telecommunication area, GSM comes with a hierarchical, complex system architecture comprising many entities, interfaces, and acronyms. Figure 4.4 gives a simplified overview of the GSM system as specified in ETSI (1991b). A GSM system consists of three subsystems, the **radio subsystem (RSS)**, the **network and switching subsystem (NSS)**, and the **operation subsystem (OSS)**. Each subsystem will be discussed in more detail in the following sections. Generally, a GSM customer only notices a very small fraction of the whole network – the mobile stations (MS) and some antenna masts of the base transceiver stations (BTS).

Figure 4.4
Functional architecture
of a GSM system



4.1.2.1 Radiosubsystem

As the name implies, the **radio subsystem (RSS)** comprises all radio specific entities, i.e., the **mobile stations (MS)** and the **base station subsystem (BSS)**. Figure 4.4 shows the connection between the RSS and the NSS via the **A interface** (solid lines) and the connection to the OSS via the **O interface** (dashed lines). The A interface is typically based on circuit-switched PCM-30 systems (2.048Mbit/s), carrying up to 3064kbit/s connections, whereas the O interface uses the Signalling System No. 7 (SS7) based on X.25 carrying management data to/from the RSS.

- **Basestation subsystem (BSS):** A GSM network comprises many BSSs, each controlled by a base station controller (BSC). The BSS performs all functions necessary to maintain radio connections to an MS, coding/decoding of voice, and rate adaptation to/from the wireless network part. Besides a BSC, the BSS contains several BTSs.
- **Base transceiver station (BTS):** A BTS comprises all radio equipment, i.e., antennas, signal processing, amplifiers necessary for radio transmission. A BTS can form a radio cell or, using sectorized antennas, several cells (see section 2.8), and is connected to MS via the **U_m interface** (ISDN U interface for mobile use), and to the BSC via the **A_{bis} interface**. The U_m interface contains all the mechanisms necessary for wireless transmission (TDMA, FDMA etc.) and will be discussed in more detail below. The A_{bis} interface consists of 16 or 64 kbit/s connections. A GSM cell can measure between some 100 m and 35 km depending on the environment (buildings, open space, mountains etc.) but also expected traffic.
- **Base station controller (BSC):** The BSC basically manages the BTSs. It reserves radio frequencies, handles the handover from one BTS to another within the BSS, and performs paging of the MS. The BSC also multiplexes the radio channels onto the fixed network connections at the A interface.

Table 4.1 gives an overview of the tasks assigned to the BSC and BTS or of tasks in which these entities support other entities in the network.

- **Mobile station (MS):** The MS comprises all user equipment and software needed for communication with a GSM network. An MS consists of user independent hard- and software and of the **subscriber identity module (SIM)**, which stores all user-specific data that is relevant to GSM.³ While an MS can be identified via the **international mobile equipment identity (IMEI)**, a user can personalize any MS using his or her SIM, i.e., user-specific mechanisms like charging and authentication are based on the SIM, not on the device itself. Device-specific mechanisms, e.g., theft protection, use the device specific IMEI. Without the SIM, only emergency calls are possible. The SIM card contains many identifiers and tables, such as card-type, serial number, a list of subscribed services, a **personal identity number (PIN)**, a **PIN unblocking key (PUK)**, an **authentication key K_i**, and the **international mobile subscriber identity (IMSI)** (ETSI, 1991c). The PIN is used to unlock the MS. Using the wrong PIN three times will lock the SIM. In such cases, the PUK is needed to unlock the SIM. The MS stores dynamic information while logged onto the GSM system, such as, e.g., the **cipher key K_c** and the location information consisting of a **temporary mobile subscriber identity (TMSI)** and the **location area identification (LAI)**. Typical MSs for GSM 900 have a transmit power of up to 2 W, whereas for GSM 1800 1 W is enough due to the smaller cell size. Apart from the telephone interface, an

³ Many additional items can be stored on the mobile device. However, this is irrelevant to GSM.

Function	BTS	BSC
Management of radio channels		X
Frequency hopping	X	X
Management of terrestrial channels		X
Mapping of terrestrial onto radio channels		X
Channel coding and decoding	X	
Rate adaptation	X	
Encryption and decryption	X	X
Paging	X	X
Uplink signal measurement	X	
Traffic measurement		X
Authentication		X
Location registry, location update		X
Handover management		X

Table 4.1 Tasks of the BTS and BSC within a BSS

MS can also offer other types of interfaces to users with display, loudspeaker, microphone, and programmable soft keys. Further interfaces comprise computer modems, IrDA, or Bluetooth. Typical MSs, e.g., mobile phones, comprise many more vendor-specific functions and components, such as cameras, fingerprint sensors, calendars, address books, games, and Internet browsers. Personal digital assistants (PDA) with mobile phone functions are also available. The reader should be aware that an MS could also be integrated into a car or be used for location tracking of a container.

4.1.2.2 Network and switching subsystem

The “heart” of the GSM system is formed by the **network and switching subsystem (NSS)**. The NSS connects the wireless network with standard public networks, performs handovers between different BSSs, comprises functions for worldwide localization of users and supports charging, accounting, and roaming of users between different providers in different countries. The NSS consists of the following switches and databases:

- **Mobile services switching center (MSC):** MSCs are high-performance digital ISDN switches. They set up connections to other MSCs and to the BSCs via the A interface, and form the fixed backbone network of a GSM system. Typically, an MSC manages several BSCs in a geographical region. A **gateway MSC (GMSC)** has additional connections to other fixed networks, such as PSTN and ISDN. Using additional **interworking functions (IWF)**, an MSC

can also connect to **public data networks (PDN)** such as X.25. An MSC handles all signaling needed for connection setup, connection release and handover of connections to other MSCs. The **standard signaling system No. 7 (SS7)** is used for this purpose. SS7 covers all aspects of control signaling for digital networks (reliable routing and delivery of control messages, establishing and monitoring of calls). Features of SS7 are number portability, free phone/toll/collect/credit calls, call forwarding, three-way calling etc. An MSC also performs all functions needed for supplementary services such as call forwarding, multi-party calls, reverse charging etc.

- **Home location register (HLR):** The HLR is the most important database in a GSM system as it stores all user-relevant information. This comprises static information, such as the **mobile subscriber ISDN number (MSISDN)**, subscribed services (e.g., call forwarding, roaming restrictions, GPRS), and the **international mobile subscriber identity (IMSI)**. Dynamic information is also needed, e.g., the current **location area (LA)** of the MS, the **mobile subscriber roaming number (MSRN)**, the current VLR and MSC. As soon as an MS leaves its current LA, the information in the HLR is updated. This information is necessary to localize a user in the worldwide GSM network. All these user-specific information elements only exist once for each user in a single HLR, which also supports charging and accounting. The parameters will be explained in more detail in section 4.1.5. HLRs can manage data for several million customers and contain highly specialized data bases which must fulfill certain real-time requirements to answer requests within certain time-bounds.
- **Visitor location register (VLR):** The VLR associated to each MSC is a dynamic database which stores all important information needed for the MS users currently in the LA that is associated to the MSC (e.g., IMSI, MSISDN, HLR address). If a new MS comes into an LA the VLR is responsible for, it copies all relevant information for this user from the HLR. This hierarchy of VLR and HLR avoids frequent HLR updates and long-distance signaling of user information. The typical use of HLR and VLR for user localization will be described in section 4.1.5. Some VLRs in existence, are capable of managing up to one million customers.

4.1.2.3 Operations subsystem

The third part of a GSM system, the **operation subsystem (OSS)**, contains the necessary functions for network operation and maintenance. The OSS possesses network entities of its own and accesses other entities via SS7 signaling (see Figure 4.4). The following entities have been defined:

- **Operation and maintenance center (OMC):** The OMC monitors and controls all other network entities via the O interface (SS7 with X.25). Typical OMC management functions are traffic monitoring, status reports of network entities, subscriber and security management, or accounting and billing. OMCs use the concept of **telecommunication management network (TMN)** as standardized by the ITU-T.

- **Authentication centre (AuC):** As the radio interface and mobile stations are particularly vulnerable, a separate AuC has been defined to protect user identity and data transmission. The AuC contains the algorithms for authentication as well as the keys for encryption and generates the values needed for user authentication in the HLR. The AuC may, in fact, be situated in a special protected part of the HLR.
- **Equipment identity register (EIR):** The EIR is a database for all IMEIs, i.e., it stores all device identifications registered for this network. As MSs are mobile, they can be easily stolen. With a valid SIM, anyone could use the stolen MS. The EIR has a blacklist of stolen (or locked) devices. In theory an MS is useless as soon as the owner has reported a theft. Unfortunately, the blacklists of different providers are not usually synchronized and the illegal use of a device in another operator's network is possible (the reader may speculate as to why this is the case). The EIR also contains a list of valid IMEIs (white list), and a list of malfunctioning devices (gray list).

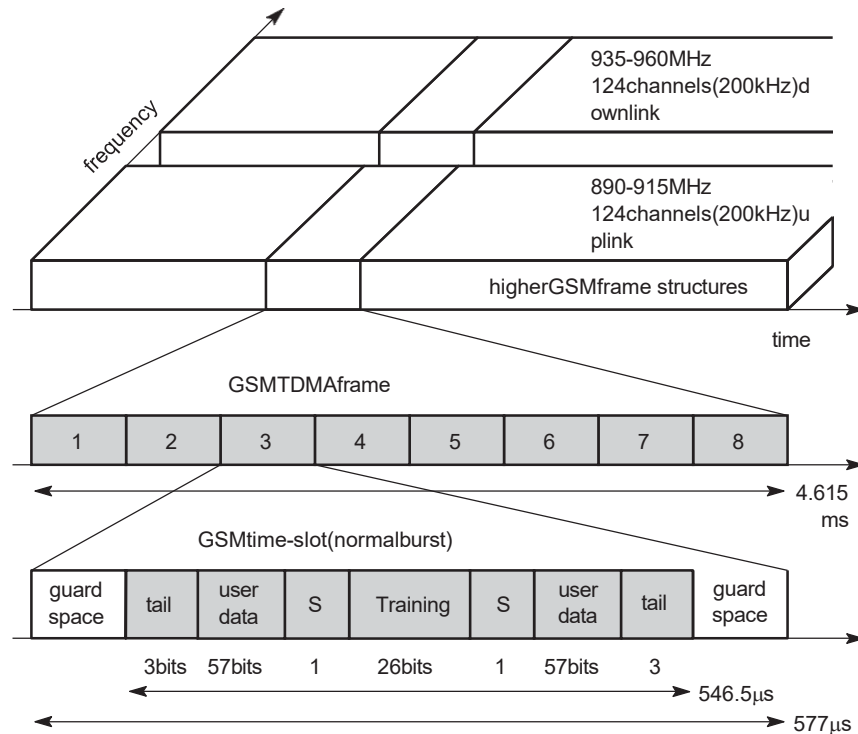
4.1.3 Radio interface

The most interesting interface in a GSM system is U_m , the radio interface, as it comprises many mechanisms presented in chapters 2 and 3 for multiplexing and media access. GSM implements SDMA using cells with BTS and assigns an MS to a BTS. Furthermore, FDD is used to separate downlink and uplink as shown in Figures 3.3 and 4.5. Media access combines TDMA and FDMA. In GSM 900, 124 channels, each 200 kHz wide, are used for FDMA, whereas GSM 1800 uses, 374 channels. Due to technical reasons, channels 1 and 124 are not used for transmission in GSM 900. Typically, 32 channels are reserved for organizational data; the remaining 90 are used for customers. Each BTS then manages a single channel for organizational data and, e.g., up to 10 channels for user data. The following example is based on the GSM 900 system, but GSM works in a similar way at 1800 and 1900 MHz.

While Figure 3.3 in chapter 3 has already shown the FDM in GSM, Figure 4.5 also shows the TDM used. Each of the 248 channels is additionally separated in time via a **GSM TDMA frame**, i.e., each 200 kHz carrier is subdivided into frames that are repeated continuously. The duration of a frame is 4.615 ms. A frame is again subdivided into **8 GSM time slots**, where each slot represents a physical TDM channel and lasts for 577 μ s. Each TDM channel occupies the 200 kHz carrier for 577 μ s every 4.615 ms.

Data is transmitted in small portions, called **bursts**. Figure 4.5 shows a so-called **normal burst** as used for data transmission inside a time slot (user and signaling data). In the diagram, the burst is only 546.5 μ s long and contains 148 bits. The remaining 30.5 μ s are used as **guard space** to avoid overlapping with other bursts due to different path delays and to give the transmitter time to turn on and off. Filling the whole slot with data allows for the transmission of

Figure 4.5
GSM TDMA frame,
slots, and bursts



156.25 bit within 577 μs. Each physical TDM channel has a raw data rate of about 33.8 kbit/s, each radio carrier transmits approximately 270 kbit/s over the U_m interface.

The first and last three bits of a normal burst (**tail**) are all set to 0 and can be used to enhance the receiver performance. The **training** sequence in the middle of a slot is used to adapt the parameters of the receiver to the current path propagation characteristics and to select the strongest signal in case of multi-path propagation. A flag **S** indicates whether the **data** field contains user or network control data. Apart from the normal burst, ETSI (1993a) defines four more bursts for data transmission: a **frequency correction** burst allows the MS to correct the local oscillator to avoid interference with neighboring channels, a **asynchronization burst** with an extended training sequence synchronizes the MS with the BTS in time, an **access burst** is used for the initial connection setup between MS and BTS, and finally a **dummy burst** is used if no data is available for a slot.

Two factors allow for the use of simple transmitter hardware: on the one hand, the slots for uplink and downlink of a physical TDM channel are separated in frequency (45 MHz for GSM 900, 95 MHz for GSM 1800 using FDD). On the other hand, the TDMA frames are shifted in time for three slots, i.e., if the BTS sends data at time t in slot n on the downlink, the MS accesses slot

one on the uplink at time $t_0 + 3 \cdot 577 \mu\text{s}$. An MS does not need a full-duplex transmitter, a simpler half-duplex transmitter switching between receiving and sending is enough.

To avoid frequency selective fading, GSM specifies an optional **slow frequency hopping** mechanism. MS and BTS may change the carrier frequency after each frame based on a common hopping sequence. An MS changes its frequency between up and downlink slots respectively.

4.1.3.1 Logical channels and frame hierarchy

While the previous section showed the physical separation of the medium into 8×124 duplex channels, this section presents logical channels and a hierarchy of frames based on the combination of these physical channels. A physical channel consists of a slot, repeated every 4.615 ms . Think of a logical channel C_1 that only takes up every fourth slot and another logical channel C_2 that uses every other slot. Both logical channels could use the same physical channel with the pattern $C_1 C_2 \times C_2 C_1 C_2 \times C_2 C_1$ etc. (The \times indicates that the physical channel still has some capacity left.)

GSM specifies two basic groups of logical channels, i.e., traffic channels and control channels:⁴

- Traffic channels (TCH):** GSM uses a TCH to transmit user data (e.g., voice, fax). Two basic categories of TCHs have been defined, i.e., **full-rate TCH (TCH/F)** and **half-rate TCH (TCH/H)**. A TCH/F has a data rate of 22.8 kbit/s , whereas TCH/H only has 11.4 kbit/s . With the voice codecs available at the beginning of the GSM standardization, 13 kbit/s were required, whereas the remaining capacity of the TCH/F (22.8 kbit/s) was used for error correction (**TCH/FS**). Improved codes allow for better voice coding and can use a TCH/H. Using these TCH/HSs doubles the capacity of the GSM system for voice transmission. However, speech quality decreases with the use of TCH/HS and many providers try to avoid using them. The standard codecs for voice are called **full rate (FR, 13 kbit/s)** and **half rate (HR, 5.6 kbit/s)**. A newer codec, **enhanced full rate (EFR)**, provides better voice quality than FR as long as the transmission error rate is low. The generated data rate is only 12.2 kbit/s . New codecs, which automatically choose the best mode of operation depending on the error rate (AMR, adaptive multi-rate), will be used together with 3G systems. An additional increase in voice quality is provided by the so-called **tandem free operation (TFO)**. This mode can be used if two MSs exchange voice data. In this case, coding to and from PCM encoded voice (standard in ISDN) can be skipped and the GSM encoded voice data is directly exchanged. Data transmission in GSM is possible at many different data rates, e.g., **TCH/F4.8** for 4.8 kbit/s , **TCH/F9.6** for 9.6 kbit/s , and, as a newer specification, **TCH/F14.4** for 14.4 kbit/s . These logical channels differ in terms of their coding schemes and error correction capabilities.

⁴More information about channels can be found in Goodman (1997) and ETSI (1993a).

- **Control channels (CCH):** Many different CCHs are used in a GSM system to control medium access, allocation of traffic channels or mobility management. Three groups of control channels have been defined, each again with subchannels (maybe you can imagine why the initial specification already needed over 5,000 pages):
 - **Broadcast control channel (BCCH):** A BTS uses this channel to signal information to all MSs within a cell. Information transmitted in this channel is, e.g., the cell identifier, options available within this cell (frequency hopping), and frequencies available inside the cell and in neighboring cells. The BTS sends information for frequency correction via the **frequency correction channel (FCCH)** and information about time synchronization via the **synchronization channel (SCH)**, where both channels are subchannels of the BCCH.
 - **Common control channel (CCCH):** All information regarding connection setup between MS and BS is exchanged via the CCCH. For calls toward an MS, the BTS uses the **paging channel (PCH)** for paging the appropriate MS. If an MS wants to set up a call, it uses the **random access channel (RACH)** to send data to the BTS. The RACH implements multiple access (all MSs within a cell may access this channel) using slotted Aloha. This is where a collision may occur with other MSs in a GSM system. The BTS uses the **access grant channel (AGCH)** to signal an MS that it can use a TCH or SDCCH for further connection setup.
 - **Dedicated control channel (DCCH):** While the previous channels have all been unidirectional, the following channels are bidirectional. As long as an MS has not established a TCH with the BTS, it uses the **stand-alone dedicated control channel (SDCCH)** with a low data rate (782 bit/s) for signaling. This can comprise authentication, registration or other data needed for setting up a TCH. Each TCH and SDCCH has a **slow associated dedicated control channel (SACCH)** associated with it, which is used to exchange system information, such as the channel quality and signal power level. Finally, if more signaling information needs to be transmitted and a TCH already exists, GSM uses a **fast associated dedicated control channel (FACCH)**. The FACCH uses the time slots which are otherwise used by the TCH. This is necessary in the case of handovers where BTS and MS have to exchange larger amounts of data in less time.

However, these channels cannot use time slots arbitrarily – GSM specifies a very elaborate multiplexing scheme that integrates several hierarchies of frames. If we take a simple TCH/F for user data transmission, each TCH/F will have an associated SACCH for slow signaling. If fast signaling is required, the FACCH uses the time slots for the TCH/F. A typical usage pattern of a physical channel for data transmission now looks like this (with T indicating the user traffic in the TCH/F and S indicating the signalling traffic in the SACCH):

TTTTTTTTTTTTSTTTTTTTTTTTTTTSTTTTTTTTTTTTTSTTTTTTTTTTTTT

x

Twelve slots with user data are followed by a signalling slot. Again 12 slots with user data follow, then an unused slot. This pattern of 26 slots is repeated over and over again. In this case, only 24 out of 26 physical slots are used for the TCH/F. Now recall that each normal burst used for data transmission carries 114 bit user data and is repeated every 4.615 ms. This results in a data rate of 24.7 kbit/s. As the TCH/F only uses 24/26 of the slots, the final data rate is 22.8 kbit/s as specified for the TCH/F. The SACCH thus has a capacity of 950 bit/s.

This periodic pattern of 26 slots occurs in all TDMA frames with a TCH. The combination of these frames is called **traffic multiframe**. Figure 4.6 shows the logical combination of 26 frames (TDMA frames with a duration of 4.615 ms) to a multiframe with a duration of 120 ms. This type of multiframe is used for TCHs, SACCHs for TCHs, and FACCHs. As these logical channels are all associated with user traffic, the multiframe is called **traffic multiframe**. TDMA frames containing (signaling) data for the other logical channels are combined to a **control multiframe**. Control multiframes consist of 51 TDMA frames and have a duration of 235.4 ms.

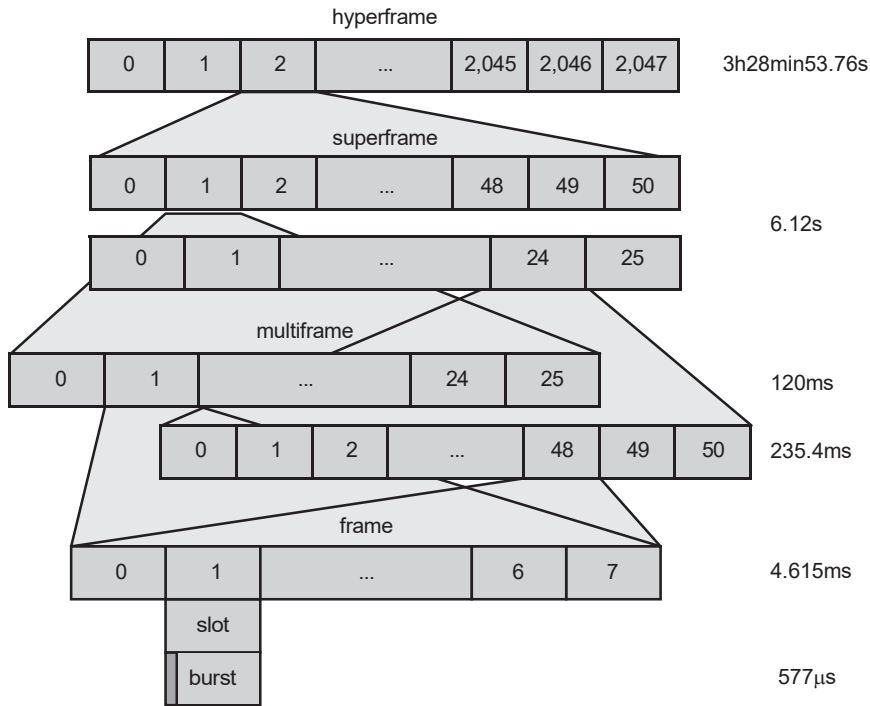


Figure 4.6
GSM structuring of time using a frame hierarchy

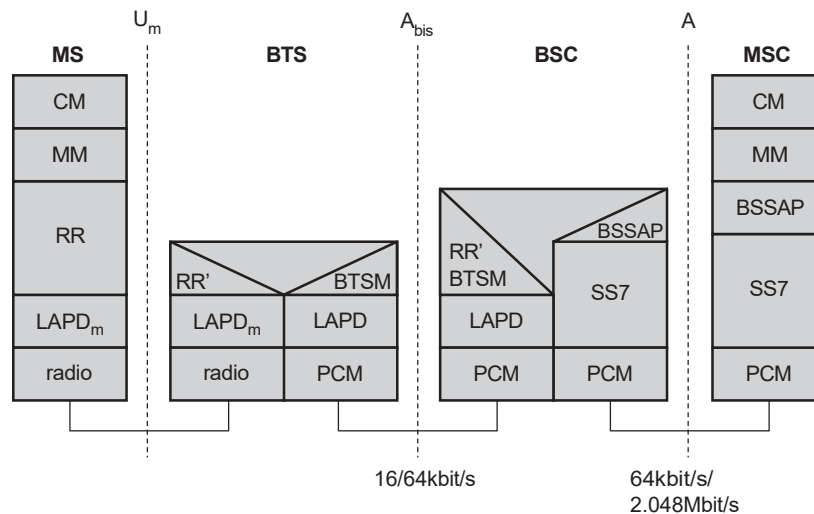
This logical frame hierarchy continues, combining 26 multiframes with 51 frames or 51 multiframes with 26 frames to form a **superframe**. 2,048 superframes build a **hyperframe** with a duration of almost 3.5 hours. Altogether, 2,715,648 TDMA frames form a hyperframe. This large logical structure is needed for encryption – GSM counts each TDMA frame, with the frame number forming input for the encryption algorithm. The frame number plus the slot number uniquely identify each time slot in GSM.

4.1.4 Protocols

Figure 4.7 shows the protocol architecture of GSM with signaling protocols, interfaces, as well as the entities already shown in Figure 4.4. The main interest lies in the U_m interface, as the other interfaces occur between entities in a fixed network. **Layer 1**, the physical layer, handles all **radio**-specific functions. This includes the creation of bursts according to the five different formats, **multiplexing** of bursts into a TDMA frame, **synchronization** with the BTS, detection of idle channels, and measurement of the **channel quality** on the downlink. The physical layer at U_m uses GMSK for digital **modulation** and performs **encryption/decryption** of data, i.e., encryption is not performed end-to-end, but only between MS and BSS over the air interface.

Synchronization also includes the correction of the individual path delay between an MS and the BTS. All MSs within a cell use the same BTS and thus must be synchronized to this BTS. The BTS generates the time-structure of frames, slots etc. A problematic aspect in this context are the different round trip times (RTT). An MS close to the BTS has a very short RTT, whereas an MS 35 km away already exhibits an RTT of around 0.23 ms. If the MS far away used the slot structure with-

Figure 4.7
Protocol architecture
for signaling



out correction, large guard spaces would be required, as 0.23 ms are already 40 per cent of the 0.577 ms available for each slot. Therefore, the BTS sends the current RTT to the MS, which then adjusts its access times so that all bursts reach the BTS within their limits. This mechanism reduces the guard space to only 30.5 μ s or five percent (see Figure 4.5). Adjusting the access is controlled via the variable **timing advance**, where a burst can be shifted up to 63 bit times earlier, with each bit having a duration of 3.69 μ s (which results in the 0.23 ms needed). As the variable timing advance cannot be extended a burst cannot be shifted earlier than 63 bit times. This results in the 35 km maximum distance between an MS and a BTS. It might be possible to receive this signal over longer distances; to avoid collisions at the BTS, access cannot be allowed.⁵

The main tasks of the physical layer comprise **channel coding and error detection/correction**, which is directly combined with the coding mechanisms. Channel coding makes extensive use of different **forward error correction (FEC)** schemes. FEC adds redundancy to user data, allowing for the detection and correction of selected errors. The power of an FEC scheme depends on the amount of redundancy, coding algorithm and further interleaving of data to minimize the effects of burst errors. The FEC is also the reason why error detection and correction occurs in layer one and not in layer two as in the ISO/OSI reference model. The GSM physical layer tries to correct errors, but it does not deliver erroneous data to the higher layer.

Different logical channels of GSM use different coding schemes with different correction capabilities. Speech channels need additional coding of voice data after analog to digital conversion, to achieve a data rate of 22.8 kbit/s (using the 13 kbit/s from the voice codec plus redundancy, CRC bits, and interleaving (Goodman, 1997)). As voice was assumed to be the main service in GSM, the physical layer also contains special functions, such as **voice activity detection (VAD)**, which transmits voice data only when there is a voice signal. This mechanism helps to decrease interference as a channel might be silent approximately 60 percent of the time (under the assumption that only one person speaks at the same time and some extra time is needed to switch between the speakers). During periods of silence (e.g., if a user needs time to think before talking), the physical layer generates **comfort noise** to fake a connection (complete silence would probably confuse a user), but no actual transmission takes place. The noise is even adapted to the current background noise at the communication partner's location.

All this interleaving of data for a channel to minimize interference due to burst errors and the recurrence pattern of a logical channel generates a **delay** for transmission. The delay is about 60 ms for a TCH/FS and 100 ms for a TCH/F9.6

⁵ A special trick allows for larger cells. If the timing advance for MSs that are further away than 35 km is set to zero, the bursts arriving from these MSs will fall into the following time slot. Reception of data is simply shifted one time slot and again the timing advance may be used up to a distance of 70 km (under simplified assumptions). Using this special trick, the capacity of a cell is decreased (near and far MSs cannot be mixed arbitrarily), but coverage of GSM is extended. Network operators may choose this approach, e.g., in coastal regions.

(within 100 ms signals in fixed networks easily travel around the globe). These times have to be added to the transmission delay if communicating with an MS instead of a standard fixed station (telephone, computer etc.) and may influence the performance of any higher layer protocols, e.g., for computer data transmission (see chapter 9).

Signaling between entities in a GSM network requires higher layers (see Figure 4.7). For this purpose, the **LAPD_m** protocol has been defined at the **U_m** interface for **layer two**. **LAPD_m**, as the name already implies, has been derived from link access procedure for the D-channel (**LAPD**) in ISDN systems, which is a version of HDLC (Goodman, 1997), (Halsall, 1996). **LAPD_m** is a lightweight **LAPD** because it does not need synchronization flags or checksumming for error detection. (The GSM physical layer already performs these tasks.) **LAPD_m** offers reliable data transfer over connections, re-sequencing of data frames, and flow control (ETSI, 1993b), (ETSI, 1993c). As there is no buffering between layer one and two, **LAPD_m** has to obey the frame structures, recurrence patterns etc. defined for the **U_m** interface. Further services provided by **LAPD_m** include segmentation and reassembly of data and acknowledged/unacknowledged data transfer.

The network layer in GSM, **layer three**, comprises several sublayers as Figure 4.7 shows. The lowest sublayer is the **radio resource management (RR)**. Only a part of this layer, **RR'**, is implemented in the BTS, the remainder is situated in the BSC. The functions of **RR'** are supported by the BSC via the **BTS management (BTSM)**. The main tasks of **RR** are setup, maintenance, and release of radio channels. **RR** also directly accesses the physical layer for radio information and offers a reliable connection to the next higher layer.

Mobility management (MM) contains functions for registration, authentication, identification, location updating, and the provision of a **temporary mobile subscriber identity (TMSI)** that replaces the **international mobile subscriber identity (IMSI)** and which hides the real identity of an MS user over the air interface. While the **IMSI** identifies a user, the **TMSI** is valid only in the current location area of a VLR. **MM** offers a reliable connection to the next higher layer.

Finally, the **call management (CM)** layer contains three entities: **call control (CC)**, **short message service (SMS)**, and **supplementary service (SS)**. **SMS** allows for message transfer using the control channels **SDCCH** and **SACCH** (if no signaling data is sent), while **SS** offers the services described in section 4.1.1.3. **CC** provides a point-to-point connection between two terminals and is used by higher layers for call establishment, call clearing and change of call parameters. This layer also provides functions to send in-band tones, called **dual tone multiple frequency (DTMF)**, over the GSM network. These tones are used, e.g., for the remote control of answering machines or the entry of PINs in electronic banking and are, also used for dialing in traditional analog telephone systems. These tones cannot be sent directly over the voice codec of a GSM MS, as the codec would distort the tones. They are transferred as signals and then converted into tones in the fixed network part of the GSM system.

Additional protocols are reused at the A_{bis} and A interfaces (the internal interfaces of a GSM system not presented here). Data transmission at the physical layer typically uses **pulse code modulation (PCM)** systems. While PCM systems offer transparent 64 kbit/s channels, GSM also allows for the submultiplexing of four 16 kbit/s channels into a single 64 kbit/s channel (16 kbit/s are enough for user data from an MS). The physical layer at the A interface typically includes leased lines with 2.048 Mbit/s capacity. LAPD is used for layer two at A_{bis} , BTSM for BTS management.

Signaling system No. 7 (SS7) is used for signaling between an MSC and a BSC. This protocol also transfers all management information between MSCs, HLR, VLRs, AuC, EIR, and OMC. An MSC can also control a BSS via a **BSS application part (BSSAP)**.

4.1.5 Localization and calling

One fundamental feature of the GSM system is the automatic, worldwide localization of users. The system always knows where a user currently is, and the same phone number is valid worldwide. To provide this service, GSM performs periodic location updates even if a user does not use the mobile station (provided that the MS is still logged into the GSM network and is not completely switched off). The HLR always contains information about the current location (only the location area, not the precise geographical location), and the VLR currently responsible for the MS informs the HLR about location changes. As soon as an MS moves into the range of a new VLR (a new location area), the HLR sends all user data needed to the new VLR. Changing VLRs with uninterrupted availability of all services is also called **roaming**. Roaming can take place within the network of one provider, between two providers in one country (national roaming is, often not supported due to competition between operators), but also between different providers in different countries (international roaming). Typically, people associate international roaming with the term roaming as it is this type of roaming that makes GSM very attractive: one device, over 190 countries!

To locate an MS and to address the MS, several numbers are needed:

- **Mobile station international ISDN number (MSISDN):**⁶ The only important number for a user of GSM is the phone number. Remember that the phone number is not associated with a certain device but with the SIM, which is personalized for a user. The MSISDN follows the ITU-T standard E.164 for addresses as it is also used in fixed ISDN networks. This number consists of the **country code (CC)** (e.g., +49 179 1234567 with 49 for Germany), the **national destination code (NDC)** (i.e., the address of the network provider, e.g., 179), and the **subscriber number (SN)**.

⁶ In other types of documentation, this number is also called 'Mobile Subscriber ISDN Number' or 'Mobile Station ISDN Number'. Even the original ETSI standards use different wordings for the same acronym. However, the term 'subscriber' is much better suited as it expresses the independence of the user-related number from the device (station).

- **International mobile subscriber identity (IMSI):** GSM uses the IMSI for internal unique identification of a subscriber. IMSI consists of a **mobile country code (MCC)** (e.g., 240 for Sweden, 208 for France), the **mobile network code (MNC)** (i.e., the code of the network provider), and finally the **mobile subscriber identification number (MSIN)**.
- **Temporary mobile subscriber identity (TMSI):** To hide the IMSI, which would give away the exact identity of the user signaling over the air interface, GSM uses the 4 byte TMSI for local subscriber identification. TMSI is selected by the current VLR and is only valid temporarily and within the location area of the VLR (for an ongoing communication TMSI and LAI are sufficient to identify a user; the IMSI is not needed). Additionally, a VLR may change the TMSI periodically.
- **Mobile station⁷ roaming number (MSRN):** Another temporary address that hides the identity and location of a subscriber is MSRN. The VLR generates this address on request from the MSC, and the address is also stored in the HLR. MSRN contains the current **visitor country code (VCC)**, the **visitor national destination code (VNDC)**, the identification of the current MSC together with the subscriber number. The MSRN helps the HLR to find a subscriber for an incoming call.

All these numbers are needed to find a subscriber and to maintain the connection with a mobile station. The interesting case is the **mobile terminated call (MTC)**, i.e., a situation in which a station calls a mobile station (the calling station could be outside the GSM network or another mobile station). Figure 4.8 shows the basic steps needed to connect the calling station with the mobile user. In step 1, a user dials the phone number of a GSM subscriber. The fixed network (PSTN) notices (looking at the destination code) that the number belongs to a user in the GSM network and forwards the call setup to the Gateway MSC (2). The GMSC identifies the HLR for the subscriber (which is coded in the phone number) and signals the call setup to the HLR (3). The HLR now checks whether the number exists and whether the user has subscribed to the requested services, and requests an MSRN from the current VLR (4). After receiving the MSRN (5), the HLR can determine the MSC responsible for the MS and forwards this information to the GMSC (6). The GMSC can now forward the call setup request to the MSC indicated (7).

From this point on, the MSC is responsible for all further steps. First, it requests the current status of the MS from the VLR (8). If the MS is available, the MSC initiates paging in all cells it is responsible for (i.e. the location area, LA, 10), as searching for the right cell would be too time consuming (but this approach puts some load on the signaling channels so optimizations exist). The

⁷ Here, a discrepancy exists between ITU-T standards and ETSI's GSM. MSCs are not mobile stations or mobile subscribers. Typically, almost all MS in GSM refer to subscribers, as identifiers are not dependent on the station, but on the subscriber identity (stored in the SIM).

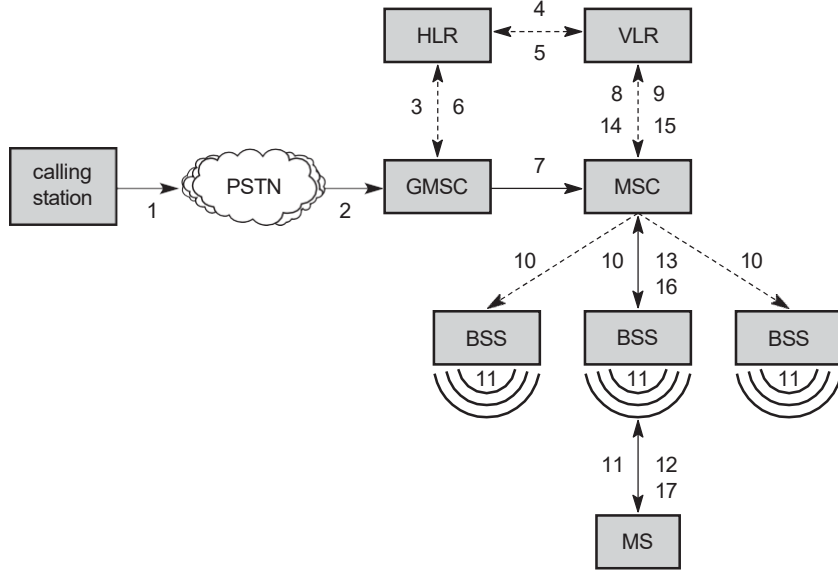


Figure 4.8
Mobile terminated call (MTC)

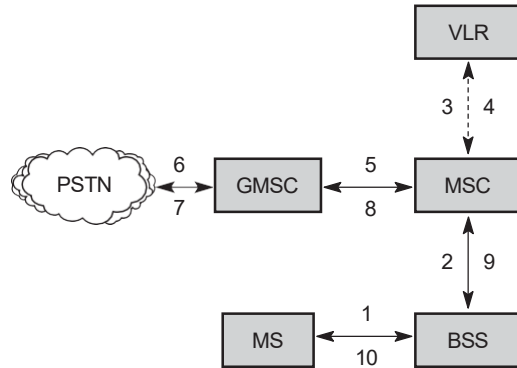
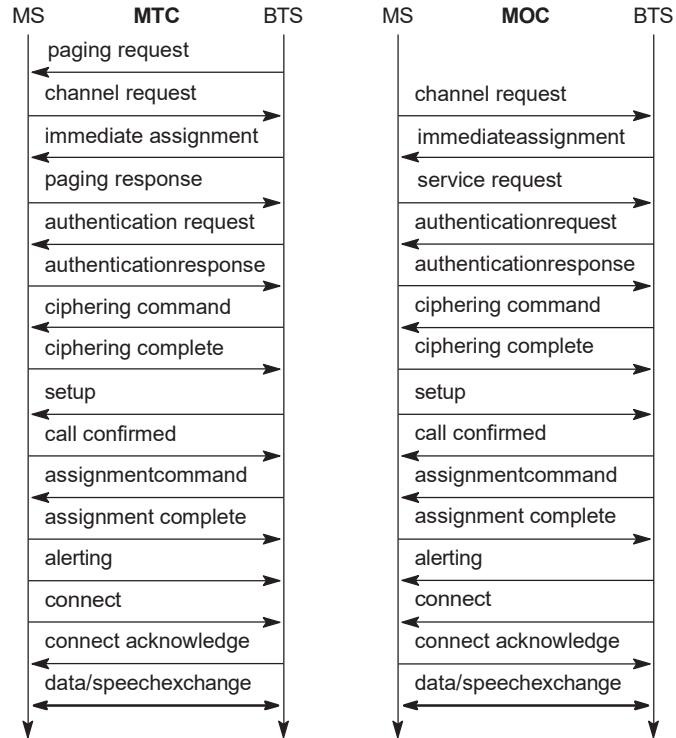


Figure 4.9
Mobile originated call (MOC)

BTSs of all BSSs transmit this paging signal to the MS (11). If the MS answers (12 and 13), the VLR has to perform security checks (set up encryption etc.). The VLR then signals to the MSC to set up a connection to the MS (steps 15 to 17).

It is much simpler to perform a **mobile originated call (MOC)** compared to a MTC (see Figure 4.9). The MS transmits a request for a new connection (1), the BSS forwards this request to the MSC (2). The MSC then checks if this user is allowed to set up a call with the requested service (3 and 4) and checks the availability of resources through the GSM network and into the PSTN. If all resources are available, the MSC sets up a connection between the MS and the fixed network.

Figure 4.10
Messageflowfor
MTC and MOC



In addition to the steps mentioned above, other messages are exchanged between an MS and BTS during connection setup (in either direction). These messages can be quite often heard in radios or badly shielded loudspeakers as crackling noise before the phone rings. Figure 4.10 shows the messages for an MTC and MOC. Paging is only necessary for an MTC, then similar message exchanges follow. The first step in this context is the channel access via the random access channel (RACH) with consecutive channel assignment; the channel assigned could be a traffic channel (TCH) or a slower signalling channel SDCCH.

The next steps, which are needed for communication security, comprise the authentication of the MS and the switching to encrypted communication. The system now assigns a TCH (if this has not been done). This has the advantage of only having to use an SDCCH during the first setup steps. If the setup fails, no TCH has been blocked. However, using a TCH from the beginning has a speed advantage.

The following steps depend on the use of MTC or MOC. If someone is calling the MS, it answers now with 'alerting' that the MS is ringing and with 'connect' that the user has pressed the connect button. The same actions

happen the other way round if the MS has initiated the call. After connection acknowledgement, both parties can exchange data.

Closing the connection comprises a user-initiated disconnect message (both sides can do this), followed by releasing the connection and the radio channel.

4.1.6 Handover

Cellular systems require **handover** procedures, as single cells do not cover the whole service area, but, e.g., only up to 35 km around each antenna on the countryside and some hundred meters in cities (Tripathi, 1998). The smaller the cell size and the faster the movement of a mobile station through the cells (up to 250 km/h for GSM), the more handovers of ongoing calls are required.

However, a handover should not cause a cut-off, also called **call drop**. GSM aims at maximum handover duration of 60 ms.

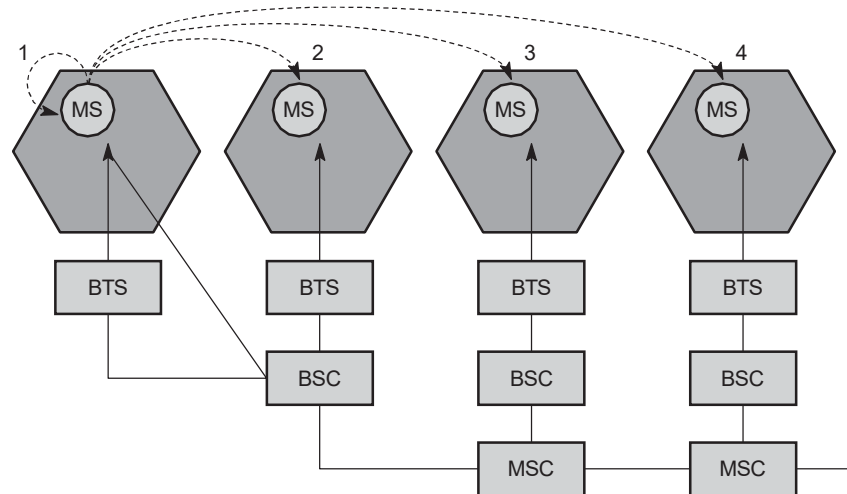
There are two basic reasons for a handover (about 40 have been identified in the standard):

- The mobile station **moves out of the range** of a BTS or a certain antenna of a BTS respectively. The received **signal level** decreases continuously until it falls below the minimal requirements for communication. The **error rate** may grow due to interference, the distance to the BTS may be too high (max. 35 km) etc. – all these effects may diminish the **quality of the radio link** and make radio transmission impossible in the near future.
- The wired infrastructure (MSC, BSC) may decide that the **traffic in one cell is too high** and shift some MS to other cells with a lower load (if possible). Handover may be due to **load balancing**.

Figure 4.11 shows four possible handover scenarios in GSM:

- **Intra-cell handover:** Within a cell, narrow-band interference could make transmission at a certain frequency impossible. The BSC could then decide to change the carrier frequency (scenario 1).
- **Inter-cell, intra-BSC handover:** This is a typical handover scenario. The mobile station moves from one cell to another, but stays within the control of the same BSC. The BSC then performs a handover, assigns a new radio channel in the new cell and releases the old one (scenario 2).
- **Inter-BSC, intra-MSC handover:** As a BSC only controls a limited number of cells; GSM also has to perform handovers between cells controlled by different BSCs. This handover then has to be controlled by the MSC (scenario 3). This situation is also shown in Figure 4.13.
- **Inter MSC handover:** A handover could be required between two cells belonging to different MSCs. Now both MSCs perform the handover together (scenario 4).

Figure 4.11
Types of handover
in GSM



To provide all the necessary information for a handover due to a weak link, MS and BTS both perform periodic measurements of the downlink and uplink quality respectively. (Link quality comprises signal level and bit error rate.) Measurement reports are sent by the MS about every half-second and contain the quality of the current link used for transmission as well as the quality of certain channels in neighboring cells (the BCCHs).

Figure 4.12 shows the typical behavior of the received signal level while an MS moves away from one BTS (BTS_{old}) closer to another one (BTS_{new}). In this case, the handover decision does not depend on the actual value of the received signal level, but on the average value. Therefore, the BSC collects all values (bit error rate and signal levels from uplink and downlink) from BTS and MS and calculates average values. These values are then compared to thresholds, i.e., the handover margin (HO_MARGIN), which includes some hysteresis to avoid a ping-pong effect (Wong, 1997). (Without hysteresis, even short-term interference, e.g., shadowing due to a building, could cause a handover.) Still, even with the HO_MARGIN, the ping-pong effect may occur in GSM – a value which is too high could cause a cut-off, and a value which is too low could cause too many handovers.

Figure 4.13 shows the typical signal flow during an inter-BSC, intra-MSC handover. The MS sends its periodic measurements reports, the BTS_{old} forwards these reports to the BSC_{old} together with its own measurements. Based on these values and, e.g., on current traffic conditions, the BSC_{old} may decide to perform a handover and sends the message HO_required to the MSC. The task of the MSC then comprises the request of the resources needed for the handover from the new BSC, BSC_{new} . This BSC checks if enough resources (typically frequencies or time slots) are available and activates a physical channel at the BTS_{new} to prepare for the arrival of the MS.

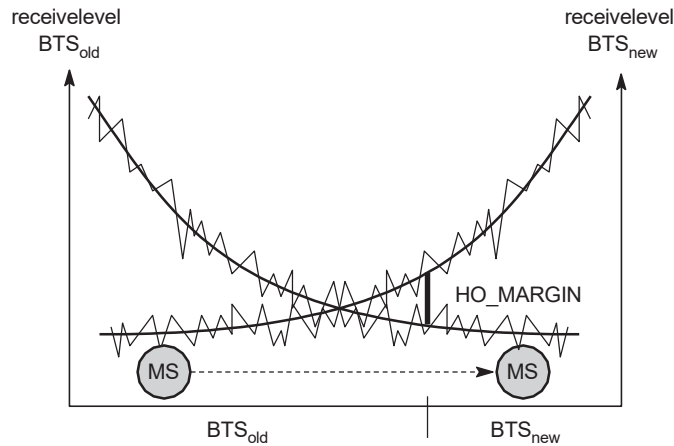


Figure 4.12 Handover decision depending on receive level

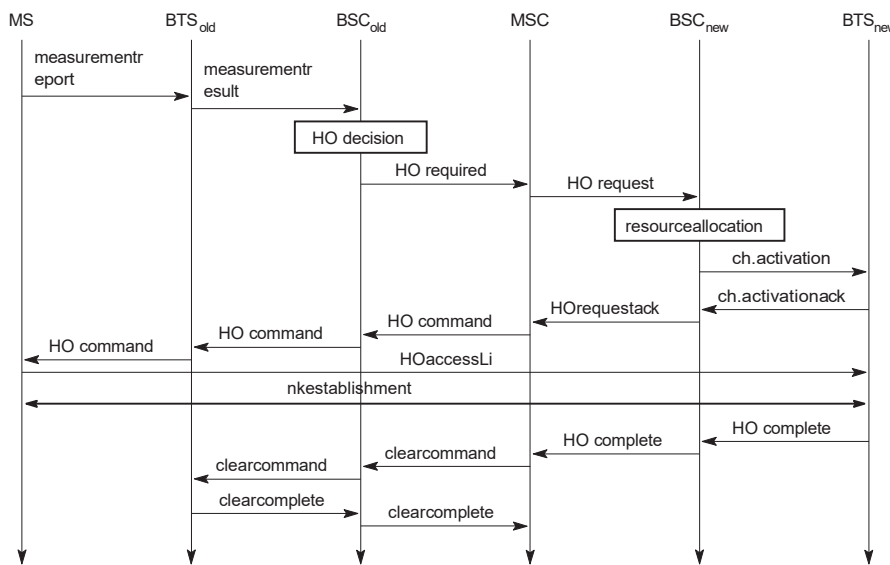


Figure 4.13 Intra-MSChandover

The BTS_{new} acknowledges the successful channel activation, BSC_{new} acknowledges the handover request. The MSC then issues a handover command that is forwarded to the MS. The MS now breaks its old radio link and accesses the new BTS. The next steps include the establishment of the link (this includes layer two link establishment and handover complete messages from the MS). Basically, the MS has then finished the handover, but it is important to release the resources at the old BSC and BTS and to signal the successful handover using the handover and clear complete messages as shown.

More sophisticated handover mechanisms are needed for seamless handovers between different systems. For example, future 3G networks will not cover whole countries but focus on cities and highways. Handover from,

e.g., UMTS to GSM without service interruption must be possible. Even more challenging is the seamless handover between wireless LANs (see chapter 7) and 2G/3G networks. This can be done using multimode mobile stations and a more sophisticated roaming infrastructure. However, it is still not obvious how these systems may scale for a large number of users and many handovers, and what handover quality guarantees they can give.

4.1.7 Security

GSM offers several security services using confidential information stored in the AuC and in the individual SIM (which is plugged into an arbitrary MS). The SIM stores personal, secret data and is protected with a PIN against unauthorized use. (For example, the secret key K_i used for authentication and encryption procedures is stored in the SIM.) These security services offered by GSM are explained below:

- **Access control and authentication:** The first step includes the authentication of a valid user for the SIM. The user needs a secret PIN to access the SIM. The next step is the subscriber authentication (see Figure 4.10). This step is based on a challenge-response scheme as presented in section 4.1.7.1.
- **Confidentiality:** All user-related data is encrypted. After authentication, BTS and MS apply encryption to voice, data, and signaling as shown in section 4.1.7.2. This confidentiality exists only between MS and BTS, but it does not exist end-to-end or within the whole fixed GSM/telephone network.
- **Anonymity:** To provide user anonymity, all data is encrypted before transmission, and user identifiers (which would reveal an identity) are not used over the air. Instead, GSM transmits a temporary identifier (TMSI), which is newly assigned by the VLR after each location update. Additionally, the VLR can change the TMSI at any time.

Three algorithms have been specified to provide security services in GSM. **Algorithm A3** is used for **authentication**, **A5** for **encryption**, and **A8** for the **generation of a cipher key**. In the GSM standard only algorithm A5 was publicly available, whereas A3 and A8 were secret, but standardized with open interfaces. Both A3 and A8 are no longer secret, but were published on the internet in 1998. This demonstrates that security by obscurity does not really work. As it turned out, the algorithms are not very strong. However, network providers can use stronger algorithms for authentication – or users can apply stronger end-to-end encryption. Algorithms A3 and A8 (or their replacements) are located on the SIM and in the AuC and can be proprietary. Only A5 which is implemented in the devices has to be identical for all providers.

4.1.7.1 Authentication

Before a subscriber can use any service from the GSM network, he or she must be authenticated. Authentication is based on the SIM, which stores the **individual authentication key K_i** , the **user identification IMSI**, and the algorithm used for authentication **A3**. Authentication uses a challenge-response method: the access

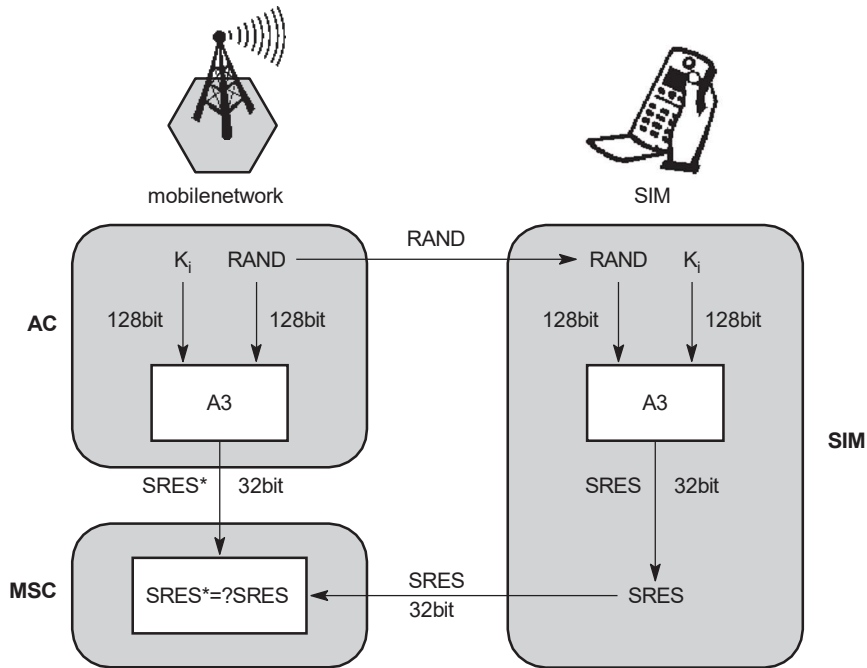


Figure 4.14
Subscriber
authentication

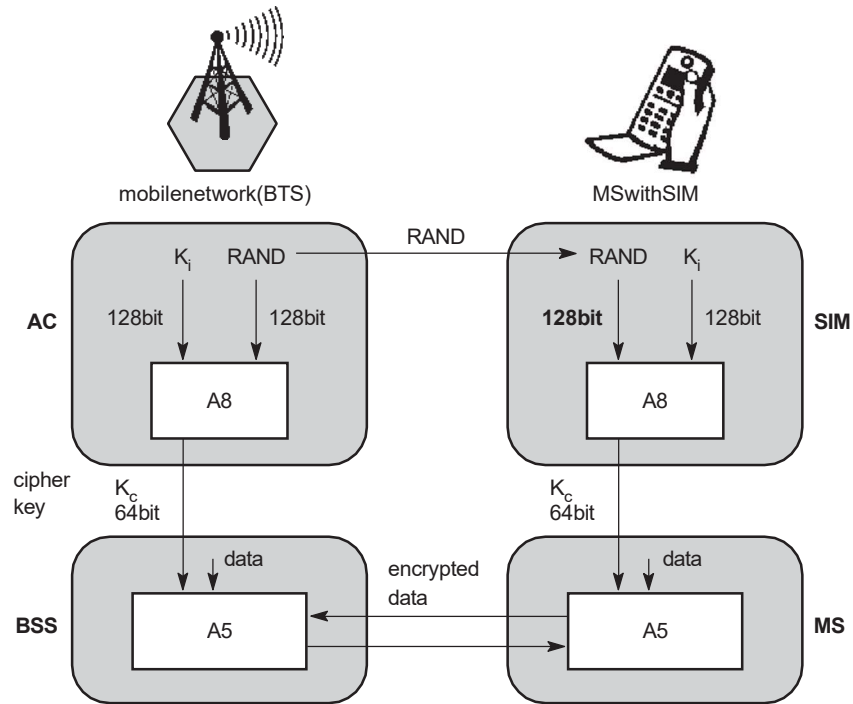
control AC generates a random number **RAND** as a challenge, and the SIM within the MS answers with **SRES** (signed response) as a response (see Figure 4.14). The AC performs the basic generation of random values **RAND**, signed responses **SRES**, and cipher keys K_c for each IMSI, and then forwards this information to the HLR. The current VLR requests the appropriate values for **RAND**, **SRES**, and K_c from the HLR.

For authentication, the VLR sends the random value **RAND** to the SIM. Both sides, network and subscriber module, perform the same operation with **RAND** and the key K_c , called **A3**. The MS sends back the **SRES** generated by the SIM; the VLR can now compare both values. If they are the same, the VLR accepts the subscriber, otherwise the subscriber is rejected.

4.1.7.2 Encryption

To ensure privacy, all messages containing user-related information are encrypted in GSM over the air interface. After authentication, MS and BSS can start using encryption by applying the cipher key K_c (the precise location of security functions for encryption, BTS and/or BSC are vendor dependent). K_c is generated using the individual key K_i and a random value by applying the algorithm **A8**. Note that the SIM in the MS and the network both calculate the same K_c based on the random value **RAND**. The key K_c itself is not transmitted over the air interface.

Figure 4.15
Data encryption



MS and BTS can now encrypt and decrypt data using the algorithm A5 and the cipher key K_c . As Figure 4.15 shows, K_c should be a 64 bit key – which is not very strong, but is at least a good protection against simple eavesdropping. However, the publication of A3 and A8 on the internet showed that in certain implementations 10 of the 64 bits are always set to 0, so that the real length of the key is thus only 54 consequently, the encryption is much weaker.

4.1.8 Newdataservices

As mentioned above, the standard bandwidth of 9.6 kbit/s (14.4 kbit/s with some providers) available for data transmission is not sufficient for the requirements of today's computers. When GSM was developed, not many people anticipated the tremendous growth of data communication compared to voice communication. At that time, 9.6 kbit/s was a lot, or at least enough for standard group 3 fax machines. But with the requirements of, e.g., web browsing, file download, or even intensive e-mail exchange with attachments, this is not enough.

To enhance the data transmission capabilities of GSM, two basic approaches are possible. As the basic GSM is based on connection-oriented traffic channels, e.g., with 9.6 kbit/s each, several channels could be combined to increase bandwidth. This system is called HSCSD and is presented in the following section. A

more progressive step is the introduction of packet-oriented traffic in GSM, i.e., shifting the paradigm from connections/telephone thinking to packets/internet thinking. The system, called GPRS, is presented in section 4.1.8.2.

4.1.8.1 HSCSD

A straightforward improvement of GSM's data transmission capabilities is **high speed circuit switched data (HSCSD)**, which is available with some providers. In this system, higher data rates are achieved by bundling several TCHs. An MS requests one or more TCHs from the GSM network, i.e., it allocates several TDMA slots within a TDMA frame. This allocation can be asymmetrical, i.e., more slots can be allocated on the downlink than on the uplink, which fits the typical user behavior of downloading more data compared to uploading. Basically, HSCSD only requires software upgrades in an MS and MSC (both have to be able to split a traffic stream into several streams, using a separate TCH each, and to combine these streams again).

In theory, an MS could use all eight slots within a TDMA frame to achieve an **air interface user rate (AIUR)** of, e.g., 8 TCH/F14.4 channels or 115.2 kbit/s (ETSI, 1998e). One problem of this configuration is that the MS is required to send and receive at the same time. Standard GSM does not require this capability – uplink and downlink slots are always shifted for three slots. ETSI (1997a) specifies the AIUR available at 57.6 kbit/s (duplex) using four slots in the uplink and downlink (Table 4.2 shows the permitted combinations of traffic channels and allocated slots for non-transparent services).

Although it appears attractive at first glance, HSCSD exhibits some major disadvantages. It still uses the connection-oriented mechanisms of GSM. These are not at all efficient for computer data traffic, which is typically bursty and asymmetrical. While downloading a larger file may require all channels reserved, typical web browsing would leave the channels idle most of the time. Allocating channels is reflected directly in the service costs, as, once the channels have been reserved, other users cannot use them.

AIUR	TCH/F4.8	TCH/F9.6	TCH/F14.4
4.8 kbit/s	1	-	-
9.6 kbit/s	2	1	-
14.4 kbit/s	3	-	1
19.2 kbit/s	4	2	-
28.8 kbit/s	-	3	2
38.4 kbit/s	-	4	-
43.2 kbit/s	-	-	3
57.6 kbit/s	-	-	4

Table 4.2 Available data rates for HSCSD in GSM

For n channels, HSCSD requires n times signaling during handover, connection setup and release. Each channel is treated separately. The probability of blocking or service degradation increases during handover, as in this case a BSC has to check resources for n channels, not just one. All in all, HSCSD may be an attractive interim solution for higher bandwidth and rather constant traffic (e.g., file download). However, it does not make much sense for bursty internet traffic as long as a user is charged for each channel allocated for communication.

4.1.8.2 GPRS

The next step toward more flexible and powerful data transmission avoids the problems of HSCSD by being fully packet-oriented. The **general packet radio service (GPRS)** provides packet mode transfer for applications that exhibit traffic patterns such as frequent transmission of small volumes (e.g., typical web requests) or infrequent transmissions of small or medium volumes (e.g., typical web responses) according to the requirement specification (ETSI, 1998a). Compared to existing data transfer services, GPRS should use the existing network resources more efficiently for packet mode applications, and should provide a selection of QoS parameters for the service requesters. GPRS should also allow for broadcast, multicast, and unicast service. The overall goal in this context is the provision of a more efficient and, thus, cheaper packet transfer service for typical internet applications that usually rely solely on packet transfer. Network providers typically support this model by charging on volume and not on connection time as is usual for traditional GSM data services and for HSCSD. The main benefit for users of GPRS is the 'always on' characteristic – no connection has to be set up prior to data transfer. Clearly, GPRS was driven by the tremendous success of the packet-oriented internet, and by the new traffic models and applications. However, GPRS, as shown in the following sections, needs additional network elements, i.e., software and hardware. Unlike HSCSD, GPRS does not only represent a software update to allow for the bundling of channels, it also represents a big step towards UMTS as the main internal infrastructure needed for UMTS (in its initial release) is exactly what GPRS uses (see section 4.4).

The main concepts of GPRS are as follows (ETSI, 1998b). For the new GPRS radio channels, the GSM system can allocate between one and eight time slots within a TDMA frame. Time slots are not allocated in a fixed, pre-determined manner but on demand. All time slots can be shared by the active users; up- and downlink are allocated separately. Allocation of the slots is based on current load and operator preferences. Depending on the coding, a transfer rate of up to 170 kbit/s is possible. For GPRS, operators often reserve at least a time slot per cell to guarantee a minimum data rate. The GPRS concept is independent of channel characteristics and of the type of channel (traditional GSM traffic or control channel), and does not limit the maximum data rate (only the GSM transport system limits the rate). All GPRS services can be used in parallel to conventional services. Table 4.3 shows the typical data rates available with GPRS if it is used together with GSM (GPRS can also be used for other TDMA systems).

Coding scheme	1slot	2slots	3slots	4slots	5slots	6slots	7slots	8slots
CS-1	9.05	18.2	27.15	36.2	45.25	54.3	63.35	72.4
CS-2	13.4	26.8	40.2	53.6	67	80.4	93.8	107.2
CS-3	15.6	31.2	46.8	62.4	78	93.6	109.2	124.8
CS-4	21.4	42.8	64.2	85.6	107	128.4	149.8	171.2

Table 4.3 GPRS data rates in kbit/s

In the beginning, only coding schemes CS-1 and CS-2 are available. The system chooses a coding scheme depending on the current error rate (CS-4 provides no error correction capabilities).

It should be noted that the real available data rate heavily depends on the current load of the cell as GPRS typically only uses idle time slots. The transfer rate depends on the capabilities of the MS as not all devices are able to send and receive at the same time. Table 4.4 gives examples for device classes together with their ability to use time slots for sending and receiving data. The maximum possible number of slots limits the transfer rate even more. For example, a class 12 device may receive data using 4 slots within a GSM time frame or it may send data using 4 slots. However, a maximum number of 5 slots may be used altogether. Using all 8 slots for data encoded using CS-4 yields the maximum rate of 171.2 kbit/s. Today, a typical MS is a class 10 device using CS-2, which results in a receiving rate of 53.6 kbit/s and a sending rate of 26.8 kbit/s.

In phase 1, GPRS offers a **point-to-point (PTP)** packet transfer service (ETSI, 1998c). One of the PTP versions offered is the **PTP connection oriented network service (PTP-CONS)**, which includes the ability of GPRS to maintain a virtual circuit upon change of the cell within the GSM network. This type of

Class	Receiving slots	Sending slots	Maximum number of slots
1	1	1	2
2	2	1	3
3	2	2	3
5	2	2	4
8	4	1	5
10	4	2	5
12	4	4	5

Table 4.4 Examples for GPRS device classes

Table 4.5 Reliability classes in GPRS according to ETSI (1998c)

Reliability class	Lost SDU probability	Duplicate SDU probability	Out of sequence SDU probability	Corrupt SDU probability
1	10^{-9}	10^{-9}	10^{-9}	10^{-9}
2	10^{-4}	10^{-5}	10^{-5}	10^{-6}
3	10^{-2}	10^{-5}	10^{-5}	10^{-2}

service corresponds to **X.25**, the typical circuit-switched packet-oriented transfer protocol available worldwide. The other PTP version offered is the **PTP connectionless network service (PTP-CLNS)**, which supports applications that are based on the Internet Protocol IP. Multicasting, called **point-to-multipoint (PTM)** service, is left for GPRS phase 2.

Users of GPRS can specify a **QoS-profile**. This determines the **service precedence** (high, normal, low), **reliability class** and **delay class** of the transmission, and **user data throughput**. GPRS should adaptively allocate radio resources to fulfill these user specifications. Table 4.5 shows the three reliability classes together with the maximum probabilities for a lost service data unit (SDU), a duplicated SDU, an SDU out of the original sequence, and the probability of delivering a corrupt SDU to the higher layer. Reliability class 1 could be used for very error-sensitive applications that cannot perform error corrections themselves. If applications exhibit greater error tolerance, class 2 could be appropriate. Finally, class 3 is the choice for error-insensitive applications or applications that can handle error corrections themselves.

Delay within a GPRS network is incurred by channel access delay, coding for error correction, and transfer delays in the fixed and wireless part of the GPRS network. The delay introduced by external fixed networks is out of scope. However, GPRS does not produce additional delay by buffering packets as store-and-forward networks do. If possible, GPRS tries to forward packets as fast as possible. Table 4.6 shows the specified maximum mean and 95 percentile delay values for packet sizes of 128 and 1,024 byte. As we can clearly see, no matter which class, all delays are orders of magnitude higher than fixed network delays. This is a very important characteristic that has to be taken into account when implementing higher layer protocols such as TCP on top of GPRS networks (see chapter 9). Typical round trip times (RTT) in fixed networks are in the order of 10 to 100 ms. Using real unloaded GPRS networks round trip times of well above 1 s for even small packets (128–512 byte) are common. Additionally, GPRS exhibits a large jitter compared to fixed networks (several 100 ms are not uncommon). This characteristic has a strong impact on user experience when, e.g., interactive Internet applications are used on top of GPRS.

Delay Class	SDU size 128 byte		SDU size 1,024 byte	
	Mean	95 percentile	Mean	95 percentile
1	<0.5s	<1.5s	<2s	<7s
2	<5s	<25s	<15s	<75s
3	<50s	<250s	<75s	<375s
4	Unspecified			

Table 4.6 Delay classes in GPRS according to ETSI (1998c)

Finally, GPRS includes several **security services** such as authentication, access control, user identity confidentiality, and user information confidentiality. Even a completely **anonymous service** is possible, as, e.g., applied for road toll systems that only charge a user via the MS independent of the user's identity.

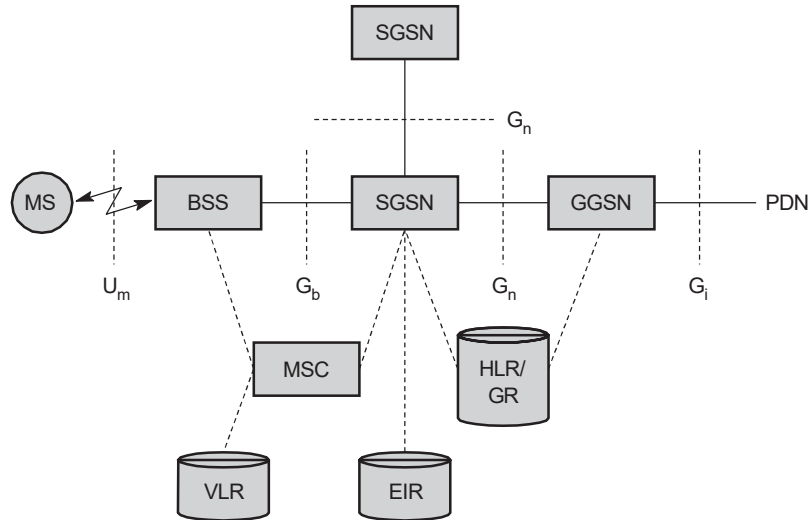
The **GPRS architecture** introduces two new network elements, which are called **GPRS support nodes (GSN)** and are in fact routers. All GSNs are integrated into the standard GSM architecture, and many new interfaces have been defined (see Figure 4.16). The **gateway GPRS support node (GGSN)** is the interworking unit between the GPRS network and external **packet data networks (PDN)**. This node contains routing information for GPRS users, performs address conversion, and tunnels data to a user via encapsulation. The GGSN is connected to external networks (e.g., IP or X.25) via the G_i interface and transfers packets to the SGSN via an IP-based GPRS backbone network (G_n interface).

The other new element is the **serving GPRS support node (SGSN)** which supports the MS via the G_b interface. The SGSN, for example, requests user addresses from the **GPRS register (GR)**, keeps track of the individual MSs' location, is responsible for collecting billing information (e.g., counting bytes), and performs several security functions such as access control. The SGSN is connected to a BSC via frame relay and is basically on the same hierarchy level as an MSC. The GR, which is typically a part of the HLR, stores all GPRS-relevant data. GGSNs and SGSNs can be compared with home and foreign agents, respectively, in a mobile IP network (see chapter 8).

As shown in Figure 4.16, packet data is transmitted from a PDN, via the GGSN and SGSN directly to the BSS and finally to the MS. The MSC, which is responsible for data transport in the traditional circuit-switched GSM, is only used for signaling in the GPRS scenario. Additional interfaces to further network elements and other PLMNs can be found in ETSI (1998b).

Before sending any data over the GPRS network, an MS must attach to it, following the procedures of the **mobility management**. The attachment procedure includes assigning a temporal identifier, called a **temporary logical link identity (TLLI)**, and a **ciphering key sequence number (CKSN)** for data encryption. For each MS, a **GPRS context** is set up and stored in the MS and in

Figure 4.16
GPRS architecture
reference model



the corresponding SGSN. This context comprises the status of the MS (which can be ready, idle, or standby; ETSI, 1998b), the CKSN, a flag indicating if compression is used, and routing data (TLLI, the routing area RA, a cell identifier, and a packet data channel, PDCH, identifier). Besides attaching and detaching, mobility management also comprises functions for authentication, location management, and ciphering (here, the scope of ciphering lies between MS and SGSN, which is more than in standard GSM). In **idle** mode an MS is not reachable and all context is deleted. In the **standby** state only movement across routing areas is updated to the SGSN but not changes of the cell. Permanent updating would waste battery power, no updating would require system-wide paging. The update procedure in standby mode is a compromise. Only in the **ready** state every movement of the MS is indicated to the SGSN.

Figure 4.17 shows the protocol architecture of the transmission plane for GPRS. Architectures for the signaling planes can be found in ETSI (1998b). All data within the GPRS backbone, i.e., between the GSNs, is transferred using the **GPRS tunneling protocol (GTP)**. GTP can use two different transport protocols, either the reliable **TCP** (needed for reliable transfer of X.25 packets) or the non-reliable **UDP** (used for IP packets). The network protocol for the GPRS backbone is **IP** (using any lower layers). To adapt to the different characteristics of the underlying networks, the **subnetwork dependent convergence protocol (SNDCP)** is used between an SGSN and the MS. On top of SNDCP and GTP, user packet data is tunneled from the MS to the GGSN and vice versa. To achieve a high reliability of packet transfer between SGSN and MS, a special LLC is used, which comprises ARQ and FEC mechanisms for PTP (and later PTM) services.

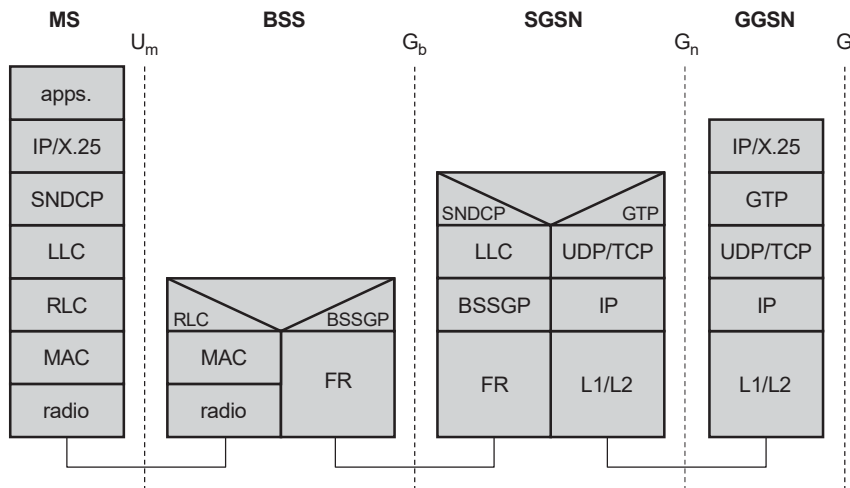


Figure 4.17
GPRS transmission
plane protocol
reference model

Abase station subsystem GPRS protocol (BSSGP) is used to convey routing and QoS-related information between the BSS and SGSN. BSSGP does not perform error correction and works on top of a **frame relay (FR)** network. Finally, radio link dependent protocols are needed to transfer data over the U_m interface. The **radio link protocol (RLC)** provides a reliable link, while the **MAC** controls access with signaling procedures for the radio channel and the mapping of LLC frames onto the GSM physical channels. The **radio interface** at U_m needed for GPRS does not require fundamental changes compared to standard GSM (Brasche, 1997), (ETSI, 1998d). However, several new logical channels and their mapping onto physical resources have been defined. For example, one MS can allocate up to eight **packet data traffic channels (PDTCHs)**. Capacity can be allocated on demand and shared between circuit-switched channels and GPRS. This allocation can be done dynamically with load supervision or alternatively, capacity can be pre-allocated.

A very important factor for any application working end-to-end is that it does not 'notice' any details from the GSM/GPRS-related infrastructure. The application uses, e.g., TCP on top of IP, IP packets are tunneled to the GGSN, which forwards them into the PDN. All PDNs forward their packets for a GPRS user to the GGSN, the GGSN asks the current SGSN for tunnel parameters, and forwards the packets via SGSN to the MS. Although MSs using GPRS may be considered as part of the internet, one should know that operators typically perform an address translation in the GGSN using NAT. All MSs are assigned private IP addresses which are then translated into global addresses at the GGSN. The advantage of this approach is the inherent protection of MSs from attacks (the subscriber typically has to pay for traffic even if it originates from an attack!) - private addresses are not routed through the internet so it is not possible to

reach an MS from the internet. This is also a disadvantage if an MS wants to offer a service using a fixed, globally visible IP address. This is difficult with IPv4 and NAT and it will be interesting to see how IPv6 is used for this purpose (while still protecting the MSs from outside attacks as air traffic is expensive).

4.2 DECT

Another fully digital cellular network is the **digital enhanced cordless telecommunications (DECT)** system specified by ETSI (2002, 1998j, k), (DECT Forum, 2002). Formerly also called **digital European cordless telephone and digital European cordless telecommunications**, DECT replaces older analog cordless phone systems such as CT1 and CT1+. These analog systems only ensured security to a limited extent as they did not use encryption for data transmission and only offered a relatively low capacity. DECT is also a more powerful alternative to the digital system CT2, which is mainly used in the UK (the DECT standard works throughout Europe), and has even been selected as one of the 3G candidates in the IMT-2000 family (see section 4.4). DECT is mainly used in offices, on campus, at trade shows, or in the home. Furthermore, access points to the PSTN can be established within, e.g., railway stations, large government buildings and hospitals, offering a much cheaper telephone service compared to a GSM system. DECT could also be used to bridge the last few hundred meters between a new network operator and customers. Using this 'small range' local loop, new companies can offer their service without having their own lines installed in the streets. DECT systems offer many different interworking units, e.g., with GSM, ISDN, or data networks. Currently, over 100 million DECT units are in use (DECT, 2002).

A big difference between DECT and GSM exists in terms of cell diameter and cell capacity. While GSM is designed for outdoor use with a cell diameter of up to 70 km, the range of DECT is limited to about 300 m from the base station (only around 50 m are feasible inside buildings depending on the walls). Due to this limited range and additional multiplexing techniques, DECT can offer its service to some 10,000 people within one km². This is a typical scenario within a big city, where thousands of offices are located in skyscrapers close together. DECT also uses base stations, but these base stations together with a mobile station are in a price range of €100 compared to several €10,000 for a GSM base station. GSM base stations can typically not be used by individuals for private networks. One reason is licensing as all GSM frequencies have been licensed to network operators. DECT can also handle handover, but it was not designed to work at a higher speed (e.g., up to 250 km/h like GSM systems). Devices handling GSM and DECT exist but have never been a commercial success.

DECT works at a frequency range of 1880–1990 MHz offering 120 full duplex channels. Time division duplex (TDD) is applied using 10 ms frames. The frequency range is subdivided into 10 carrier frequencies using FDMA, each frame being divided into 24 slots using TDMA. For the TDD mechanism,

12 slots are used as uplink, 12 slots as downlink (see Figure 3.4). The digital modulation scheme is GMSK—each station has an average transmission power of only 10 mW with a maximum of 250 mW.

4.2.1 System architecture

A DECT system, may have various different physical implementation depending on its actual use. Different DECT entities can be integrated into one physical unit; entities can be distributed, replicated etc. However, all implementations are based on the same logical reference model of the system architecture as shown in Figure 4.18. **Aglobal network** connects the local communication structure to the outside world and offers its services via the interface D_1 . Global networks could be integrated services digital networks (ISDN), public switched telephone networks (PSTN), public land mobile networks (PLMN), e.g., GSM, or packet switched public data network (PSPDN). These services offered by these networks include transportation of data and the translation of addresses and routing of data between the local networks.

Local networks in the DECT context offer local telecommunication services that can include everything from simple switching to intelligent call forwarding, address translation etc. Examples for such networks are analog or digital private branch exchanges (PBXs) or LANs, e.g., those following the IEEE 802.x family of LANs. As the core of the DECT system itself is quite simple, all typical network functions have to be integrated in the local or global network, where the databases **home data base (HDB)** and **visitor data base (VDB)** are also located. Both databases support mobility with functions that are similar to those in the HLR and VLR in GSM systems. Incoming calls are automatically forwarded to the current subsystem responsible for the DECT user, and the current VDB informs the HDB about changes in location.

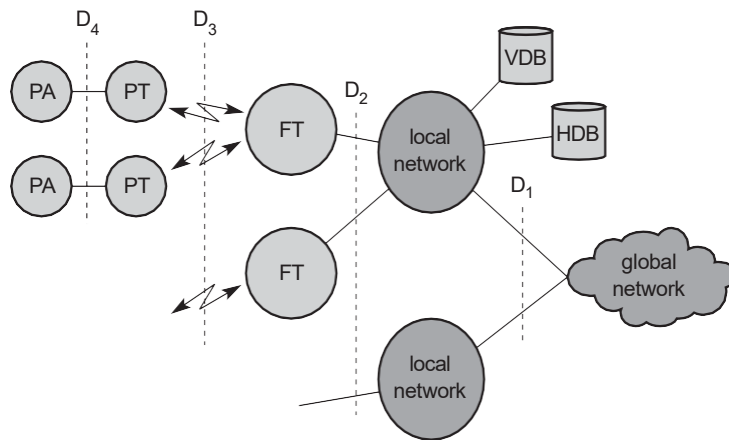


Figure 4.18
DECT system
architecture reference
model

The DECT core network consists of the **fixed radiotermination (FT)** and the **portable radiotermination (PT)**, and basically only provides a multiplexing service. FT and PT cover layers one to three at the fixed network side and mobile network side respectively. Additionally, several portable applications (PA) can be implemented on a device.

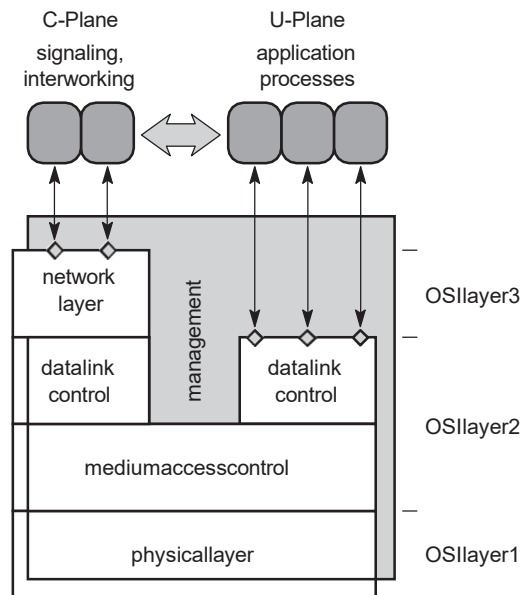
4.2.2 Protocol architecture

The DECT protocol reference architecture follows the OSI reference model. Figure 4.19 shows the layers covered by the standard: the physical layer, medium access control, and data link control⁸ for both the **control plane (C-Plane)** and the **user plane (U-Plane)**. An additional network layer has been specified for the C-Plane, so that user data from layer two is directly forwarded to the U-Plane. A management plane vertically covers all lower layers of a DECT system.

4.2.2.1 Physical layer

As in all wireless networks, the **physical layer** comprises all functions for modulation/demodulation, incoming signal detection, sender/receiver synchronization, and collection of status information for the management plane. This layer generates the physical channel structure with a certain, guaranteed throughput. On request from the MAC layer, the physical layer assigns a channel for data transmission.

Figure 4.19
DECT protocol
layers



⁸ Strictly speaking, the name "data link control" for the upper part of layer two is wrong in this architecture. According to the OSI reference model, the data link control (layer two) comprises the logical link control (layer 2b) and the medium access control (layer 2a).

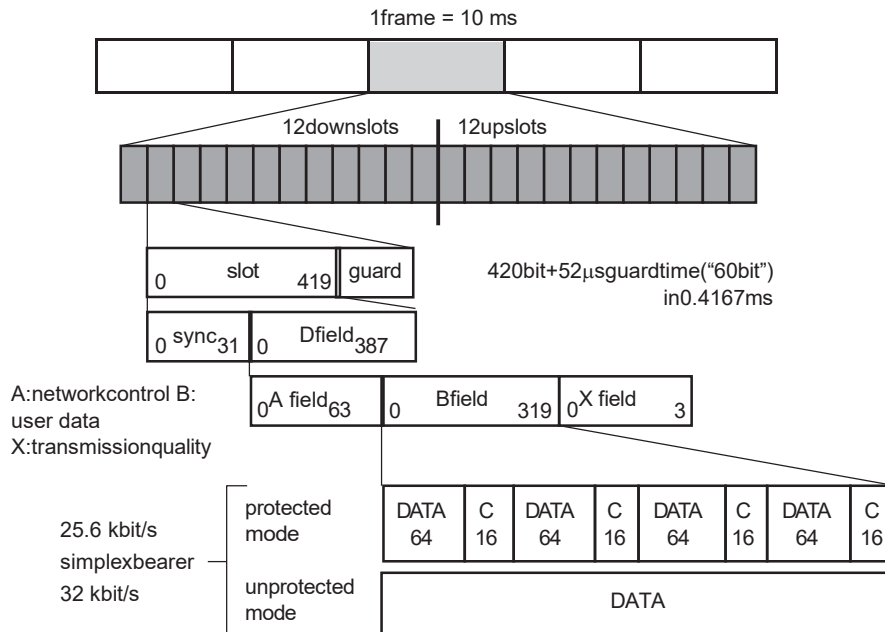


Figure 4.20
DECT multiplexed frame
structure

Figure 4.20 shows the standard TDMA frame structure used in DECT and some typical data packets. Each frame has a duration of 10 ms and contains 12 slots for the downlink and 12 slots for the uplink in the **basic connection** mode. If a mobile node receives data in slot s , it returns data in slot $s+12$. An **advanced connection** mode allows differential allocation schemes. Each slot has a duration of 0.4167 ms and can contain several different physical packets. Typically, 420 bits are used for data; the remaining 52 μ s are left as **guard space**. The 420 data bits are again divided into a 32 bits **synchronization pattern** followed by the **data** field D.

The fields for data transmission now use these remaining 388 bits for **network control** (A field), **user data** (B field), and the transfer of the **transmission quality** (X field). While network control is transmitted with a data rate of 6.4 kbit/s (64 bit each 10 ms), the user data rate depends on additional error correction mechanisms. The **simplex bearer** provides a data rate of 32 kbit/s in an **unprotected mode**, while using a 16 bit CRC **checksum** C for a data block of 64 bit in the **protected mode** reduces the data rate to 25.6 kbit/s. A **duplex bearer** service is produced by combining two simplex bearers. DECT also defines bearer types with higher throughputs by combining slots, e.g., the **double duplex bearer** offers 80 kbit/s full-duplex.

4.2.2.2 Mediumaccesscontrollayer

The **medium access control (MAC)** layer establishes, maintains, and releases channels for higher layers by activating and deactivating physical channels. MAC multiplexes several logical channels onto physical channels. Logical channels exist for signaling network control, user data transmission, paging, or sending broadcast messages. Additional services offered include segmentation/reassembly of packets and error control/error correction.

4.2.2.3 Datalinkcontrollayer

The **data link control (DLC)** layer creates and maintains reliable connections between the mobile terminal and the base station. Two services have been defined for the **C-Plane**: a **connectionless broadcast** service for paging (called **Lb**) and a **point-to-point** protocol similar to LAPD in ISDN, but adapted to the underlying MAC (called **LAPC+Lc**).

Several services exist for the **U-Plane**, e.g., a transparent unprotected service (basically a null service), a forward error correction service, rate adaptation services, and services for future enhancements. If services are reused, e.g., to transfer ISDN data at 64 kbit/s, then DECT also tries to transfer 64 kbit/s. However, in case of errors, DECT raises the transfer rate to 72 kbit/s, and includes FEC and a buffer for up to eight blocks to perform ARQ. This buffer then introduces an additional delay of up to 80 ms.

4.2.2.4 Network layer

The **network layer** of DECT is similar to those in ISDN and GSM and only exists for the **C-Plane**. This layer provides services to request, check, reserve, control, and release resources at the fixed station (connection to the fixed network, wireless connection) and the mobile terminal (wireless connection). The **mobility management (MM)** within the network layer is responsible for identity management, authentication, and the management of the location data bases. **Call control (CC)** handles connection setup, release, and negotiation. Two message services, the **connection oriented message service (COMS)** and the **connectionless message service (CLMS)** transfer data to and from the interworking unit that connects the DECT system with the outside world.

4.3 TETRA

Trunked radio systems constitute another method of wireless data transmission. These systems use many different radio carriers but only assign a specific carrier to a certain user for a short period of time according to demand. While, for example, taxi services, transport companies with fleet management systems and rescue teams all have their own unique carrier frequency in traditional systems, they can share a whole group of frequencies in trunked radio systems for better frequency reuse via FDMA and TDM techniques. These types of radio systems typically offer

interfaces to the fixed telephone network, i.e., voice and data services, but are not publicly accessible. These systems are not only simpler than most other networks, they are also reliable and relatively cheap to set up and operate, as they only have to cover the region where the local users operate, e.g., a city taxi service.

To allow a common system throughout Europe, ETSI standardized the **TETRA** system (**terrestrial trunked radio**)⁹ in 1991 (ETSI, 2002), (TETRAMoU, 2002). This system should replace national systems, such as MODACOM, MOBILTEX and COGNITO in Europe that typically connect to an X.25 packet network. (An example system from the US is ARDIS.) TETRA offers two standards: the **Voice+Data (V+D)** service (ETSI, 1998) and the **packet data optimized (PDO)** service (ETSI, 1998). While V+D offers circuit-switched voice and data transmission, PDO only offers packet data transmission, either connection-oriented to connect to X.25 or connectionless for the ISO CLNS (connectionless network service). The latter service can be point-to-point or point-to-multipoint, the typical delay for a short message (128 byte) being less than 100 ms. V+D connection modes comprise unicast and broadcast connections, group communication within a certain protected group, and a direct ad hoc mode without a base station. However, delays for short messages can be up to 500 ms or higher depending on the priority.

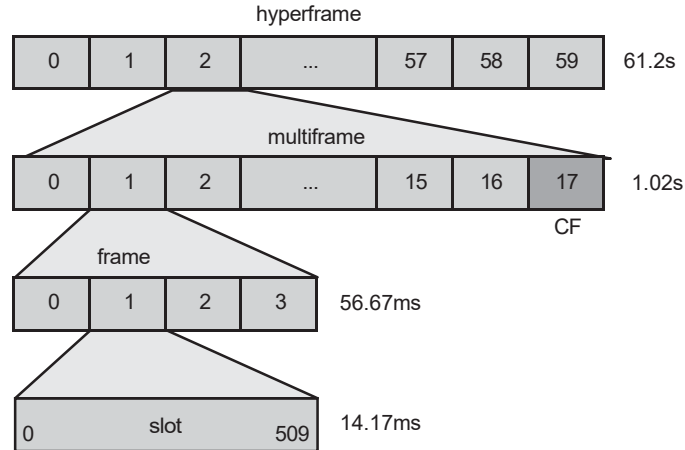
TETRA also offers bearer services of up to 28.8 kbit/s for unprotected data transmission and 9.6 kbit/s for protected transmission. Examples for end-to-end services are call forwarding, call barring, identification, call hold, call priorities, emergency calls and group joins. The system architecture of TETRA is very similar to GSM. Via the radio interface U_m , the **mobile station (MS)** connects to the **switching and management infrastructure (SwMI)**, which contains the user data bases (HDB, VDB), the base station, and interfaces to PSTN, ISDN, or PDN. The system itself, however, is much simpler in real implementation compared to GSM, as typically no handover is needed. Taxis usually remain within a certain area which can be covered by one TETRA cell.

Several frequencies have been specified for TETRA which uses FDD (e.g., 380–390 MHz uplink/390–400 MHz downlink, 410–420 MHz uplink/420–430 MHz downlink). Each channel has a bandwidth of 25 kHz and can carry 36 kbit/s. Modulation is DQPSK. While V+D uses up to four TDMA voice or data channels per carrier, PDO performs statistical multiplexing. For accessing a channel, slotted Aloha is used.

Figure 4.21 shows the typical **TDMA frame structure** of TETRA. Each **frame** consists of four slots (four channels in the V+D service per carrier), with a frame duration of 56.67 ms. Each **slot** carries 510 bits within 14.17 ms, i.e., 36 kbit/s. 16 frames together with one **control frame (CF)** form a **multiframe**, and finally, a **hyperframe** contains 60 multiframe. To avoid sending and receiving at the same time, TETRA shifts the uplink for a period of two slots compared to the downlink.

⁹ Formerly known as trans-European trunked radio, but worldwide marketing is better without "Europe" in the name (see DECT).

Figure 4.21
TETRA frame
structure



TETRA offers **traffic channels (TCH)** and **control channels (CCH)** similar to GSM. Typical TCHs are TCH/S for voice transmission, and TCH/7.2, TCH/4.8, TCH/2.4 for data transmission (depending on the FEC mechanisms required).

However, in contrast to GSM, TETRA offers additional services like group call, acknowledged group call, broadcast call, and discreet listening. Emergency services need a sub-second group-call setup in harsh environments which possibly lack all infrastructure. These features are currently not available in GSM or other typical mobile telephony networks, so TETRA is complementary to other systems. TETRA has been chosen by many government organizations in Europe and China.

4.4 UMTS and IMT-2000

A lot has been written about third generation (or 3G) networks in the last few years. After a lot of hype and frustration these networks are currently deployed in many countries around the world. But how did it all start? First of all, the International Telecommunication Union (ITU) made a request for proposals for radio transmission technologies (RTT) for the **international mobile telecommunications (IMT) 2000** program (ITU, 2002), (Callendar, 1997), (Shafi, 1998). IMT-2000, formerly called future public land mobile telecommunication system (FPLMTS), tried to establish a common worldwide communication system that allowed for terminal and user mobility, supporting the idea of universal personal telecommunication (UPT). Within this context, ITU has created several recommendations for FPLMTS systems, e.g., network architectures for FPLMTS (M.817), Requirements for the Radio Interface(s) for FPLMTS (M.1034), or Framework for Services Supported by FPLMTS (M.816). The number 2000 in IMT-2000 should indicate the start of the system (year 2000+x) and the spec-

trum used (around 2000 MHz). IMT-2000 includes different environments such as indoor use, vehicles, satellites and pedestrians. The world radio conference (WRC) 1992 identified 1885–2025 and 2110–2200 MHz as the frequency bands that should be available worldwide for the new IMT-2000 systems (Recommendation ITU-R M.1036). Within these bands, two times 30 MHz have been reserved for mobile satellite services (MSS).

Figure 4.22 shows the ITU frequency allocation (from the world administrative radio conference, 1992) together with examples from several regions that already indicate the problem of worldwide common frequency bands. In Europe, some parts of the ITU's frequency bands for IMT-2000 are already allocated for DECT (see section 4.2). The remaining frequencies have been split into bands for UTRA-FDD (uplink: 1920–1980 MHz, downlink: 2110–2170 MHz) and UTRA-TDD (1900–1920 MHz and 2010–2025 MHz). The technology behind UTRA-FDD and -TDD will subsequently be explained in more detail as they form the basis of UMTS. Currently, no other system is planned for IMT-2000 in Europe. More bandwidth is available in China for the Chinese 3G system TD-SCDMA or possibly other 3G technologies (such as W-CDMA or cdma2000 – it is still open which system will dominate the Chinese market; Chen, 2002). Against slightly different frequencies are used by the 3G services in Japan, which are based on W-CDMA (like UTRA-FDD) or cdma2000. An open question is the future of 3G in the US as the ITU's frequency bands have already been allocated for 2G networks or are reserved for other use. In addition to the original frequency allocations, the world radio conference (WRC) allocated new terrestrial IMT-2000 bands in the range of 800–1000 MHz, 1700–1900 MHz and 2500–2700 MHz in 2000. This approach includes the reuse of 2G spectrum (Evci, 2001).

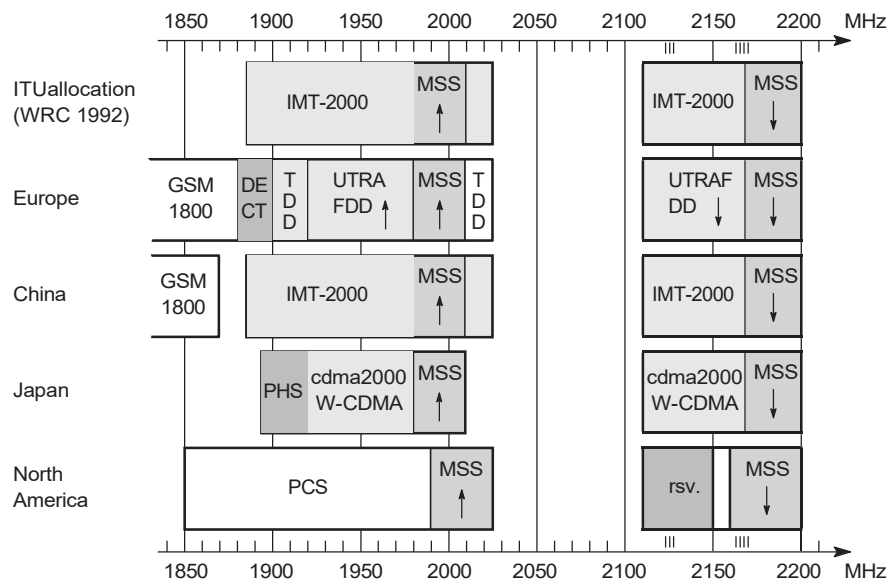


Figure 4.22
IMT-2000 frequencies

Now the reader might be confused by all the different technologies mentioned in the context of IMT-2000. Wasn't the plan to have a common global system? This was the original plan, but after many political discussions and fights about patents this idea was dropped and a so-called family of 3G standards was adopted.

For the RTT, several proposals were received in 1998 for indoor, pedestrian, vehicular, and satellite environments. These came from many different organizations, e.g., **UWC-136** from the Universal Wireless Communications Consortium (US) that extends the IS-136 standard into the third generation systems, **cdma2000** that is based on the IS-95 system (US), and wideband packet CDMA (WP-CDMA) which tries to align to the European UTRA proposal. Basically, three big regions were submitting proposals to the ITU: ETSI for Europe, ARIB (Association of Radio Industries and Broadcasting) and TTC (Telecommunications Technology Council) for Japan, and ANSI (American National Standards Institute) for the US.

The European proposal for IMT-2000 prepared by ETSI is called **universal mobile telecommunications system (UMTS)** (Dasilva, 1997), (Ojanperä, 1998), the specific proposal for the radio interface RTT is **UMTS (now: universal) terrestrial radio access (UTRA)** (ETSI, 1998n), (UMTS Forum, 2002). UMTS as initially proposed by ETSI rather represents an evolution from the second generation GSM system to the third generation than a completely new system. In this way, many solutions have been proposed for a smooth transition from GSM to UMTS, saving money by extending the current system rather than introducing a new one (GSMMoU, 1998).

One initial enhancement of GSM toward UMTS was **enhanced data rates for global (or: GSM) evolution (EDGE)**, which uses enhanced modulation schemes (8 PSK instead of GSM's GMSK, see chapter 2) and other techniques for data rates of up to 384 kbit/s using the same 200 kHz wide carrier and the same frequencies as GSM (i.e., a data rate of 48 kbit/s per time slot is available). EDGE can be introduced incrementally offering some channels with EDGE enhancement that can switch between EDGE and GSM/GPRS. In Europe, EDGE was never used as a step toward UMTS but operators directly jumped onto UMTS. However, EDGE can also be applied to the US IS-136 system and may be a choice for operators that want to enhance their 2G systems (3G Americas, 2002).

Besides enhancing data rates, new additions to GSM, like **customized application for mobile enhanced logic (CAMEL)** introduce intelligent network support. This system supports, for example, the creation of a **virtual home environment (VHE)** for visiting subscribers. GSMMoU (1998) provides many proposals covering QoS aspects, roaming, services, billing, accounting, radio aspects, core networks, access networks, terminal requirements, security, application domains, operation and maintenance, and several migration aspects.

UMTS fits into a bigger framework developed in the mid-nineties by ETSI, called **global multimedia mobility (GMM)**. GMM provides an architecture to integrate mobile and fixed **terminals**, many different **access networks** (GSM BSS, DECT, ISDN, UMTS, LAN, WAN, CATV, MBS), and several **core transport**

networks(GSM+IN,ISDN+IN,B-ISDN+TINA,TCP/IP)(ETSI,2002).Within this framework,ETSI developed **basic requirements** for UMTS and for UTRA, the radio interface (ETSI, 1998h). Key requirements are minimum data rates of 144 kbit/s for rural outdoor access (with the goal of 384 kbit/s) at a maximum speed of 500 km/h.¹⁰ For suburban outdoor use a minimum of 384 kbit/s should be achieved with the goal of 512 kbit/s at 120 km/h. For indoor or city use with relatively short ranges, up to 2 Mbit/s are required at 10 km/h (walking).

UMTS should also provide several bearer services, real-time and non real-time services, circuit and packet switched transmission, and many different data rates. Handovers should be possible between UMTS cells, but also between UMTS and GSM or satellite networks. The system should be compatible with GSM, ATM, IP, and ISDN-based networks. To reflect the asymmetric bandwidth needs of typical users, UMTS should provide a variable division of uplink and down-link data rates. Finally, UMTS has to fit into the IMT-2000 framework (this is probably the decisive factor for its success). As the global UMTS approach is rather ambitious, a more realistic alternative for the initial stages would be UMTS cells in cities providing a subset of services.

Several companies and interest groups have handed in proposals for UTRA (ETSI, 1998i), of which ETSI selected two for UMTS in January 1998. For the **paired band** (using FDD as a duplex mechanism), ETSI adopted the **W-CDMA** (Wideband CDMA) proposal, for the **unpaired band** (using TDD as duplex mechanism) the **TD-CDMA** (Time Division CDMA) proposal is used (Adachi, 1998), (Dahlman, 1998), (ETSI, 1998n). The paired band is typically used for public mobile network providers (wide area, see GSM), while the unpaired band is often used for local and indoor communication (see DECT). The following sections will present key properties of the initial UMTS system.

What happened to the IMT-2000 family? Figure 4.23 gives an overview. As a single standard could not be found, the ITU standardized five groups of 3G radio access technologies.

- **IMT-DS:** The **direct spread** technology comprises wideband CDMA (**W-CDMA**) systems. This is the technology specified for UTRA-FDD and used by all European providers and the Japanese NTT DoCoMo for 3G wide area services. To avoid complete confusion ITU's name for the technology is IMT-DS, ETSI called it UTRA-FDD in the UMTS context, and technology used is called W-CDMA (in Japan this is promoted as FOMA, freedom of mobile multimedia access). Today, standardization of this technology takes place in 3GPP (Third generation partnership project, 3GPP, 2002a). Section 4.4.1 provides more detail about the standardization process.
- **IMT-TC:** Initially, this family member, called **time code**, contained only the UTRA-TDD system which uses time-division CDMA (**TD-CDMA**). Later on, the Chinese proposal, TD-synchronous CDMA (**TD-SCDMA**) was added.

¹⁰ This speed is a problem as currently, only DAB can provide higher bit rates at high speeds.

Both standards have been combined and 3GPP fosters the development of this technology. It is unclear when and to what extent this technology will be introduced. The initial UMTS installations are based on W-CDMA.

- **IMT-MC:** cdma2000 is a **multi-carrier** technology standardized by 3GPP2 (Third generation partnership project 2, 3GPP2, 2002), which was formed shortly after 3GPP to represent the second main stream in 3G technology. Version cdma2000 EV-DO has been accepted as the 3G standard.
- **IMT-SC:** The enhancement of the US TDMA systems, UWC-136, is a **single carrier** technology originally promoted by the Universal Wireless Communications Consortium (UWCC). It is now integrated into the 3GPP efforts. This technology applies EDGE, among others, to enhance the 2G IS-136 standard.
- **IMT-FT:** As **frequency time** technology, an enhanced version of the cordless telephone standard DECT has also been selected for applications that do not require high mobility. ETSI is responsible for the standardization of DECT.

The main driving forces in the standardization process are 3GPP and 3GPP2. ETSI has moved its GSM standardization process to 3GPP and plays a major role there. 3GPP tends to be dominated by European and Japanese manufacturers and standardization bodies, while 3GPP2 is dominated by the company Qualcomm and CDMA network operators. The quarrels between Qualcomm and European manufacturers (e.g., Nokia, Ericsson) regarding CDMA patents (UMTS and cdma2000 use CDMA) even escalated into the political arena back in 1998 (US vs EU). Everything cooled down when hundreds of patents had been exchanged and the systems had been harmonized (e.g., CDMA chipping rates).

Figure 4.23
The IMT-2000
family

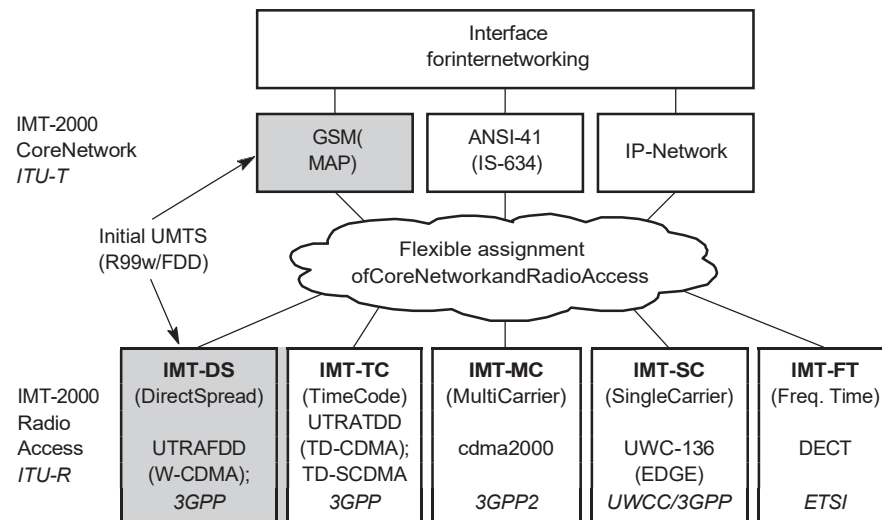


Figure 4.23 shows more than just the radio access technologies. One idea of IMT-2000 is the flexible assignment of a core network to a radio access system. The classical core network uses SS7 for signaling which is enhanced by ANSI-41 (cdmaOne, cdma2000, TDMA) or MAP (GSM) to enable roaming between different operators. The evolution toward 4G systems is indicated by the use of all-IP core networks (see Chapter 11). Obviously, internet-working functions have to be provided to enable cross-system data transfer, roaming, billing etc.

4.4.1 UMTS releases and standardization

UMTS as discussed today and introduced in many countries relies on the initial release of the UMTS standard called **release 99** or **R99** for short. This release of the specification describes the new radio access technologies UTRA FDD and UTRA TDD, and standardizes the use of a GSM/GPRS network as core within 440 separate specifications. This enables a cost effective migration from GSM to UMTS. The initial installations will even offer the FDD mode only as indicated in Figure 4.23. This release was (almost) finalized in 1999 – hence the name R99. The following sections will focus on this release as it is unclear when, and to what extent, the following releases will be realized.

After R99 the release 2000 or R00 followed. However, in September 2000 3GPP realized that it would be impossible to finalize the standard within the year 2000. 3GPP decided to split R2000 into two standards and call them release 4 (Rel-4) and release 5 (Rel-5). The version of all standards finalized for R99 start with 3.x.y (a reason for renaming R99 into Rel-3), Rel-4 and Rel-5 versions start with 4.x.y and 5.x.y, respectively. The standards are grouped into series. For example, radio aspects are specified in series 25, technical realization in series 23, and codecs in series 26. The complete standard number (e.g., TS 25.401 V3.10.0) then identifies the series (25), the standard itself (401), the release (3), and the version within the release (10.0). All standards can be downloaded from www.3gpp.org (the example given is the UTRAN overall description, release 99, from June 2002).

Release 4 introduces quality of service in the fixed network plus several execution environments (e.g., MExE, mobile execution environment, see chapter 10) and new service architectures. Furthermore, the Chinese proposal, TD-SCDMA was added as low chip rate option to UTRA-TDD (only 1.28 Mchip/s occupying only 1.6 MHz bandwidth). This release already consists of over 500 specifications and was frozen in March 2001.

Release 5 specifies a radically different core network. The GSM/GPRS based network will be replaced by an almost all-IP-core. While the radio interfaces remain the same, the changes in the core are tremendous for telecommunication network operators who have used traditional telephone technologies for many years. The content of this specification was frozen March 2002. This standard integrates IP-based multimedia services (IMS) controlled by the IETF's session initiation protocol (SIP, RFC 3261; Rosenberg, 2002; SIP Forum, 2002). A high speed downlink packet access (HSDPA) with speeds in the order of

8–10 Mbit/s was added as well as a wideband 16 kHz AMR codec for better audio quality. Additional features are end-to-end QoS messaging and several data compression mechanisms.

3GPP is currently working on **release 6** (and thinking of release 7) which is expected to be frozen in March 2003. This release comprises the use of multiple input multiple output (MIMO) antennas, enhanced MMS, security enhancements, WLAN/UMTS interworking, broadcast/multicast services, enhanced IMS, IP emergency calls, and many more management features (3GPP, 2002a).

The reader should not forget that many companies still have to make any money from, release 99, so it is not clear at what time and to what extent the new releases will be implemented. The following describes the initial UMTS standard, release 99, which is currently deployed.

4.4.2 UMTS system architecture

Figure 4.24 shows the very simplified UMTS reference architecture which applies to both UTRA solutions (3GPP, 2000). The **UTRA network (UTRAN)** handles cell level mobility and comprises several **radio network subsystems (RNS)**. The functions of the RNS include radio channel ciphering and deciphering, hand-

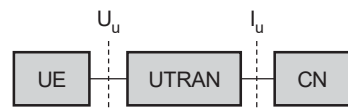
over control, radio resource management etc. The UTRAN is connected to the **user equipment (UE)** via the radio interface U_u (which is comparable to the U_m interface in GSM). Via the I_u interface (which is similar to the A interface in GSM),

UTRAN communicates with the **core network (CN)**. The CN contains functions for inter-system handover, gateways to other networks (fixed or wireless), and performs location management if there is no dedicated connection between UE and UTRAN.

UMTS further subdivides the above simplified architecture into so-called **domains** (see Figure 4.25). The **user equipment domain** is assigned to a single user and comprises all the functions that are needed to access UMTS services. Within this domain are the **USIM domain** and the **mobile equipment domain**. The **USIM domain** contains the SIM for UMTS which performs functions for encryption and authentication of users, and stores all the necessary user-related data for UMTS. Typically, this USIM belongs to a service provider and contains a micro processor for an enhanced program execution environment (USAT, UMTS SIM application toolkit). The end device itself is in the **mobile equipment domain**. All functions for radio transmission as well as user interfaces are located here.

The **infrastructure domain** is shared among all users and offers UMTS services to all accepted users. This domain consists of the **access network domain**, which contains the radio access networks (RAN), and the **core network domain**, which contains access network independent functions. The **core network domain** can be separated into three domains with specific tasks. The **servicing network domain** comprises all functions currently used by a user for accessing

Figure 4.24
Main components
of the UMTS
reference
architecture



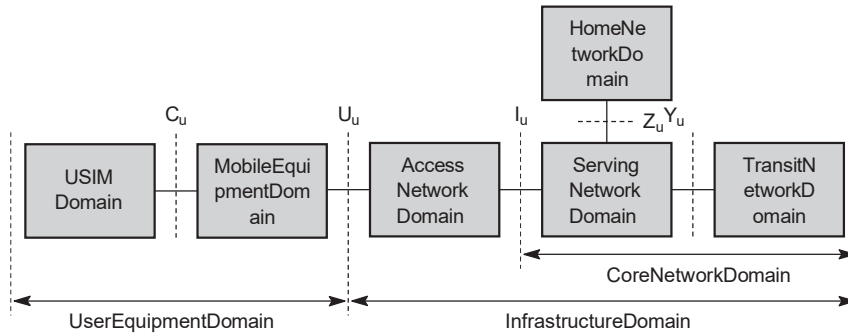


Figure 4.25
UMTS domains
and interfaces

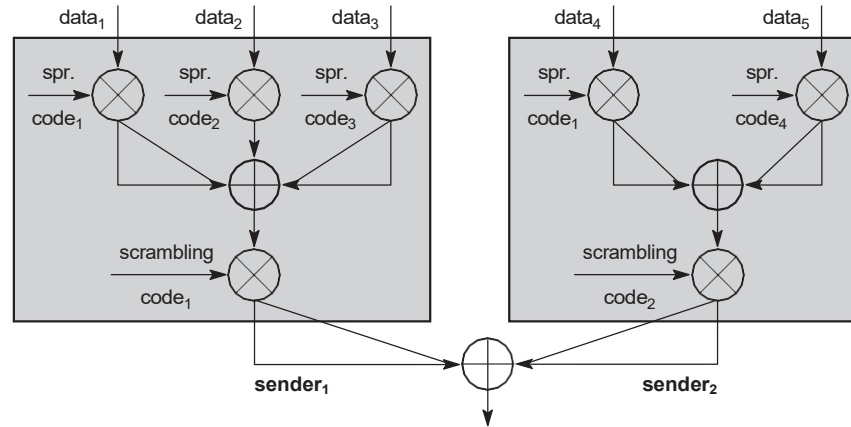
UMTS services. All functions related to the home network of a user, e.g., user data look-up, fall into the **home network** domain. Finally, the **transit network** domain may be necessary if, for example, the serving network cannot directly contact the home network. All three domains within the core network may in fact be the same physical network. These domains only describe functionalities.

4.4.3 UMTS radio interface

The biggest difference between UMTS and GSM comes with the new radio interface (U_u). The duplex mechanisms are already well known from GSM (FDD) and DECT (TDD). However, the direct sequence (DS) CDMA used in UMTS is new (for European standards, not in the US where CDMA technology has been available since the early nineties). DS-SS was introduced in chapters 2 and 3. This technology multiplies a stream of bits with a chipping sequence. This spreads the signal and, if the chipping sequence is unique, can separate different users. All signals use the same frequency band (in UMTS/IMT-2000 5 MHz-wide bands have been specified and licensed to network operators). To separate different users, the codes used for spreading should be (quasi) orthogonal, i.e., their cross-correlation should be (almost) zero.

UMTS uses a constant **chipping rate** of 3.84 Mchip/s. Different user data rates can be supported using different spreading factors (i.e., the number of chips per bit). Figure 4.26 shows the basic ideas of spreading and separation of different senders in UMTS. The first step in a sender is spreading of user data (data_u) using orthogonal **spreading** codes. Using orthogonal codes separate the different data streams of a sender. UMTS uses so-called **orthogonal variable spreading factor (OVSF)** codes. Figure 4.27 shows the basic idea of OVSF. Orthogonal codes are generated by doubling a chipping sequence X with and without flipping the sign of the chips. This results in X and $-X$, respectively. Doubling the chipping sequence also results in spreading a bit twice as much as before. The spreading factor $SF = nb$ becomes $2n$. Starting with a spreading factor of 1, Figure 4.27 shows the generation of orthogonal codes with different spreading factors. Two codes are orthogonal as long as one code is never a part of the other code. Looking at the coding tree in Figure 4.27 and considering the construction of the codes, orthogonality is guaranteed if one code has not been generated based on another. For example, if a sender uses the code $(1, -1)$

Figure 4.26
Spreading and scrambling of user data



with spreading factor 2, it is not allowed to use any of the codes located in the subtrees generated out of $(1,-1)$. This means that, e.g., $(1,-1,1,-1)$, $(1,-1,-1,1,1,-1,1,-1)$, or $(1,-1,-1,1,-1,1,1,-1,-1,1,1,-1,1,-1,-1,1)$ cannot be used anymore. However, it is no problem to use codes with different spreading factors if one code has not been generated using the other. Thus, $(1,-1)$ block only the lower subtree in Figure 4.27, many other codes from the upper part can still be used. An example for a valid combination in OVFS is $(1,-1), (1,1,-1,-1), (1,1,1,1,1,1,1,1), (1,1,1,1,-1,-1,-1,-1,1,1,1,1,-1,-1,-1,-1), (1,1,1,1,-1,-1,-1,-1,-1,-1,-1,1,1,1,1)$. This combination occupies the whole code spaces and allows for the transmission of data with different spreading factors (2, 4, 8, and $2 \cdot 16$). This example shows the tight coupling of available spreading factors and orthogonal codes.

Now remember that UMTS uses a constant chipping rate (3.84 Mchip/s). Using different spreading factors this directly translates into the support of different data rates. If the chipping rate is constant, doubling the spreading factor means dividing the data rate by two. But this also means that UMTS can only support a single data stream with SF=1 as then no other code may be used. Using the example combination above, a stream with half the maximum data rate, one with a fourth, one with an eighth, and two with a sixteenth are supported at the same time.

Each sender uses OVFS to spread its data streams as Figure 4.26 shows. The spreading codes chosen in the senders can be the same. Using different spreading codes in all senders within a cell would require a lot of management and would increase the complexity. After spreading all chip streams are added and scrambled. **Scrambling** does not spread the chip sequence any further but XORs chips based on a code. In the **FDD** mode, this scrambling code is unique for each sender and separates all senders (UE and base station) in a cell. After scrambling, the signals of different senders are quasi orthogonal. Quasi-orthogonal signals have the nice feature that they stay quasi-orthogonal even if they are not synchronized. Using orthogonal codes would require chip-synchronous reception and tight synchronization (this is done in other CDMA networks). For **TDD** the scrambling code is cell specific, i.e., all stations in a cell use the same scrambling code and cells are separated using different codes. The scrambled chips are **QPSK** modulated and transmitted.

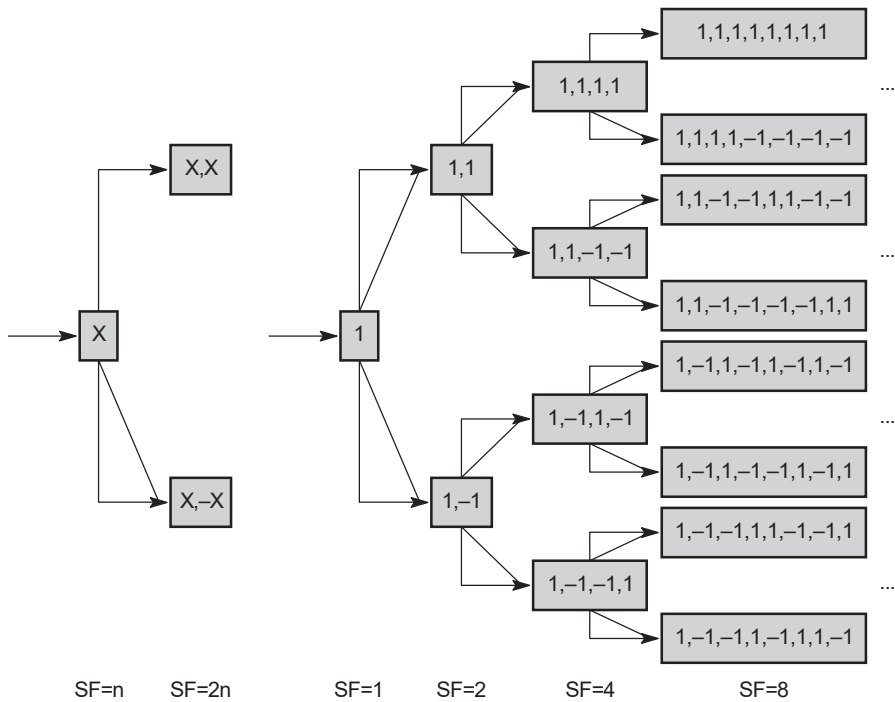


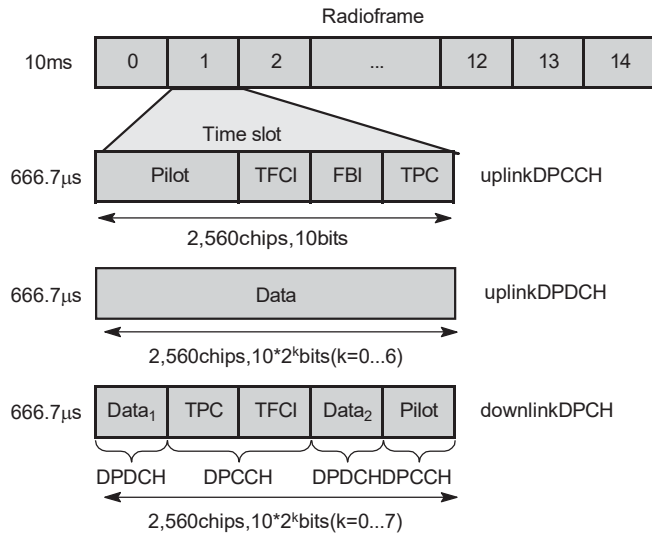
Figure 4.27
OVSF codetree used
for orthogonal spreading

4.4.3.1 UTRA-FDD(W-CDMA)

The FDD mode for UTRA uses **wideband CDMA (W-CDMA)** with direct sequence spreading. As implied by FDD, uplink and downlink use different frequencies. A mobile station in Europe sends via the uplink using a carrier between 1920 and 1980 MHz, the base station uses 2110 to 2170 MHz for the downlink (see Figure 4.22). Figure 4.28 shows a radio frame comprising 15 time slots. Time slots in W-CDMA are not used for user separation but to support periodic functions (note that this is in contrast to, e.g., GSM, where time slots are used to separate users!). A radio frame consists of 38,400 chips and has a duration of 10 ms. Each time slot consists of 2,560 chips, which roughly equals $666.6 \mu s$.¹¹ The occupied bandwidth per W-CDMA channel is 4.4 to 5 MHz (channel spacing can be varied to avoid interference between channels of different operators). These 5 MHz bands of the spectrum have been sold in many countries using an auction or a beauty contest. In Germany, the FDD spectrum was sold for over 50 billion Euros during an auction! But that was at a time when marketing people tried to convince everyone that UMTS would bring

¹¹ Early version of W-CDMA specified a chipping rate of 4.096 M chip/s and 16 time slots per frame. This was changed during the harmonization process which was necessary to avoid patent conflicts and to enable devices that can handle different CDMA standards. The harmonization process is fostered by the operators harmonization group (OHG), which is an informal steering group of wireless operator companies promoting 3G harmonization. The OHG was founded in 1999.

Figure 4.28
UTRA FDD (W-CDMA)
frame structure



high-bandwidth application to any mobile device with high profits for all. Today, most people are much more realistic and know that data rates will be quite low in the beginning (150 kbit/s per user are realistic, 2 Mbit/s are not). The capacity of a cell under realistic assumptions (interference etc.), i.e., the sum of all data rates, will rather be 2 Mbit/s. To provide high data rates a lot of money has to be invested in the infrastructure: UTRA FDD requires at least twice as many base stations as GSM; cell diameters of 500 m will become common place. This shows clearly that this technology will not cover whole countries in the near future but cities and highway only. People in the countryside will have to rely on GSM/GPRS for many more years to come.

Back to the frame structure shown in Figure 4.28. Similar to GSM, UMTS defines many logical and physical channels, and their mapping. The figure shows three examples of physical channels as they are used for data transmission. Two physical channels are shown for the uplink.

- Dedicated physical data channel (DPDCH):** This channel conveys user or signaling data. The spreading factor of this channel can vary between 4 and 256. This directly translates into the data rate this channel can offer: 960 kbit/s (spreading factor 4, 640 bits per slot, 15 slots per frame, 100 frames per second), 480, 240, 120, 60, 30, and 15 kbit/s (spreading factor 256). This also shows one of the problems of using OVSF for spreading: only certain multiples of the basic data rate of 15 kbit/s can be used. If, for example, 250 kbit/s are needed the device has to choose 480 kbit/s, which wastes bandwidth. In each connection in layer 1 it can have between zero and six DPDCHs. This results in a theoretical maximum data rate of 5,740 kbit/s (UMTS describes UEs with a maximum of 1,920 kbit/s only). Table 4.7 shows typical user data rates together with the required data rates on the physical channels.

User data rate [kbit/s]	12.2 (voice)	64	144	384
DPDCH [kbit/s]	60	240	480	960
DPCCH [kbit/s]	15	15	15	15
Spreading	64	16	8	4

Table 4.7 Typical UTRA-FDD uplink data rates

- Dedicated physical control channel (DPCCH):** In each connection layer 1 needs exactly one DPCCH. This channel conveys control data for the physical layer only and uses the constant spreading factor 256. The **pilot** is used for channel estimation. The **transport format combination identifier (TFCI)** specifies the channels transported within the DPDCHs. Signaling for a soft handover is supported by the **feedback information field (FBI)**. The last field, **transmit power control (TPC)** is used for controlling the transmission power of a sender. Power control is performed in each slot, thus 1,500 power control cycles are available per second. Tight power control is necessary to mitigate near-far effects as explained in chapter 2. Six different DPCCH bursts have been defined which differ in the size of the fields.
- Dedicated physical channel (DPCH):** The downlink time multiplexes control and user data. Spreading factors between 4 and 512 are available. Again, many different burst formats (17 altogether) have been defined which differ in the size of the fields shown in Figure 4.28. The available data rates for data channels (DPDCH) within a DPCH are 6 (SF=512), 24, 51, 90, 210, 432, 912, and 1,872 kbit/s (SF=4).

While no collisions can occur on the downlink (only the base station sends on the downlink), medium access on the uplink has to be coordinated. A **physical random access channel (PRACH)** is used for this purpose. UTRA-FDD defines 15 random access slots within 20 ms; within each access slot 16 different access preambles can be used for random access. Using slotted Aloha, a UE can access an access slot by sending a preamble. The UE starts with the lowest available transmission power to avoid interfering with other stations. If no positive acknowledgement is received, the UE tries another slot and another preamble with the next higher power level (power ramping). The number of available access slots can be defined per cell and is transmitted via a broadcast channel to all UEs.

A UE has to perform the following steps during the **search for a cell** after power on:

- Primary synchronization:** A UE has to synchronize with the help of a 256 chip primary synchronization code. This code is the same for all cells and helps to synchronize with the time slot structure.

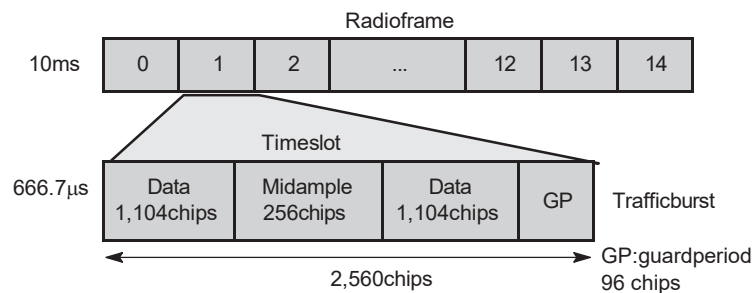
- **Secondary synchronization:** During this second phase the UE receives a secondary synchronization code which defines the group of scrambling codes used in this cell. The UE is now synchronized with the frame structure.
- **Identification of the scrambling code:** The UE tries all scrambling codes within the group of codes to find the right code with the help of a correlator. After these three steps the UE can receive all further data over a broadcast channel.

4.4.3.2 UTRA-TDD(TD-CDMA)

The second UTRA mode, UTRA-TDD, separates up and downlink in time using a radio frame structure similar to FDD. 15 slots with 2,560 chips per slot form a radio frame with a duration of 10 ms. The chipping rate is also 3.84 Mchip/s. To reflect different user needs in terms of data rates, the TDD frame can be **symmetrical or asymmetrical**, i.e., the frame can contain the same number of uplink and downlink slots or any arbitrary combination. The frame can have only one **switching point** from uplink to downlink or several switching points. However, at least one slot must be allocated for the uplink and downlink respectively.

The system can change the spreading factor (1, 2, 4, 8, 16) as a function of the desired data rate. Using the burst type shown in Figure 4.29 results in data rates of 6,624, 3,312, 1,656, 828, and 414 kbit/s respectively (if all slots are used for data transmission). The figure shows a burst of type 2 which comprises two data fields of 1,104 chips each. Spreading is applied to these data fields only. Additionally, a **midamble** is used for training and channel estimation. As TDD uses the same scrambling codes for all stations, the stations must be tightly synchronized and the spreading codes are available only once per slot. This results in a maximum number of 16 simultaneous sending stations. To loosen the tight synchronization a little bit, a **guard period (GP)** has been introduced at the end of each slot. Due to the tight synchronization and the use of orthogonal codes, a simpler power control scheme with less power control cycles (e.g., 100 per second) is sufficient.

Figure 4.29
UTRA-TDD(TD-CDMA)
frame structure



UTRA TDD occupies 5 MHz bandwidth per channel as UTRA FDD does per direction (FDD needs 2x5 MHz). Compared to the license for FDD, TDD was quite cheap. Germany paid less than €300 million. Figure 4.22 shows the location of the spectrum for this UMTS mode, but it is unclear to what extent this system will be deployed. The coverage per cell is even less than using FDD, UEs must not move too fast – this sounds like the characteristics of WLANs which are currently deployed in many places.

4.4.4 UTRAN

Figure 4.30 shows the basic architecture of the UTRA network (UTRAN; 3GPP, 2002b). This consists of several **radio network subsystems (RNS)**. Each RNS is controlled by **radio network controller (RNC)** and comprises several components that are called node B. An RNC in UMTS can be compared with the BSC; a node B is similar to a BTS. Each node B can control several antennas which make a radio cell. The mobile device, UE, can be connected to one or more antennas as will subsequently be explained in the context of handover. Each RNC is connected with the core network (CN) over the interface I_u (similar to the role of the A interface in GSM) and with a node B over the interface I_{ub} . A new interface, which has no counterpart in GSM, is the interface I_{ur} connecting two RNCs with each other. The use of this interface is explained together with the UMTS handover mechanisms.

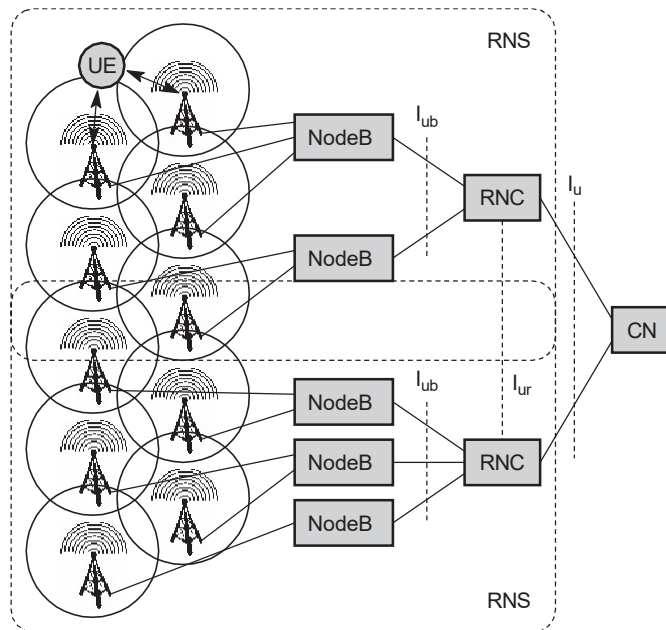


Figure 4.30
Basic architecture
of the UTRAN network

4.4.4.1 Radio network controller

An RNC in UMTS has a broad spectrum of tasks as listed in the following:

- **Call admission control:** It is very important for CDMA systems to keep the interference below a certain level. The RNC calculates the traffic within each cell and decides, if additional transmissions are acceptable or not.
- **Congestion control:** During packet-oriented data transmission, several stations share the available radio resources. The RNC allocates bandwidth to each station in a cyclic fashion and must consider the QoS requirements.
- **Encryption/decryption:** The RNC encrypts all data arriving from the fixed network before transmission over the wireless link (and vice versa).
- **ATM switching and multiplexing, protocol conversion:** Typically, the connections between RNCs, node Bs, and the CN are based on ATM. An RNC has to switch the connections to multiplex different data streams. Several protocols have to be converted – this is explained later.
- **Radio resource control:** The RNC controls all radio resources of the cells connected to it via a node B. This task includes interference and load measurements. The priorities of different connections have to be obeyed.
- **Radio bearer setup and release:** An RNC has to set-up, maintain, and release a logical data connection to a UE (the so-called UMTS radio bearer).
- **Code allocation:** The CDMA codes used by a UE are selected by the RNC. These codes may vary during a transmission.
- **Power control:** The RNC only performs a relatively loose power control (the outer loop). This means that the RNC influences transmission power based on interference values from other cells or even other RNCs. But this is not the tight and fast power control performed 1,500 times per second. This is carried out by a node B. This outer loop of power control helps to minimize interference between neighbouring cells or controls the size of a cell.
- **Handover control and RNS relocation:** Depending on the signal strengths received by UEs and node Bs, an RNC can decide if another cell would be better suited for a certain connection. If the RNC decides for handover it informs the new cell and the UE as explained in subsection 4.4.6. If a UE moves further out of the range of one RNC, a new RNC responsible for the UE has to be chosen. This is called RNS relocation.
- **Management:** Finally, the network operator needs a lot of information regarding the current load, current traffic, error states etc. to manage its network. The RNC provides interfaces for this task, too.

4.4.4.2 Node B

The name node B was chosen during standardization until a new and better name was found. However, no one came up with anything better so it remained. A node B connects to one or more antennas creating one or more cells (or sectors in GSM speak), respectively. The cells can either use FDD or TDD

or both. An important task of a node B is the inner loop power control to mitigate near-far effects. This node also measures connection qualities and signal strengths. A node B can even support a special case of handover, also called softer handover which takes place between different antennas of the same node B (see section 4.4.6).

4.4.4.3 User equipment

The UE shown in Figure 4.30 is the counterpart of several nodes of the architecture.

- As the counterpart of a node B, the UE performs signal quality measurements, inner loop power control, spreading and modulation, and rate matching.
- As a counterpart of the RNC, the UE has to cooperate during handover and cell selection, performs encryption and decryption, and participates in the radio resource allocation process.
- As a counterpart of the CN, the UE has to implement mobility management functions, performs bearer negotiation, or requests certain services from the network.

This list of tasks of a UE, which is not at all exhaustive, already shows the complexity such a device has to handle. Additionally, users also want to have organizers, games, cameras, operating system etc. and the stand-by time should be high.

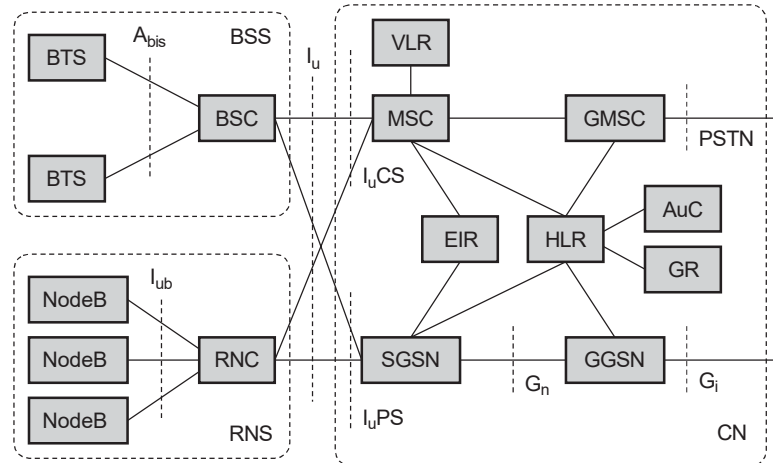
4.4.5 Core network

Figure 4.31 shows a high-level view of the UMTS release 99 core network architecture together with a UTRAN RNS and a GSM BSS (see section 4.1). This shows the evolution from GSM/GPRS to UMTS. The core network (CN) shown here is basically the same as already explained in the context of GSM (see Figure 4.4) and GPRS (see Figure 4.16). The **circuit switched domain (CSD)** comprises the classical circuit switched services including signaling. Resources are reserved at connection setup and the GSM components MSC, GMSC, and VLR are reused. The CSD connects to the RNS via a part of the I_q interface called I_qCS . The CSD components can still be part of a classical GSM network connected to a BSS but need additional functionalities (new protocols etc.).

The **packet switched domain (PSD)** uses the GPRS components SGSN and GGSN and connects to the RNS via the I_qPS part of the I_q interface. Both domains need the data-bases EIR for equipment identification and HLR for location management (including the AuC for authentication and GR for user specific GPRS data).

Reusing the existing infrastructure helps to save a lot of money and may convince many operators to use UMTS if they already use GSM. The UMTS industry pushes their technology with the help of the market dominance of GSM. This is basically the same as cdma2000, which is an evolution of cdmaOne. The real flexible core network comes with releases 5 and 6, where the GSM

Figure 4.31
UMTS core network
together with 3G RNS
and 2G BSS



circuit switched part is being replaced by an all-IP core. Chapter 11 presents this idea in the context of 4G networks. It is not yet clear when this replacement of GSM will take place as many questions are still open (quality of service and security being the most important).

Figure 4.32 shows the protocol stacks of the user planes of the circuit switched and packet switched domains, respectively. The CS uses the **ATM adaptation layer 2 (AAL2)** for user data transmission on top of ATM as transport technology. The RNC in the UTRAN implements the radio link control (RLC) and the MAC layer, while the physical layer is located in the node B. The **AAL2 segmentation and reassembly layer (SAR)** is, for example, used to segment data packets received from the RLC into small chunks which can be transported in ATM. AAL2 and ATM has been chosen, too, because these protocols can transport and multiplex low bit rate voice data streams with low jitter and latency (compared to the protocols used in the PSD).

In the PSD several more protocols are needed. Basic data transport is performed by different lower layers (e.g., ATM with AAL5, frame relay). On top of these lower layers UDP/IP is used to create a UMTS internal IP network. All packets (e.g., IP, PPP) destined for the UE are encapsulated using the **GPRS tunneling protocol (GTP)**. The RNC performs protocol conversion from the combination GTP/UDP/IP into the **packet data convergence protocol (PDCP)**. This protocol performs header compression to avoid redundant data transmission using scarce radio resources. Comparing Figure 4.32 with Figure 4.17 (GPRS protocol reference model) shows a difference with respect to the tunnel. In UMTS the RNC handles the tunneling protocol GTP, while in GSM/GPRS GTP is used between an SGSN and GGSN only. The BSC in GSM is not involved in IP protocol processing.

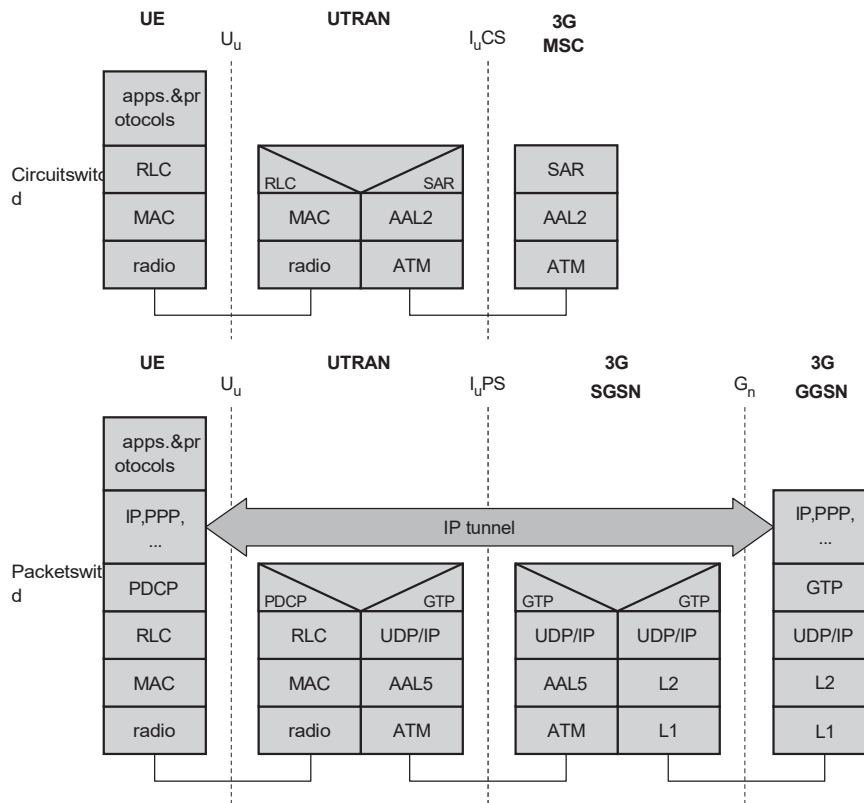


Figure 4.32 Userplane protocol stacks (circuit and packet switched)

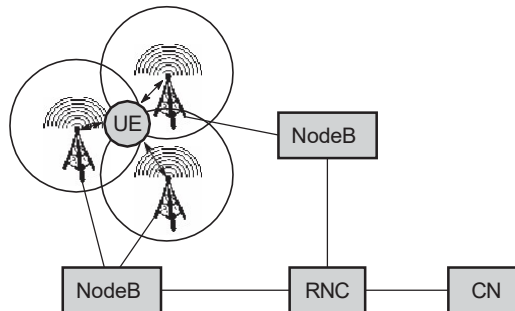
The **radio layer** (physical layer) depends on the UTRA mode (see sections 4.4.3.1 and 4.4.3.2). The **medium access control (MAC)** layer coordinates medium access and multiplexes logical channels onto transport channels. The MAC layer also helps to identify mobile devices and may encrypt data. The **radio link control (RLC)** layer offers three different transport modes. The **acknowledged mode** transfer uses ARQ for error correction and guarantees one-time in-order delivery of data packets. The **unacknowledged mode** transfer does not perform ARQ but guarantees at least one-time delivery of packets with the help of sequence numbers. The **transparent mode** transfer simply forwards MAC data without any further processing. The system then has to rely on the FEC which is always used in the radio layer. The RLC also performs segmentation and reassembly and flow control. For certain services the RLC also encrypts.

4.4.6 Handover

UMTS knows two basic classes of handovers:

- Hard handover:** This handover type is already known from GSM and other TDMA/FDMA systems. Switching between different antennas or different systems is performed at a certain point in time. **UTRA TDD** can only use this type. Switching between TDD cells is done between the slots of different frames. **Interfrequency handover**, i.e., changing the carrier frequency, is a hard handover. Receiving data at different frequencies at the same time requires a more complex receiver compared to receiving data from different sources at the same carrier frequency. Typically, all **inter system handovers** are hard handovers in UMTS. This includes handovers to and from GSM or other IMT-2000 systems. A special type of handover is the handover to a satellite system (inter-segment handover), which is also a hard handover, as different frequencies are used. However, it is unclear what technology will be used for satellite links if it will ever come. To enable a UE to listen into GSM or other frequency bands, UMTS specifies **acompressed mode** transmission for UTRA FDD. During this mode a UE stops all transmission. To avoid data loss, either the spreading factor can be lowered before and after the break in transmission (i.e., more data can be sent in shorter time) or less data is sent using different coding schemes.
- Soft handover:** This is the real new mechanism in UMTS compared to GSM and is only available in the FDD mode. Soft handovers are well known from traditional CDMA networks as they use **macrodiversity**, a basic property of CDMA. As shown in Figure 4.33, a UE can receive signals from up to three different antennas, which may belong to different node Bs. Towards the UE the RNC splits the data stream and forwards it to the node Bs. The UE combines the received data again. In the other direction, the UE simply sends its data which is then received by all node Bs involved. The RNC combines the data streams received from the node Bs. The fact that a UE receives data from different antennas at the same time makes a handover soft. Moving from one cell to another is a smooth, not an abrupt process.

Figure 4.33
Macro-diversity
supporting soft
handovers



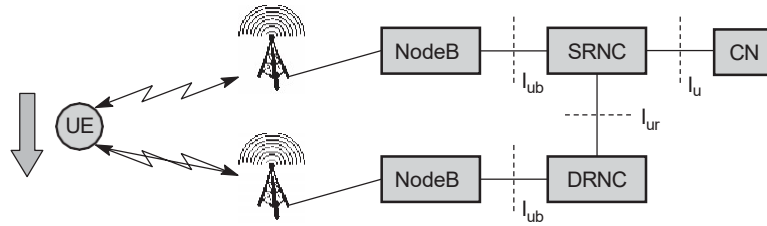


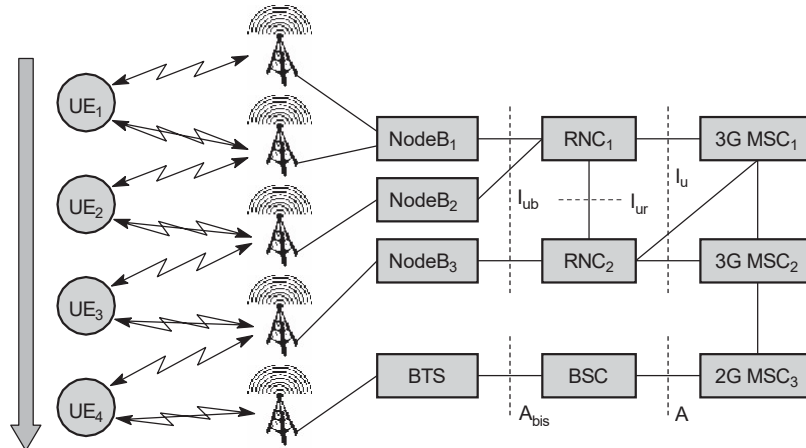
Figure 4.34
Serving RNC and
drift RNC

Macro-diversity makes the transmission more robust with respect to fast fading, multi-path propagation, and shading. If one path is blocked by an obstacle the chances are good that data can still be received using another antenna. During a soft handover a UE receives power control commands from all involved node Bs. The UE then lowers transmission power as long as it still receives a command to lower the power. This avoids interference if, for example, the UE is in the transmission area of two antennas, one close, one further away. Without the above mechanism the UE's signal may be too strong when listening to the antenna further away. The lower the interference a UE introduces into a cell, the higher the capacity. Without this control, cell breathing would be even more problematic than it already is in CDMA networks.

As soft handover is not supported by the CN, all mechanisms related to soft handover must be located within UTRAN. Figure 4.34 shows a situation where a soft handover is performed between two node Bs that do not belong to the same RNC. In this case one RNC controls the connection and forwards all data to and from the CN. If the UE moves in the example from the upper cell to the lower cell, the upper RNC acts as a **serving RNC (SRNC)** while the other is the **drift RNC (DRNC)**. (If the whole RNS is considered, the terms are serving RNS and drift RNS, respectively.) The SRNC forwards data received from the CN to its node B and to the DRNC via the I_{ur} interface (splitting). This mechanism does not exist in, e.g., GSM. Data received by the lower node B is forwarded by the DRNC to the SRNC. The SRNC combines both data streams and forwards a single stream of data to the CN. The CN does not notice anything from the simultaneous reception. If the UE moves further down and drops out of the transmission area of the upper node B, two RNCs reserve resources for data transmission, SRNC and DRNC, although none of SRNC's node Bs transmit data for this UE. To avoid wasting resources, SRNC relocation can be performed. This involves the CN so is a hard handover.

Figure 4.35 gives an overview of several common handover types in a combined UMTS/GSM network (UMTS specifies ten different types which include soft and hard handover). The combination of a UTRA-FDD/GSM device will be the most common case in the beginning as coverage of 3G networks will be poor.

Figure 4.35
Overview of different
handover types



- **Intra-node B, intra-RNC:** UE₁ moves from one antenna of node B₁ to another antenna. This type of handover is called **softer handover**. In this case node B₁ performs combining and splitting of the data streams.
- **Inter-node B, intra-RNC:** UE₂ moves from node B₁ to node B₂. In this case RNC₁ supports the soft handover by combining and splitting data.
- **Inter-RNC:** When UE₃ moves from node B₂ to node B₃, two different types of handover can take place. The **internal inter-RNC** handover is not visible for the CN, as described in Figure 4.34. RNC₁ can act as SRNC, RNC₂ will be the DRNC. The CN will communicate via the same interface I_u all the time. As soon as a relocation of the interface I_u takes place (relocation of the controlling RNC), the handover is called an **external inter-RNC** handover. Communication is still handled by the same MSC₁, but the external handover is now a hard handover.
- **Inter-MSC:** It could be also the case that MSC₂ takes over and performs a hard handover of the connection.
- **Inter-system:** UE₄ moves from a 3G UMTS network into a 2G GSM network. This hard handover is important for real life usability of the system due to the limited 3G coverage in the beginning.

4.5 Summary

This chapter has, for the most part, presented GSM as the most successful second generation digital cellular network. Although GSM was primarily designed for voice transmission, the chapter showed the evolution toward a more data-oriented transfer via HSCSD and GPRS. This evolution also includes the transition from a circuit-switched network to a packet-switched system that comes closer to the internet model. Other systems presented include DECT, the

digital standard for cordless phones, and TETRA, a trunked radio system. DECT can be used for wireless data transmission on a campus or indoors, but also for wireless local loops (WLL). For special scenarios, e.g., emergencies, trunked radio systems such as TETRA can be the best choice. They offer a fast connection setup (even within communication groups) and can work in an ad hoc network, i.e., without a base station.

The situation in the US is different from Europe. Based on the analog AMPS system, the US industry developed the TDMA system IS-54 that adds digital traffic channels. IS-54 uses dual mode mobile phones and incorporates several GSM ideas, such as, associated control channels, authentication procedures using encryption, and mobile assisted handover (called handoff). The Japanese PDC system was designed using many ideas in IS-54.

The next step, IS-136, includes digital control channels (IS-54 uses analog AMPS control channels) and is more efficient. Now fully digital phones can be used, several additional services are offered, e.g., voice mail, call waiting, identification, group calling, or SMS. IS-136 is also called North American TDMA (NA-TDMA) or Digital AMPS (D-AMPS) and operates at 800 and 1,900 MHz. Enhancements of D-AMPS/IS-136 toward IMT-2000 include advanced modulation techniques for the 30 kHz radio carrier, shifting data rates up to 64 kbit/s (first phase, called 136+). The second phase, called 136HS (High Speed) comprises a new air interface specification based on the EDGE technology.

IS-95 (promoted as cdmaOne) is based on CDMA, which is a completely different medium access method. Before deployment, the system was proclaimed as having many advantages over TDMA systems, such as its much higher capacity of users per cell, e.g., 20 times the capacity of AMPS. Today, CDMA providers are making more realistic estimates of around five times as many users. IS-95 offers soft handover, avoiding the GSM ping-pong effect (Wong, 1997). However, IS-95 needs precise synchronization of all base stations (using GPS satellites which are military satellites, so are not under control of the network provider), frequent power control, and typically, dual mode mobile phones due to the limited coverage. The basic ideas of CDMA have been integrated into most 3G systems.

This chapter also presented an overview of current and future third generation systems. UMTS, a proposal of operators and companies involved in the GSM business, was discussed in more detail. This standard is more an evolutionary approach than a revolution. To avoid even higher implementation costs, UMTS tries to reuse as much infrastructure as possible while introducing new services and higher data rates based on CDMA technology. The initial installations will basically use the GSM/GPRS infrastructure and offer only moderate data rates. The initial capacity of a UMTS cell is approximately 2 Mbit/s; cell diameters are in the order of 500 m. UMTS will be used to offload GSM networks and to offer enhanced data rates in cities as a first step. Future releases aim to replace the infrastructure by an (almost) all-IP network. These ideas will be presented together with a look at fourth generation systems in chapter 11. It

is quite clear that it will take a long time before 3G services are available in many places. It took GSM 10 years to become the most successful 2G mobile communication system. A similar period of time will be needed for 3G systems to succeed. Meanwhile, customers will need multiple mode phones offering, e.g., GSM 900/1800/1900 and UMTS UTRA-FDD services. It is not clear if and when UTRA-TDD will succeed. Providers already using cdmaOne will take the evolutionary path via cdma2000 1x toward the 3G system cdma2000 1x EV-DO. Several tests have already been conducted for 3G satellite services in the MSS spectrum (e.g., satellite based multicast, Nussli, 2002). However, right now many companies will wait before investing money in satellite services (see chapter 5). The main problem of multi-mode systems is the inter-system handover. While this chapter introduces handover scenarios within UMTS and GSM, and between GSM and UMTS, even more complex scenarios could comprise wireless LANs (see chapter 7) or other packet-oriented networks (Pahlavan, 2000).

4.6 Review exercises

- 1 Name some key features of the GSM, DECT, TETRA, and UMTS systems. Which features do the systems have in common? Why have the three older different systems been specified? In what scenarios could one system replace another? What are the specific advantages of each system?
- 2 What are the main problems when transmitting data using wireless systems that were made for voice transmission? What are the possible steps to mitigate the problems and to raise efficiency? How can this be supported by billing?
- 3 Which types of different services does GSM offer? Give some examples and reasons why these services have been separated.
- 4 Compared to the TCHs offered, standard GSM could provide a much higher data rate (33.8 kbit/s) when looking at the air interface. What lowers the data rates available to a user?
- 5 Name the main elements of the GSM system architecture and describe their functions. What are the advantages of specifying not only the radio interface but also all internal interfaces of the GSM system?
- 6 Describe the functions of the MS and SIM. Why does GSM separate the MS and SIM? How and where is user-related data represented/stored in the GSM system? How is user data protected from unauthorized access, especially over the air interface? How could the position of an MS (not only the current BTS) be localized? Think of the MS reports regarding signal quality.
- 7 Looking at the HLR/VLR database approach used in GSM - how does this architecture limit the scalability in terms of users, especially moving users?
- 8 Why is a new infrastructure needed for GPRS, but not for HSCSD? Which components are new and what is their purpose?

- 9 What are the limitations of a GSM cell in terms of diameter and capacity (voice, data) for the traditional GSM, HSCSD, GPRS? How can the capacity be increased?
- 10 What multiplexing schemes are used in GSM and for what purpose? Think of other layers apart from the physical layer.
- 11 How is synchronization achieved in GSM? Who is responsible for synchronization and why is it so important?
- 12 What are the reasons for the delays in a GSM system for packet data traffic? Distinguish between circuit-switched and packet-oriented transmission.
- 13 Where and when can collisions occur while accessing the GSM system? Compare possible collisions caused by data transmission in standard GSM, HSCSD, and GPRS.
- 14 Why and when are different signaling channels needed? What are the differences?
- 15 How is localization, location update, roaming, etc. done in GSM and reflected in the data bases? What are typical roaming scenarios?
- 16 Why are so many different identifiers/addresses (e.g., MSISDN, TMSI, IMSI) needed in GSM? Give reasons and distinguish between user-related and system-related identifiers.
- 17 Give reasons for a handover in GSM and the problems associated with it. What are the typical steps for handover, what types of handover can occur? Which resources need to be allocated during handover for data transmission using HSCSD or GPRS respectively? What about QoS guarantees?
- 18 What are the functions of authentication and encryption in GSM? How is system security maintained?
- 19 How can higher data rates be achieved in standard GSM, how is this possible with the additional schemes HSCSD, GPRS, EDGE? What are the main differences of the approaches, also in terms of complexity? What problems remain even if the data rate is increased?
- 20 What limits the data rates that can be achieved with GPRS and HSCSD using real devices (compared to the theoretical limit in a GSM system)?
- 21 Using the best delay class in GPRS and a data rate of 115.2 kbit/s - how many bytes are in transit before a first acknowledgement from the receiver could reach the sender (neglect further delays in the fixed network and receiver system)? Now think of typical web transfer with 10 kbyte average transmission size - how would a standard TCP behave on top of GPRS (see chapters 9 and 10)? Think of congestion avoidance and its relation to other round-trip time. What changes are needed?
- 22 How much of the original GSM network does GPRS need? Which elements of the network perform the data transfer?
- 23 What are typical data rates in DECT? How are they achieved considering the TDMA frames? What multiplexing schemes are applied in DECT and for what purposes? Compare the complexity of DECT with that of GSM.

- 24 Who would be the typical users of a trunked radio system? What makes trunked radio systems particularly attractive for these user groups? What are the main differences to existing systems for that purpose? Why are trunked radio systems cheaper compared to, e.g., GSM systems for their main purposes?
 - 25 Summarize the main features of third-generation mobile phone systems. How do they achieve higher capacities and higher data rates? How does UMTS implement asymmetrical communication and different data rates?
 - 26 Compare the current situation of mobile phone networks in Europe, Japan, China, and North America. What are the main differences, what are efforts to find a common system or at least interoperable systems?
 - 27 What disadvantage does OVSF have with respect to flexible data rates? How does UMTS offer different data rates (distinguish between FDD and TDD mode)?
 - 28 How are different DPCHs from different UEs within one cell distinguished in UTRA FDD?
 - 29 Which components can perform combining/splitting at what handover situation? What is the role of the interface U_r ? Why can CDMA systems offer soft handover?
 - 30 How does UTRA-FDD counteract the near-far effect? Why is this not a problem in GSM?
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